

[54] **GENERAL PURPOSE CALCULATOR WITH CAPABILITY FOR PERFORMING INTERDISCIPLINARY BUSINESS CALCULATIONS**

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[52] U.S. Cl. **235/156**

[51] Int. Cl. **G06f 7/38**

[58] Field of Search 235/156, 159, 160, 164, 235/92 AC, 92 CP; 340/122.5; 444/1

[56] **References Cited**
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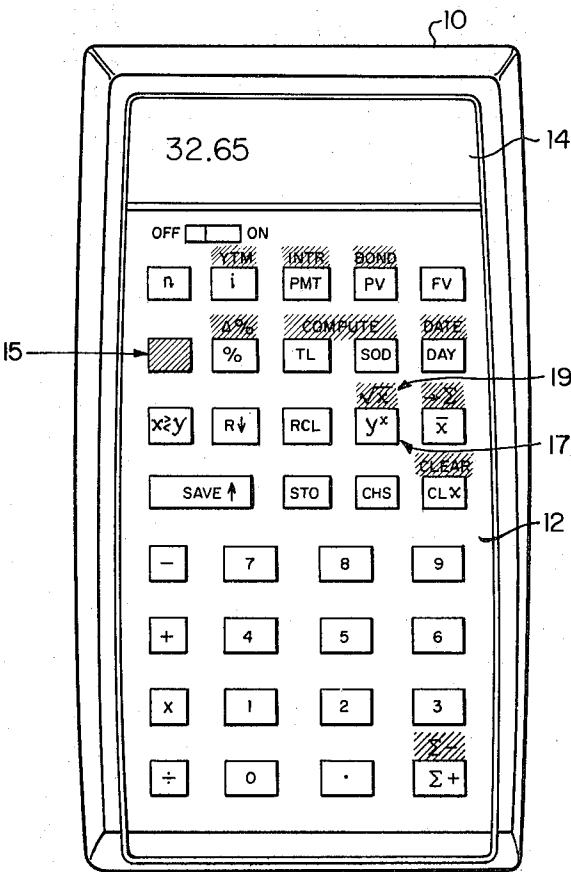
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Primary Examiner—Eugene G. Botz
Assistant Examiner—David H. Malzahn
Attorney, Agent, or Firm—Roland I. Griffin; F. David LaRiviere

[57] **ABSTRACT**

A battery-powered, hand-held, calculator employs MOS/LSI calculator circuits to perform arithmetic and financial calculations. Data and commands are input to the calculator from a keyboard having a prefix key to double the functions of selected keys. A 15-digit, seven-segment light emitter diode (LED) display serves as the output for the calculator. The calculator circuits include a read-only memory circuit in which the algorithms for performing the arithmetic and financial calculations are stored; a control and timing circuit for scanning the keyboard, retaining status information about the condition of the calculator or of an algorithm, and generating the next read-only memory address; and an arithmetic and register circuit containing an adder, a group of working registers, a group of data storage registers forming a stack for roll down operation, and a constant storage register. These circuits are interconnected by a multiple line buss system.

34 Claims, 36 Drawing Figures



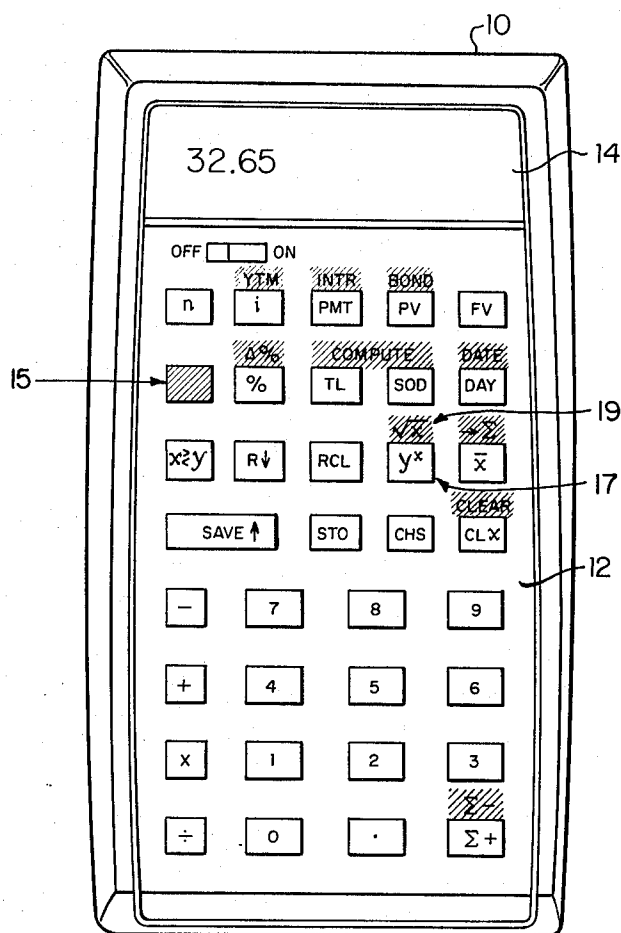


Figure 1

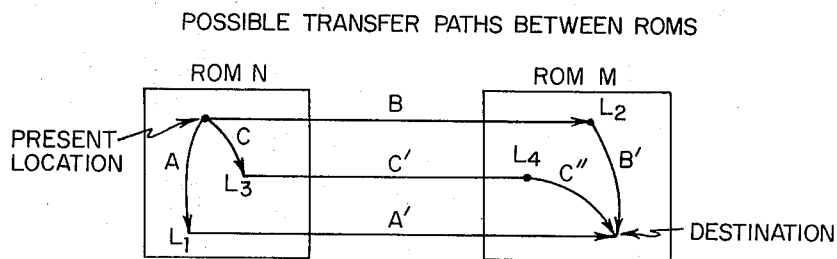


Figure 30

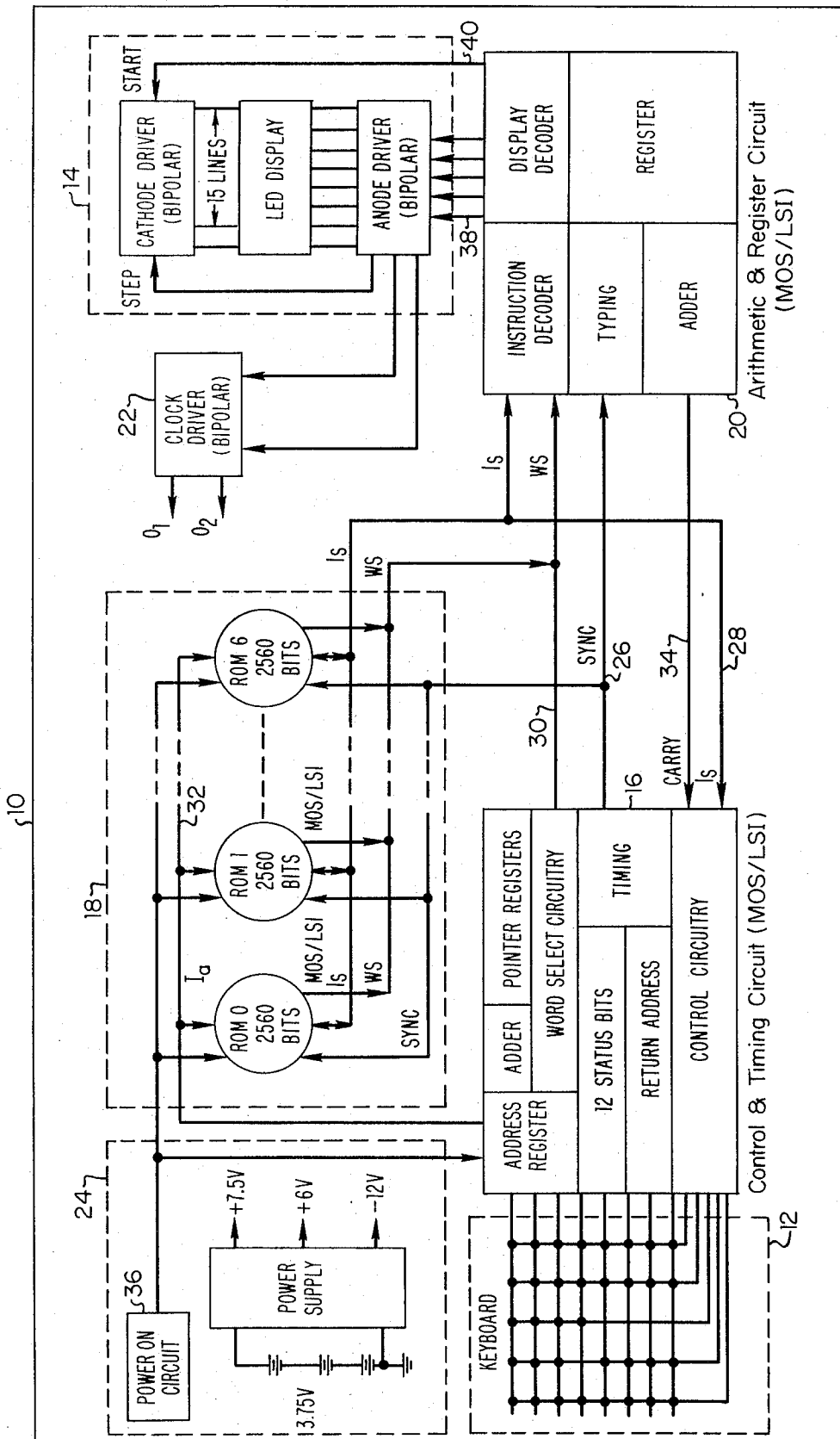


Figure 2

System Timing Signals

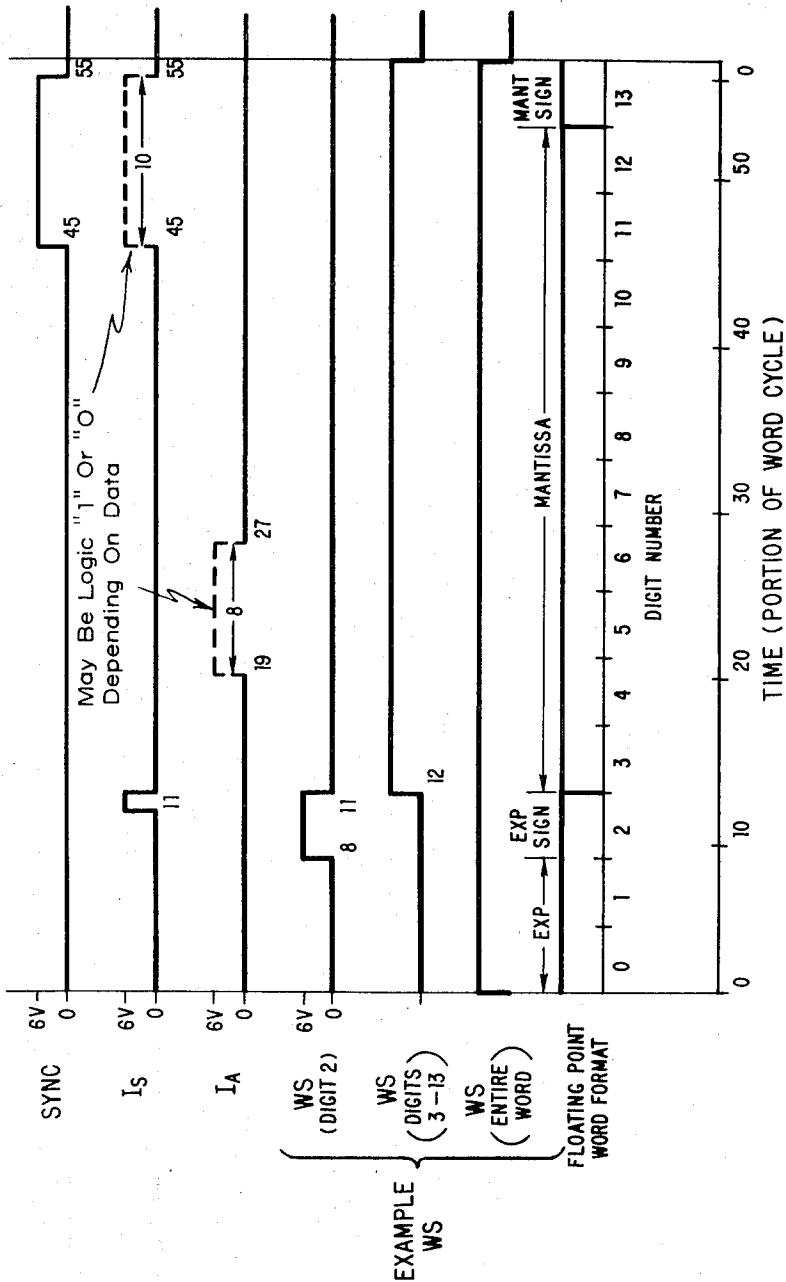


FIG.3

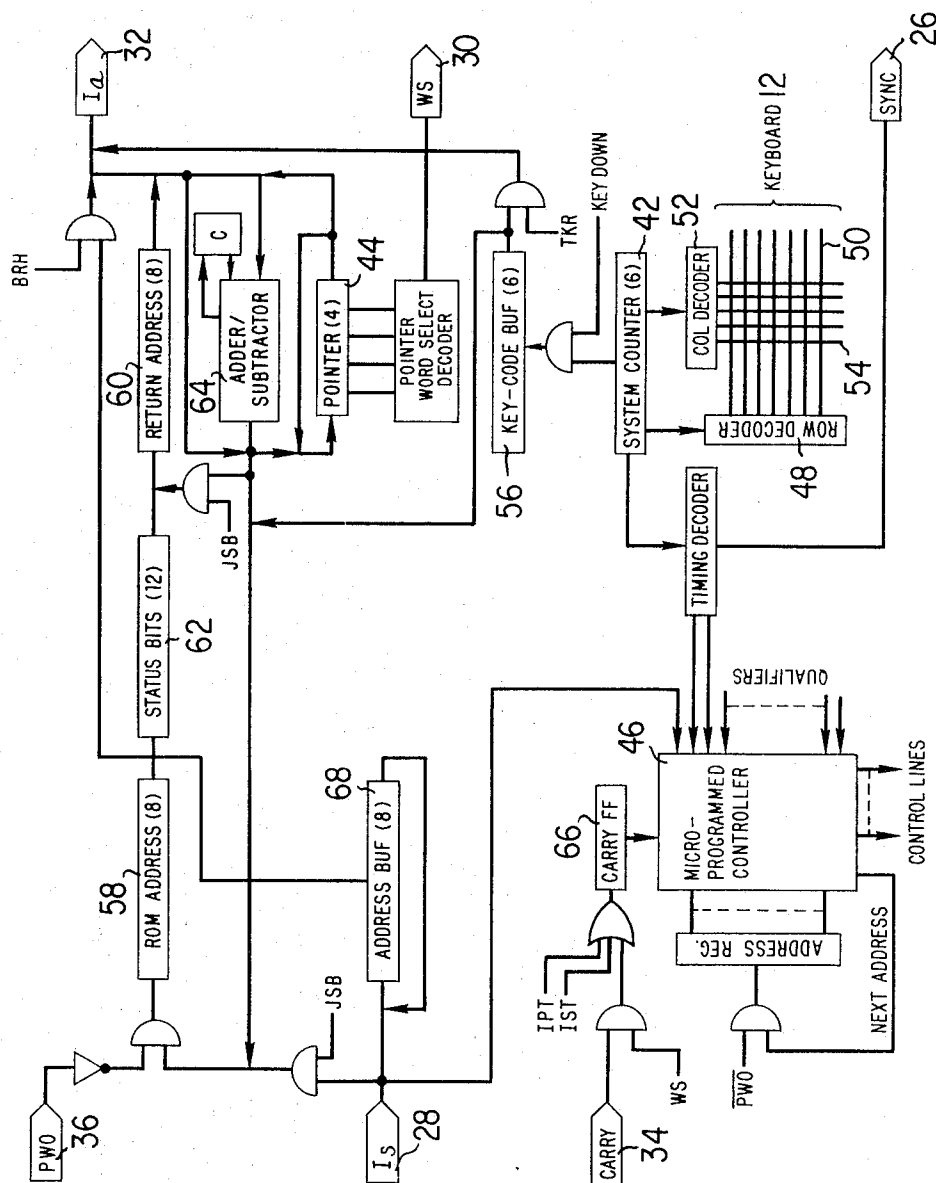


FIG. 4

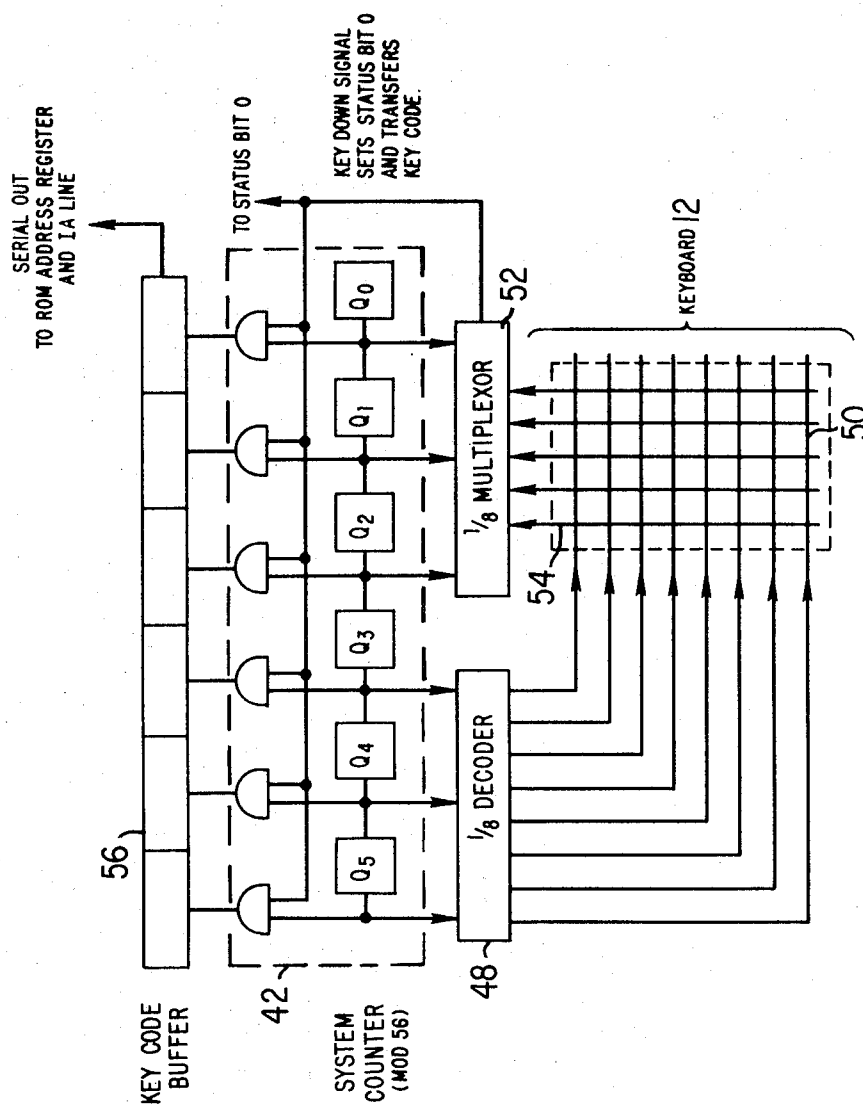


FIG. 5

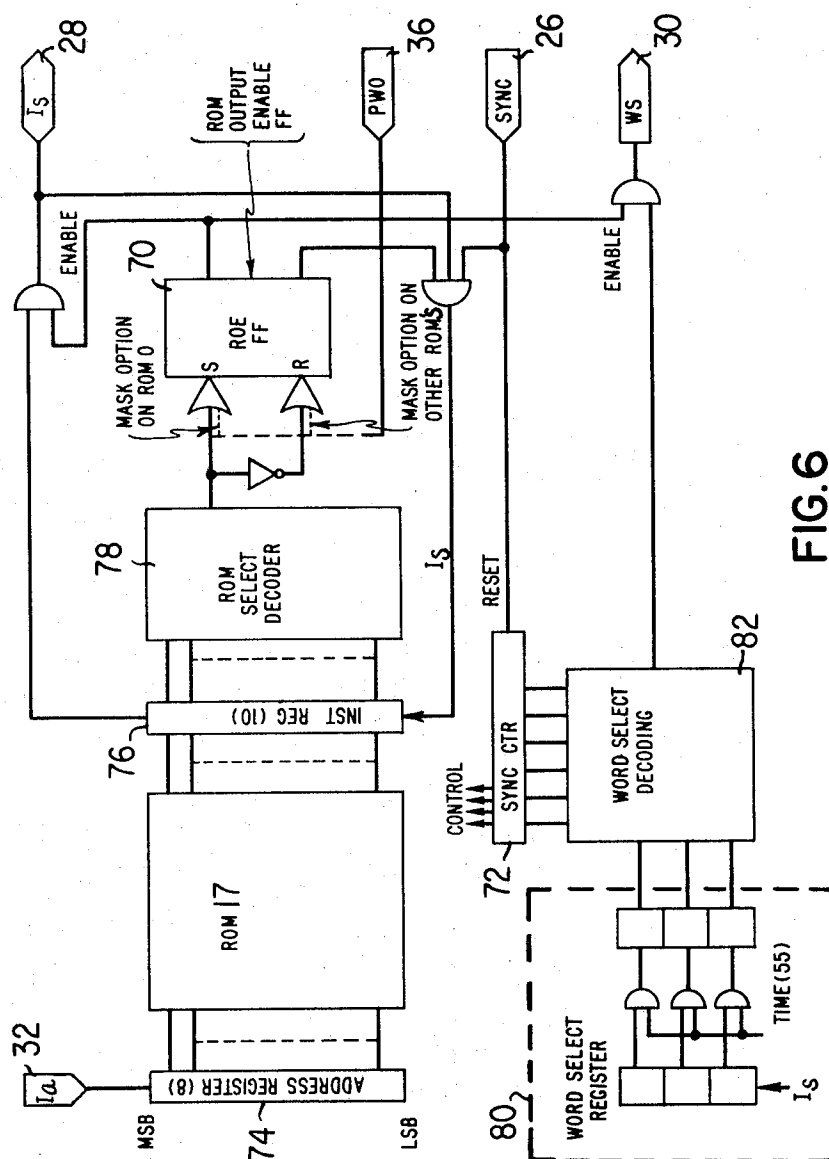


FIG. 6

Typical Address And Instruction Signals

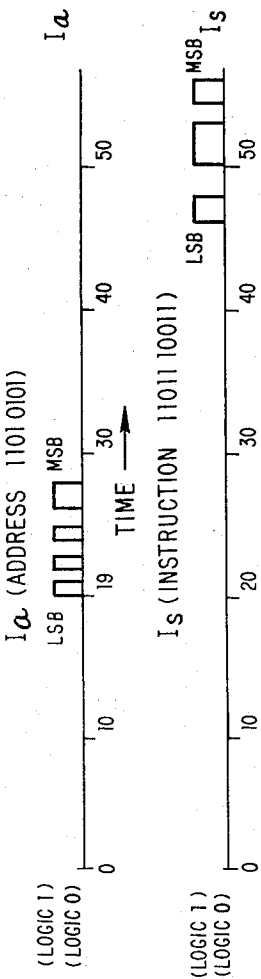


FIG. 7

Addressing And Readout Timing

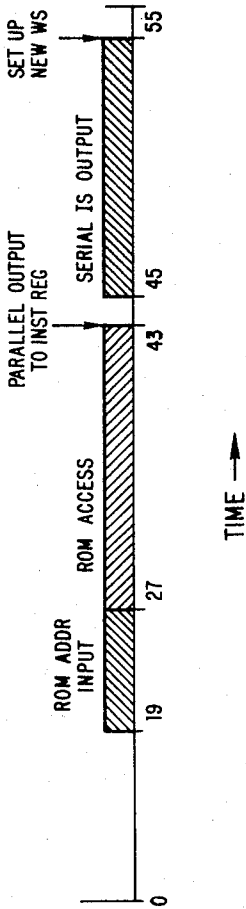


FIG. 8

Word Select Signals

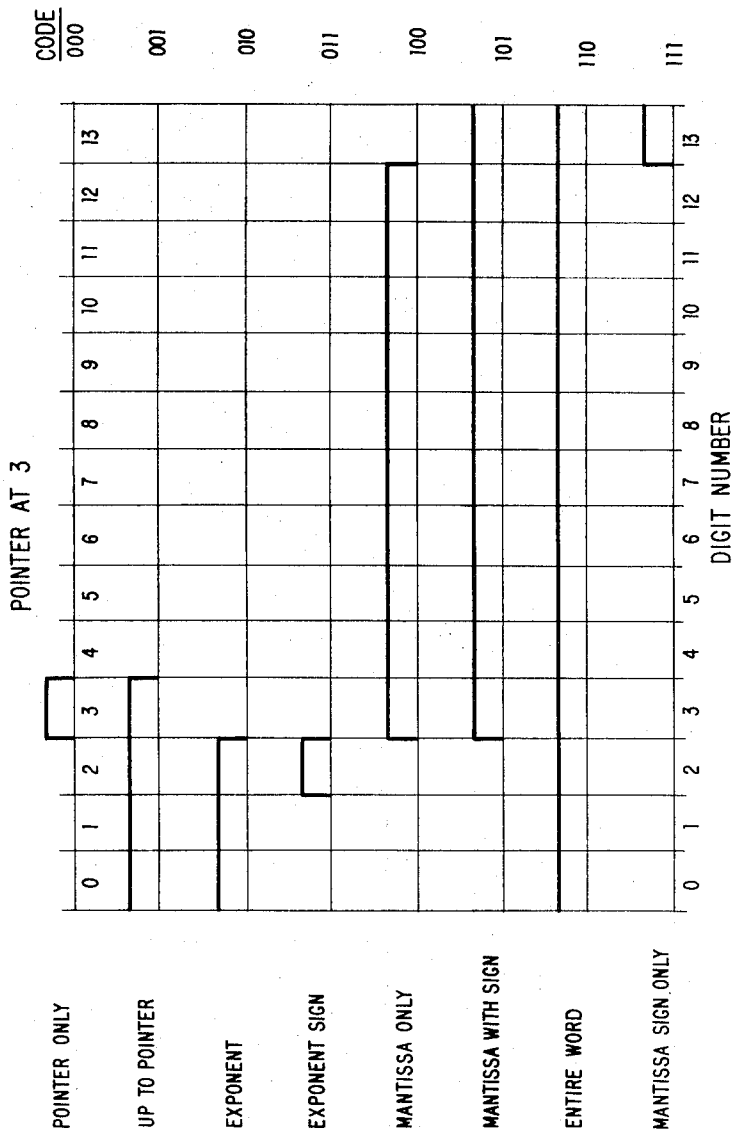


FIG. 9

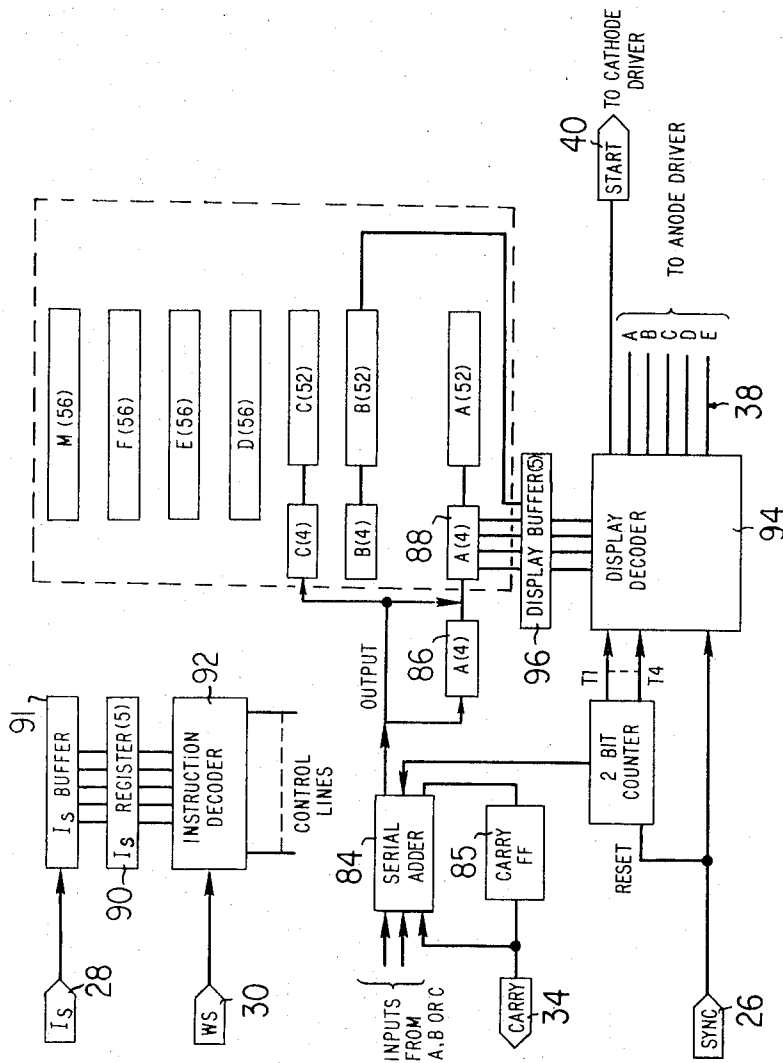


FIG. 10

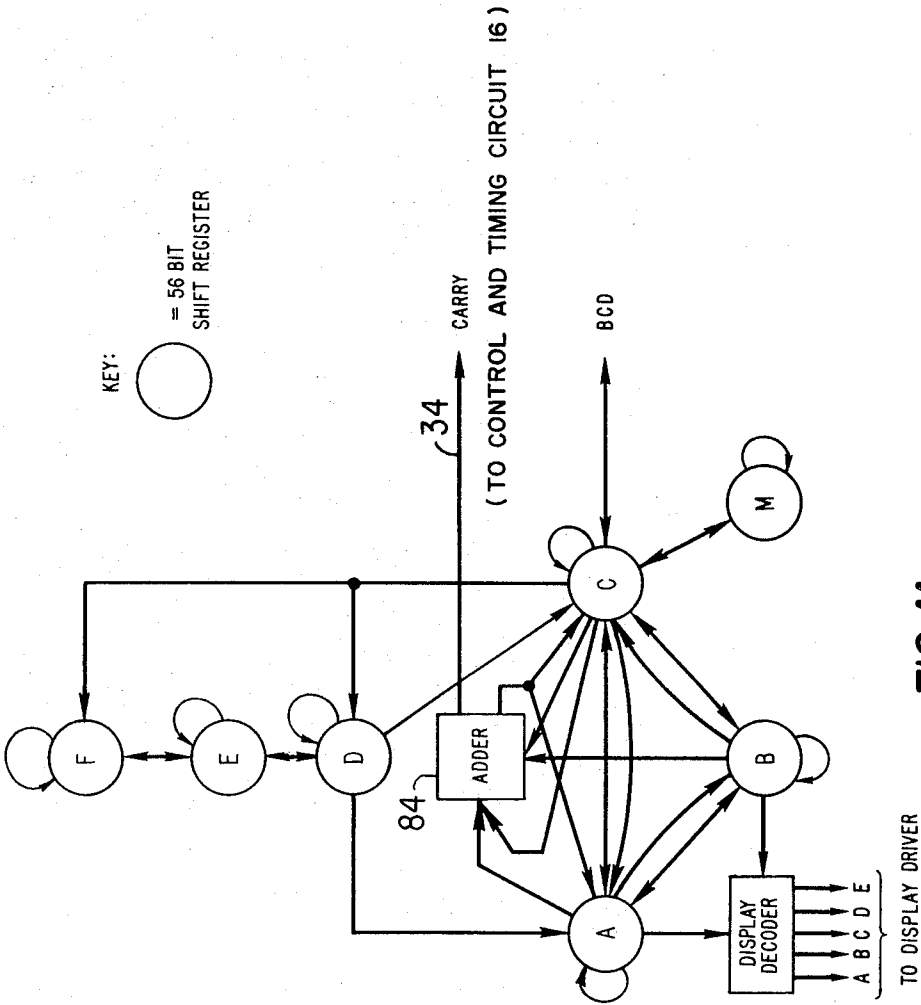


FIG. 11

Display Decoding From Arithmetic And Register Circuit

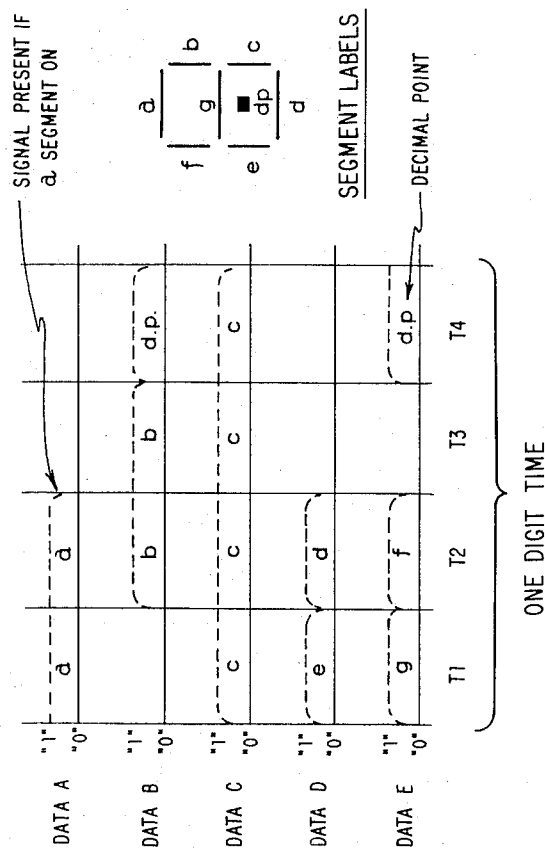


FIG.12

Example Output For Digit 9

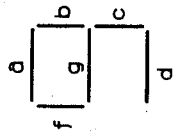
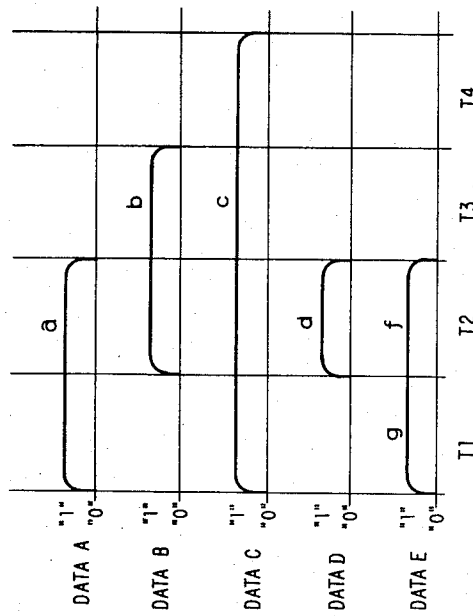


FIG.13

Start Signal

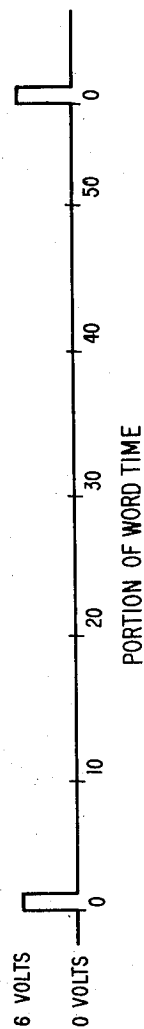


FIG. 14

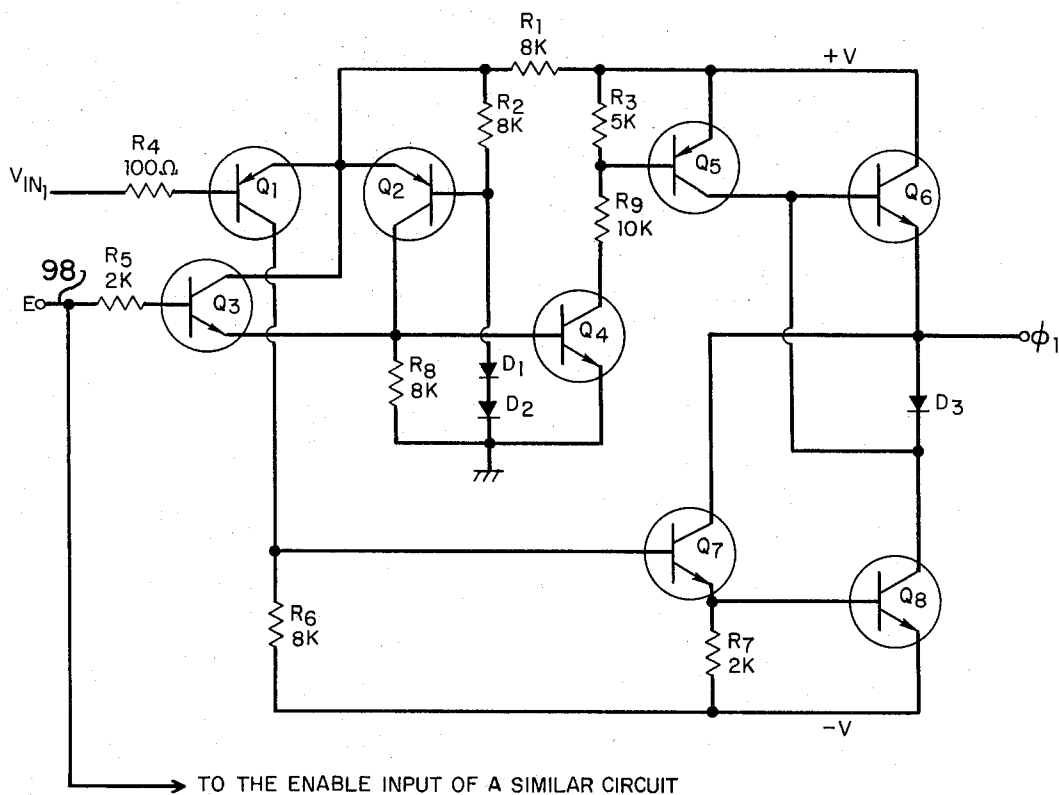


Figure 15

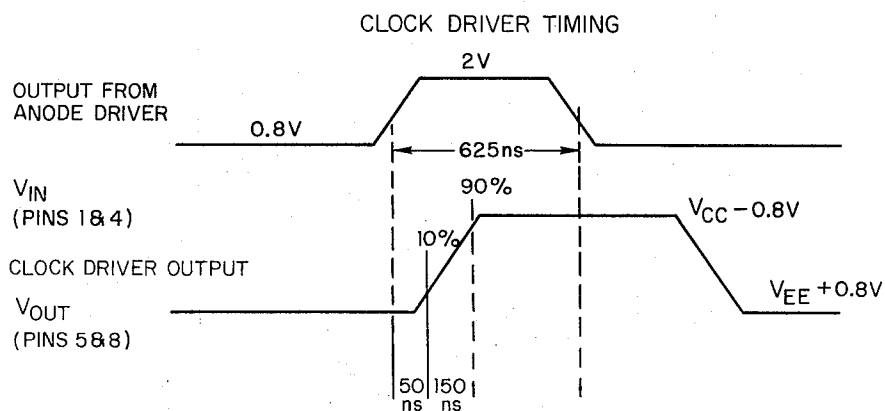


Figure 16

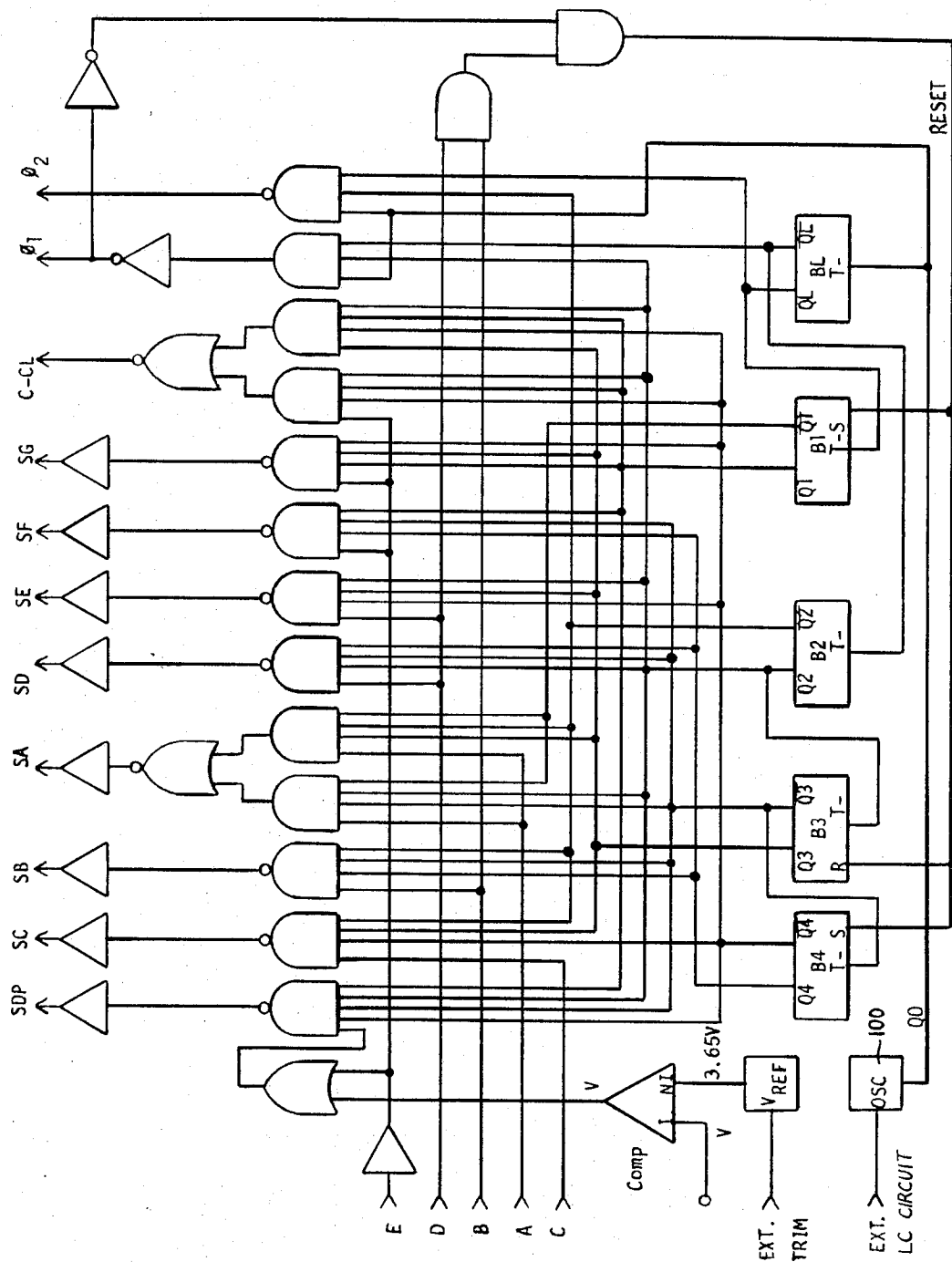


FIG.17

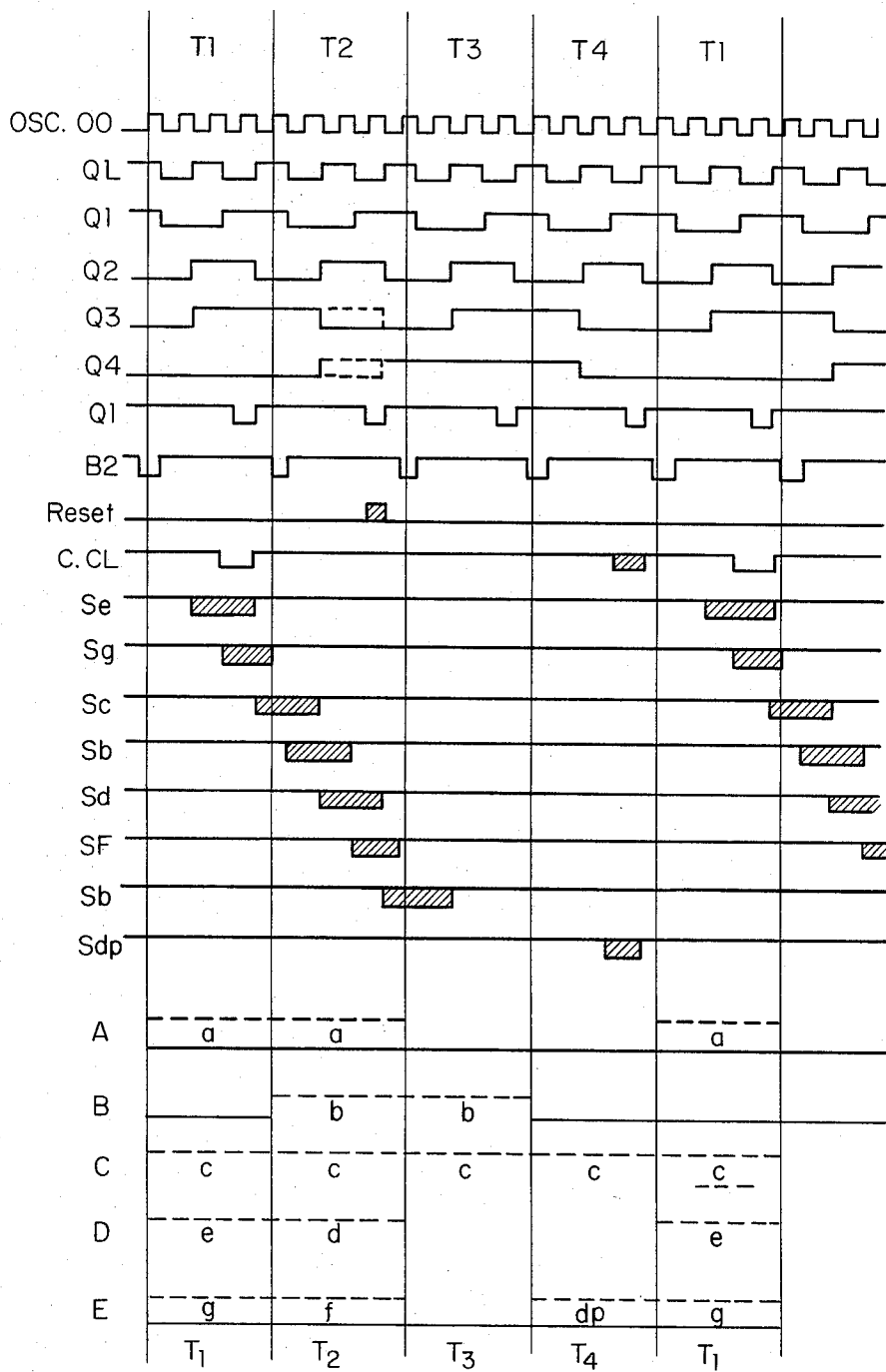


Figure 18

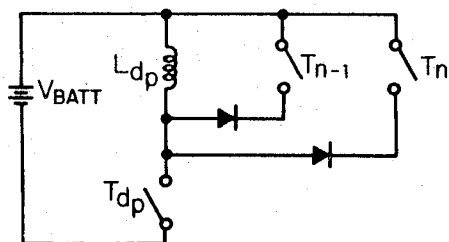


FIG. 19

Timing for Decimal Point Drive

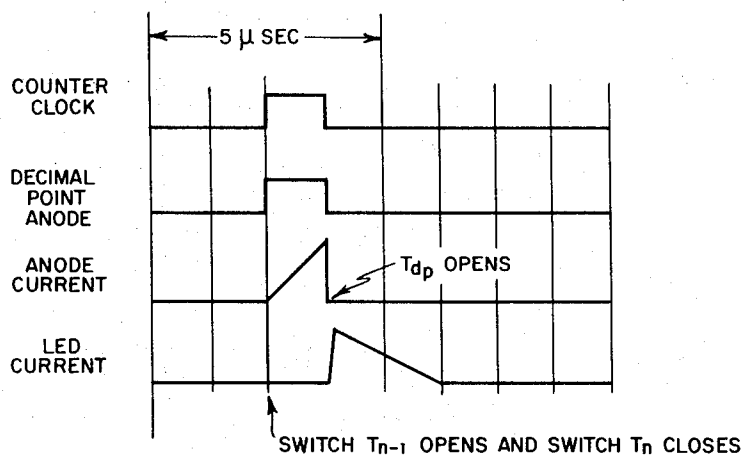


FIG. 20

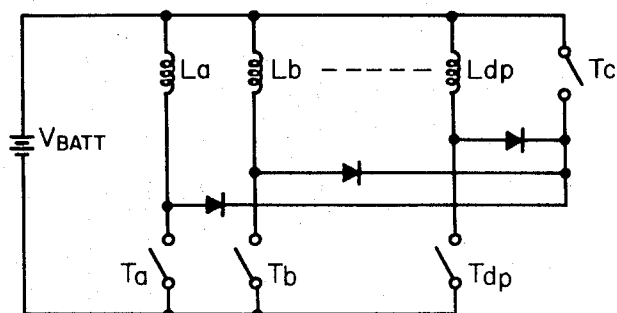


FIG. 21

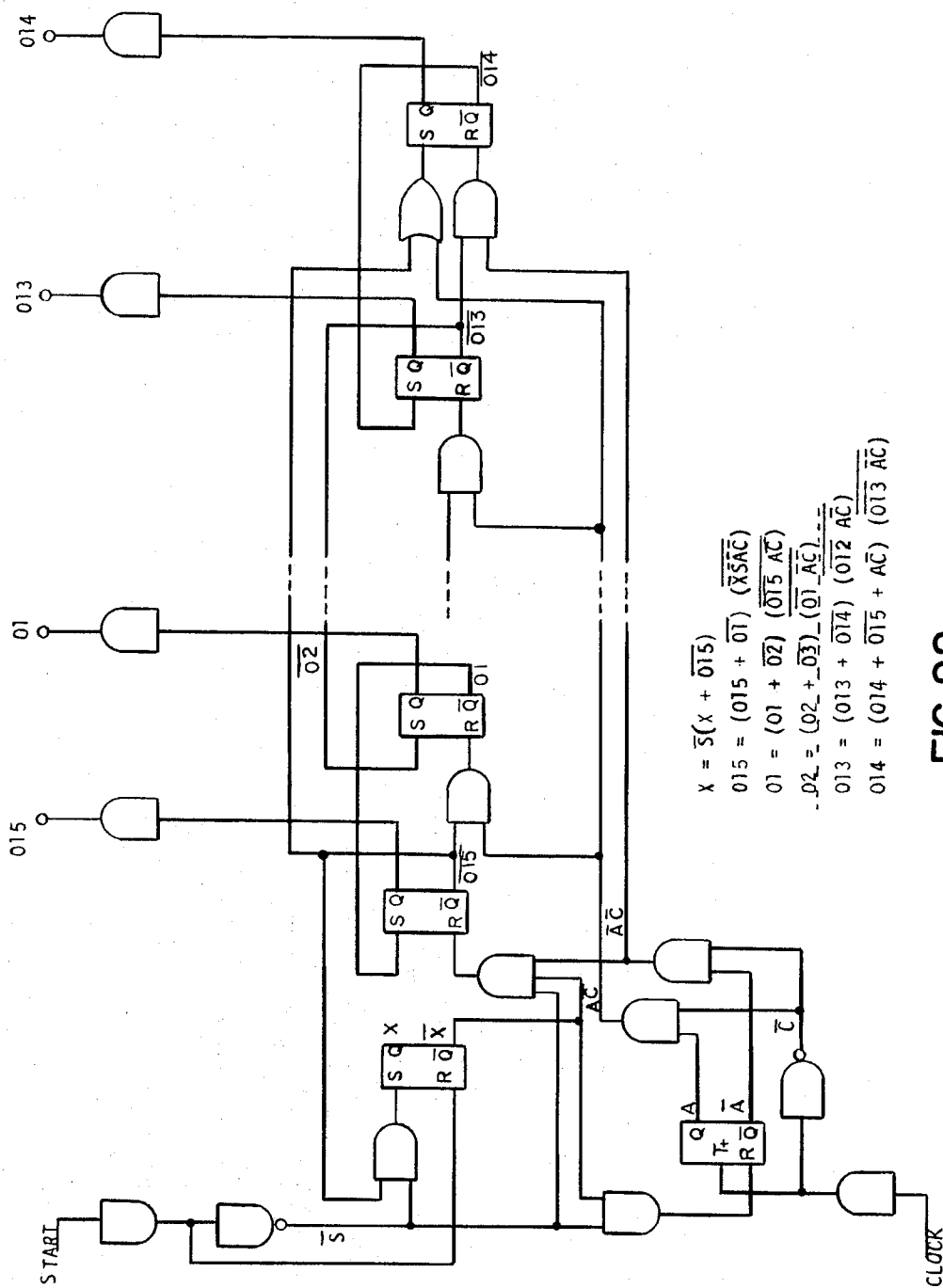


FIG. 22

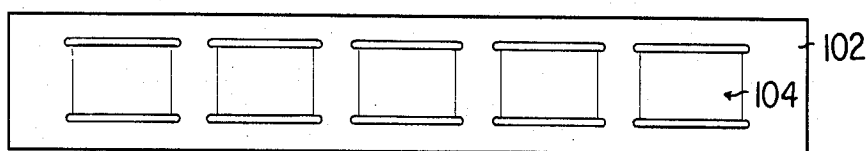


FIG. 23

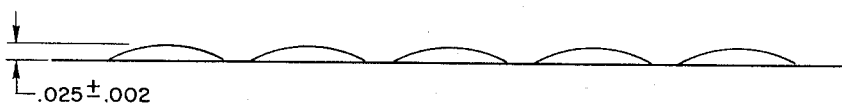


FIG. 24

Keyboard Force-Deflection Curve

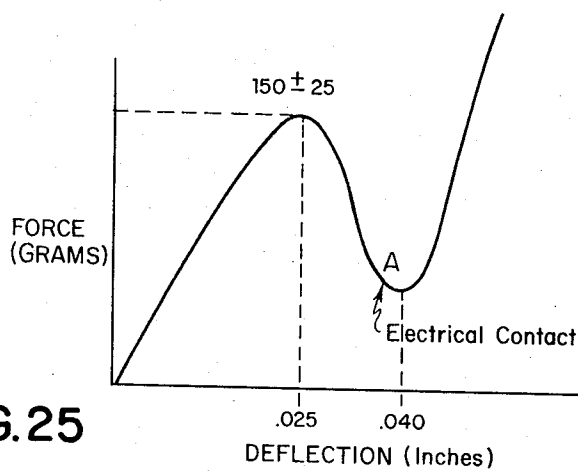


FIG. 25

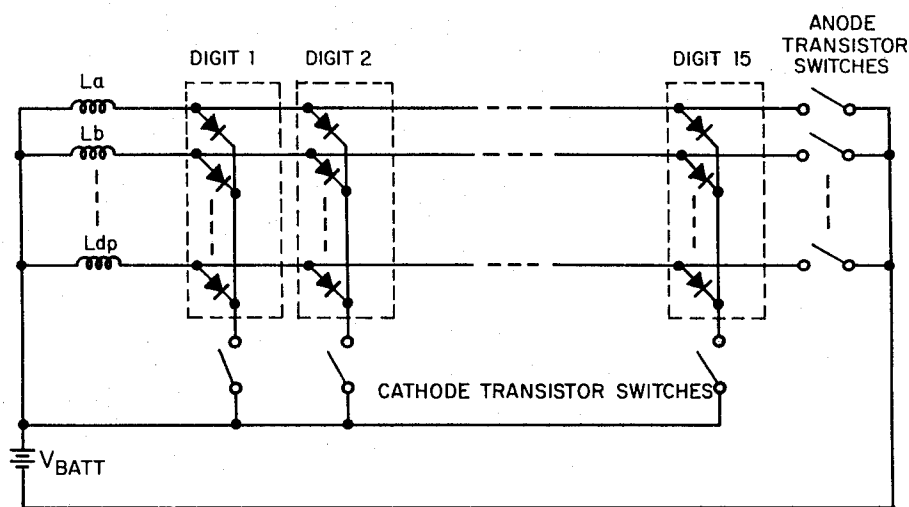


FIG. 26

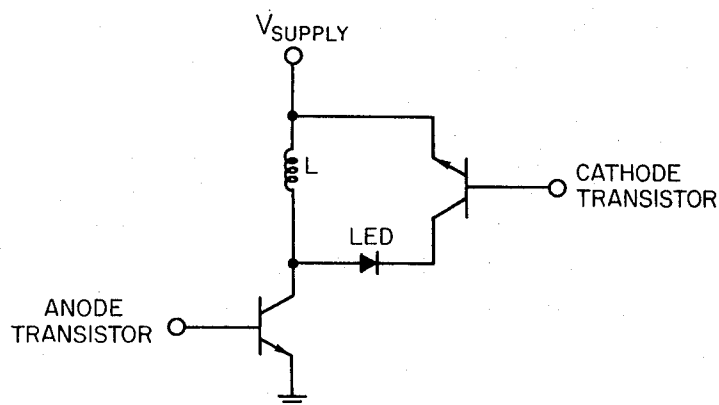
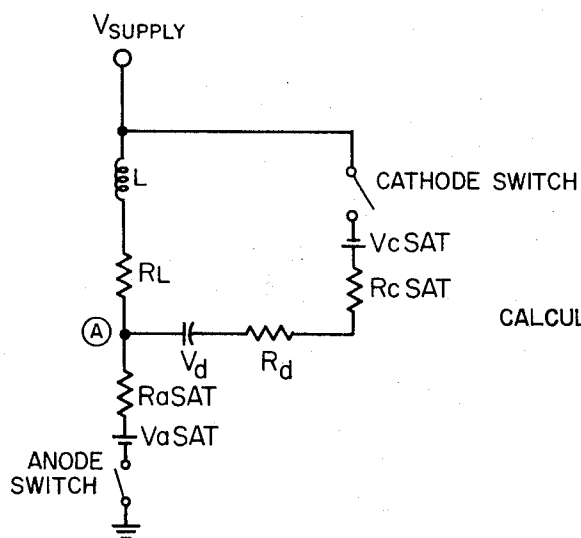


FIG. 27



CALCULATOR PARAMETERS

$L = 130 \mu h$
 $R_L = 4 \Omega$
 $R_{aSAT} = 1 \Omega$
 $R_d = 1 \Omega$
 $R_{cSAT} = 1 \Omega$
 $V_d = 1.7 \text{ Volts}$
 $V_{aSAT} = .3 \text{ Volts}$
 $V_{cSAT} = .45 \text{ Volts}$

FIG. 28

Inductor Current And LED Anode Voltages

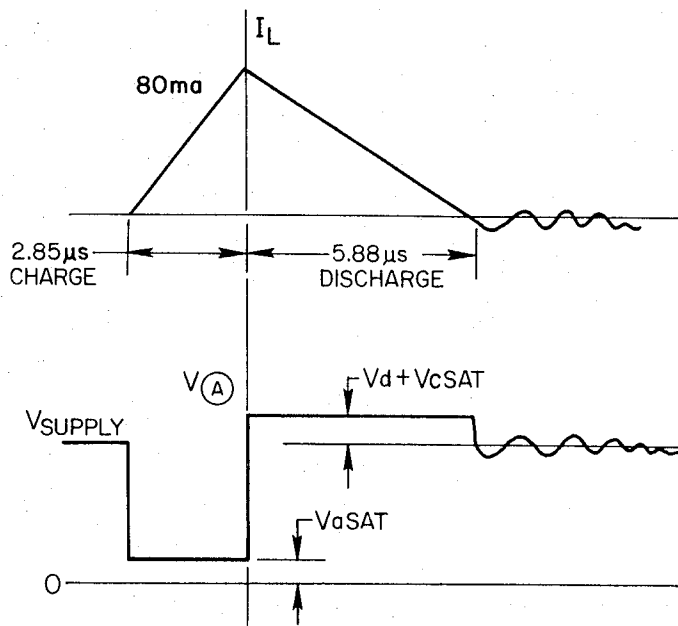


FIG. 29

FLOW DIAGRAM OF DISPLAY WAIT LOOP

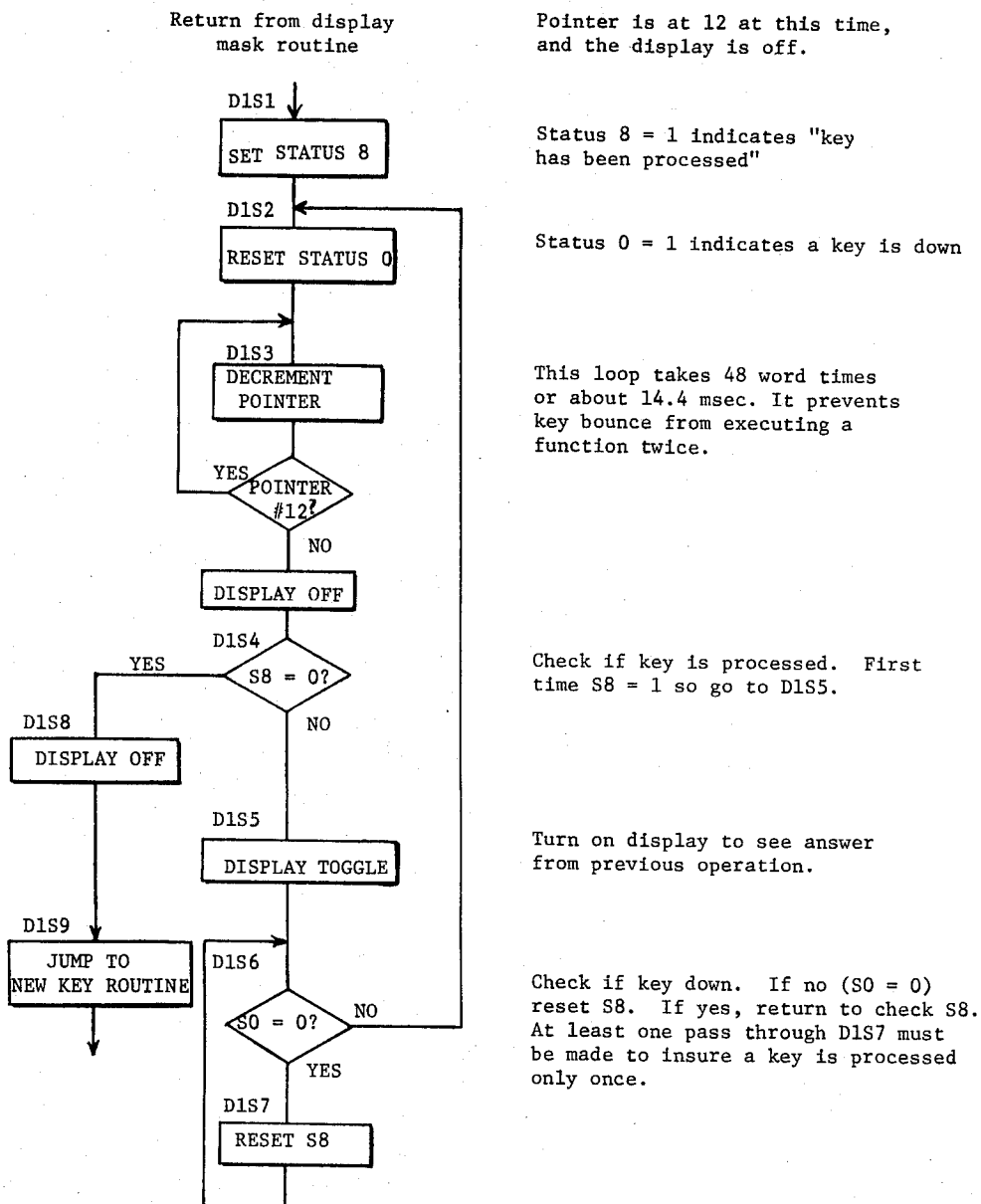


FIG. 31

6.1 PRICE OF A BOND (PV)

N - UNCOMPENSATED DAYS

R - INTEREST RATE %

C - COUPON RATE %

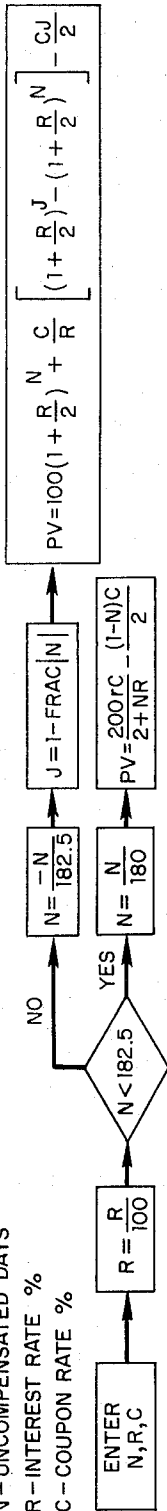


Fig. 33

6.2 YIELD TO MATURITY OF A BOND (R)

N - UNCOMPENSATED DAYS

C - COUPON RATE %

P - PRICE OF THE BOND

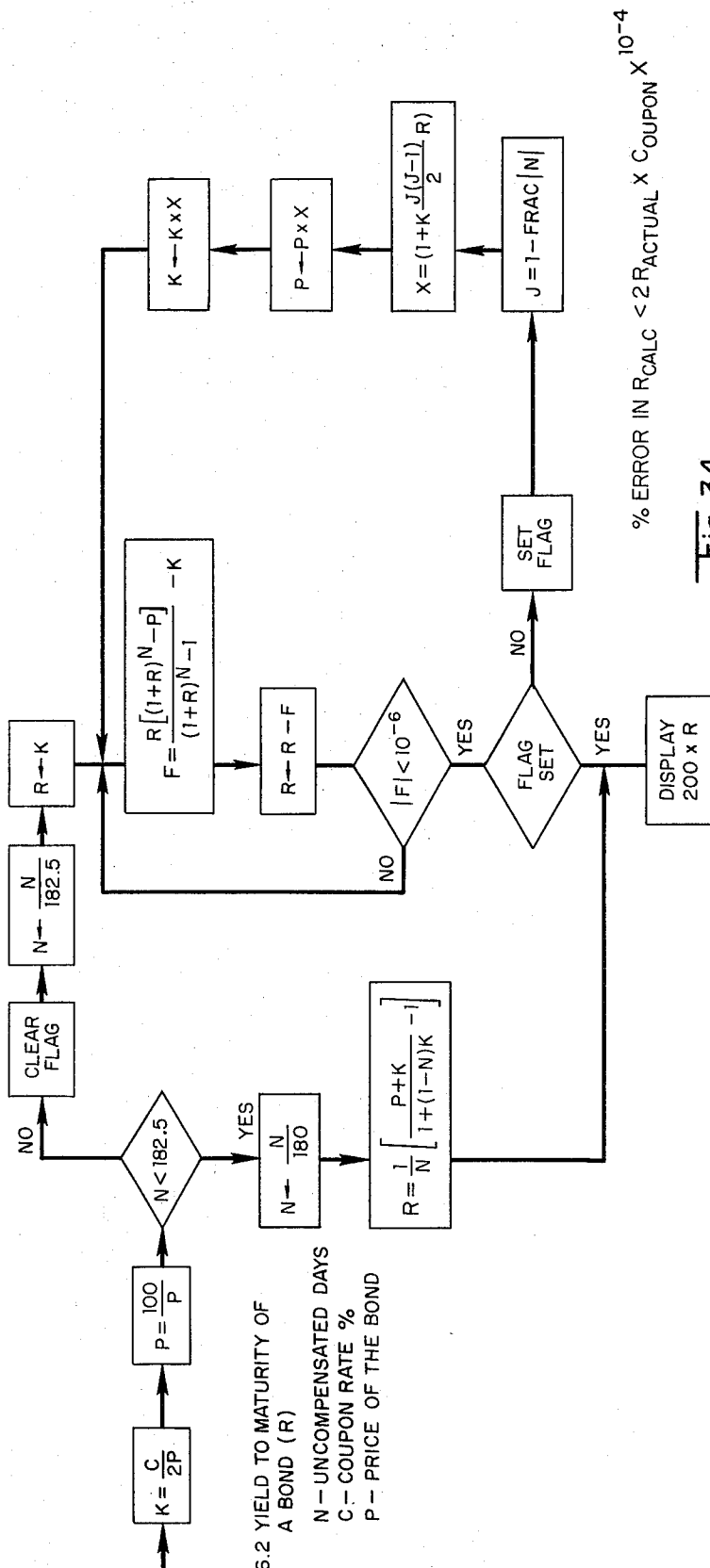
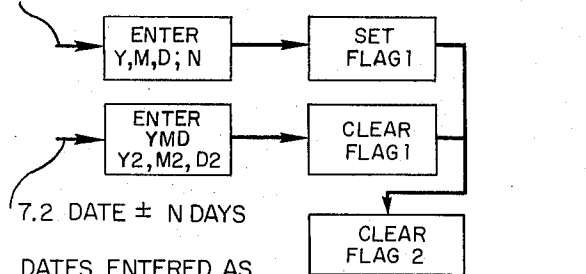


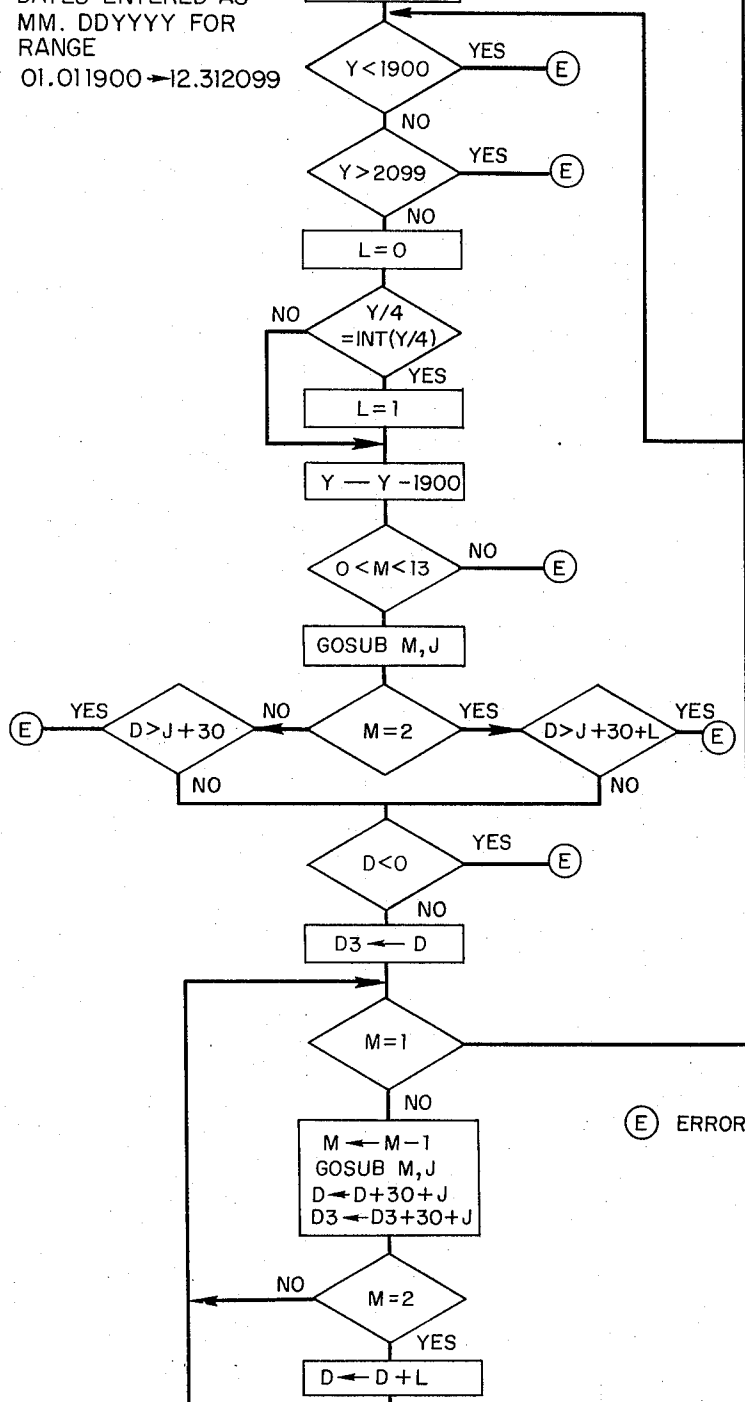
Fig. 34

7.1 DATE1 - DATE2



7.2 DATE ± N DAYS

DATES ENTERED AS
MM. DDYYYY FOR
RANGE
01.011900 → 12.312099



(E) ERROR

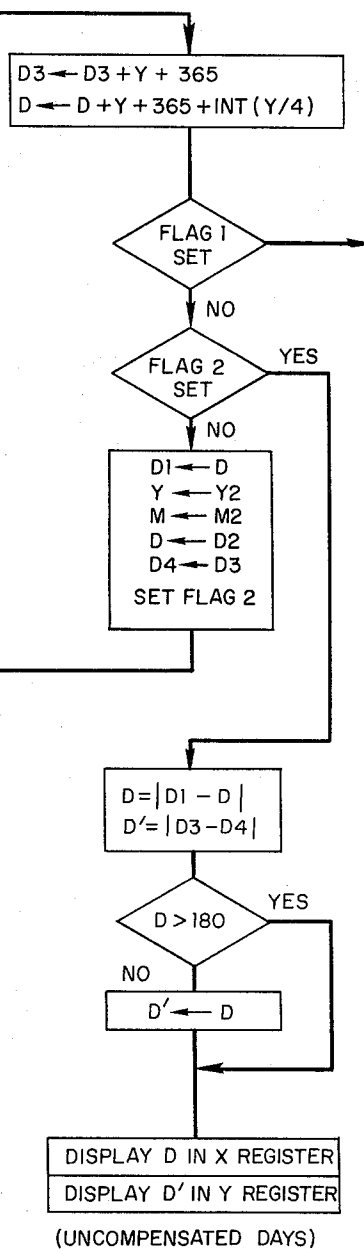


Figure 35A

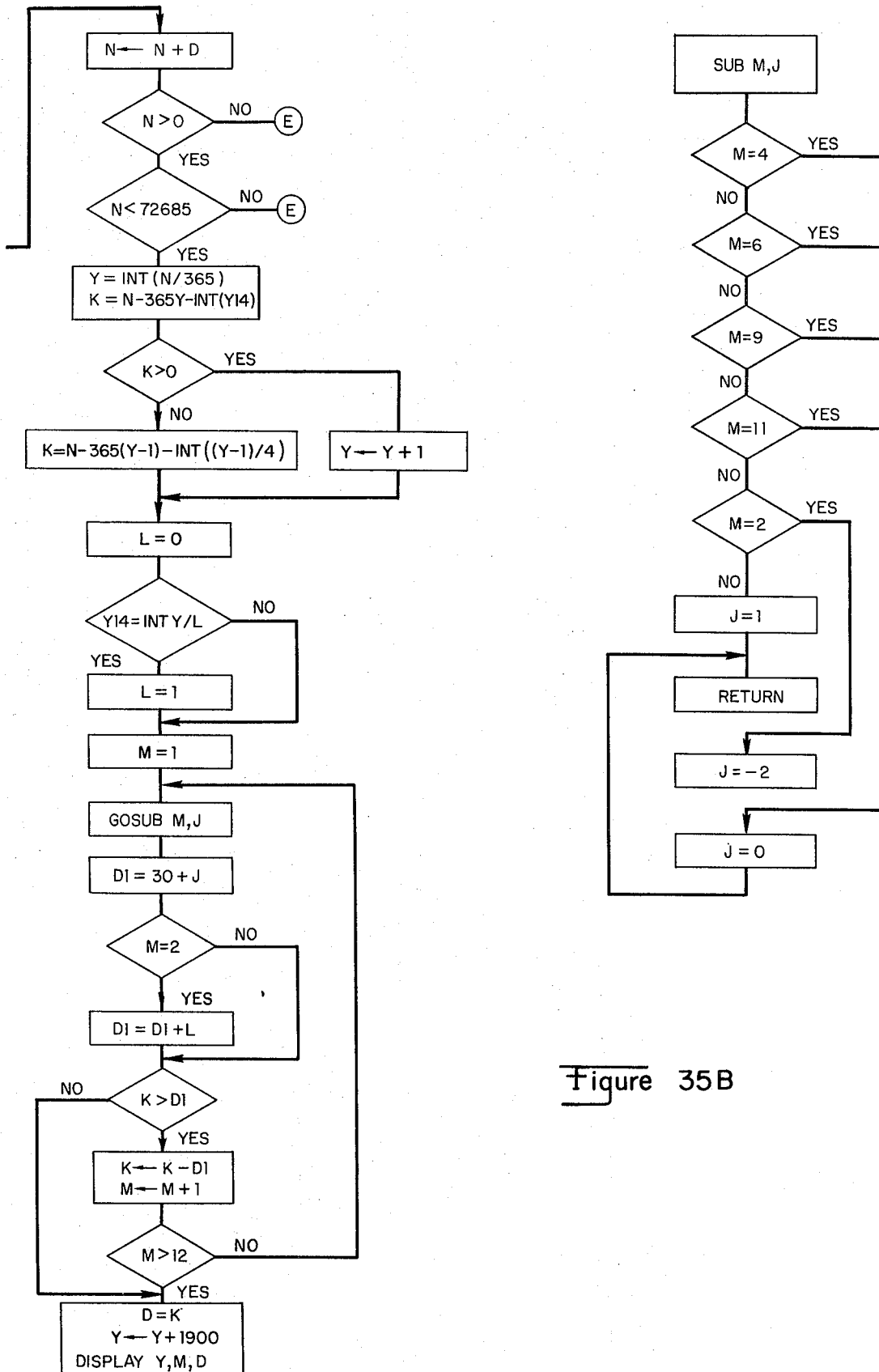


Figure 35B

GENERAL PURPOSE CALCULATOR WITH CAPABILITY FOR PERFORMING INTERDISCIPLINARY BUSINESS CALCULATIONS

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Read-Only Memory Circuit

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Clock Driver

Anode Driver

Cathode Driver

Keyboard

LED Display

Instruction Set

Detailed Listing of Routines and Subroutines of Instructions

Functions

Operating Instructions

BACKGROUND AND SUMMARY OF THE INVENTION

This invention relates generally to calculators and improvements therein and more particularly to non-programmable business calculators.

Conventional business calculators generally have less capability and flexibility than is required to meet the needs of the business user. They are usually designed to solve the most elementary calculation of one business discipline (i.e., banking or real estate, or finance, etc) and lack the capability and flexibility for interdisciplinary business calculations. For example, there are special calculators for financiers to solve bond yield and bond price problems, and calculators for realtors to solve mortgage amortization and depreciation problems. However, a financier who wishes to quickly compare the rate of return between bonds and real estate will either need two expensive calculators or will have to compromise the degree of accuracy of the calculation with gross mathematical approximations performed on a single purpose calculator. This limitation of single purpose calculators can lead to critical errors in decision making. Because conventional single purpose business calculators are designed for special applications by specialists in that area, the keyboards are generally not selfexplanatory and appear as a befuddling collection of buttons and switches with special symbols. This requires a longer user orientation period before productive usage begins.

Due to the high cost and the limited capabilities of the available business calculators, and sometimes, simply because there is no calculator available to perform certain calculations, the majority of the everyday business calculations are still made with the aid of published tables. Published tables are the only convenient means available for solving certain financial problems, such as calculations for the discount amount in discounted notes and the equivalent interest rate between accrued interest notes and discounted notes. The main disadvantage of using tables is the inherent restriction to the discrete values given in the table. The accuracy of the calculation is limited to the accuracy of the tables and the need for interpolation further compro-

mises the calculation. For example, a widely used bond value table has discrete values for bond yield to two decimal places and the interest rate is given in one-eighth of one percent increments. The use of tables with this limited accuracy could lead to errors of several thousands of dollars in a 50 million dollar bond issue.

Another disadvantage of using tables is the requirement that the user must have a working knowledge of both the problem area and the mathematical formulas to set up the problem in a specific manner before the tables are applicable. Even then, it is often necessary to take a reciprocal or multiply by a constant before the answer is usable. This limits the use of the table to only those with a certain level of expertise in the problem area. Thus, one who performs a great variety of business calculations, from asset depreciation to sale forecasting must have: (1) an expensive collection of special purpose calculators; or (2) a library of tables close at hand; or (3) the mathematical and financial expertise to set up and solve the problem correctly.

The principle object of this invention is to provide a general purpose business calculator that has vastly greater capability and flexibility than conventional business calculators and that is small, inexpensive and easier to use than conventional business calculators. This calculator was designed to incorporate into one small calculator the capability of performing the majority of the calculations used in the many disciplines of business and performing these calculations with up to ten-digit accuracy. It replaces the special calculators designed for banking, or accounting, or finance, or real estate and other businesses and eliminates the need for all commonly used financial tables. It also allows the user to make inter-disciplinary analysis, for example, between real estate or bond investment programs, quickly and with one calculator. Furthermore, with the present invention, a sophisticated user can incorporate the mathematical formulas of several business disciplines to solve a complex problem involving several disciplines.

Another object of this invention is to provide a small business calculator which does not require a high level of user expertise or a working knowledge of the problem area and the necessary mathematical formulas before the problem can be set up and solved. Keys relating to a general class of problems are grouped together and designated in accordance with the generally accepted business symbols (e.g., *i* for interest per period, PMT for payment per period, etc.). The key layout and the keying sequence are such that they suggest to the non-expert user the information necessary to solve a given problem. For example, in solving the general class of compound interest and annuity problems with this calculator, the five possible variables, number of time periods, interest rate per period, payment per period, the present value and the future value are all located on the top row. A user can key in any of the three variables in the prescribed left to right sequence and the calculator will solve either of the remaining unknowns as requested. This procedure does not require any previous knowledge of compound interest or annuity mathematics, and any of the five variables can be solved without any intermediate steps. Hence, all that is required of the user is that he be able to define the variables of the problem, and the unique keying se-

quence of the invention will carry out the necessary mathematical manipulation.

A calculator constructed according to the preferred embodiment of this invention is small enough to hold in one hand, capable of displaying data as it is entered and a numerical result as it is calculated, and incorporates many complex functions in order to perform the number and kind of calculations and mathematical operations required for different business disciplines. The limits of miniaturization and sophistication are realized, however, if the keyboard of such a calculator becomes so small and so crowded with keys that the human hand can no longer physically or conveniently manipulate them. One solution to this problem is to reduce the number of functions the calculator can perform. A better solution is to assign more than one function to each key, thus reducing the number of keys necessary to incorporate all the functional capabilities of the calculator.

As more functions are assigned to each key, however, the clarity of labelling a key's various functions becomes important. Moreover, not only must the labelling clearly refer to the particular key, but the functions each key causes the machine to perform should be easily understood and learned by the user from scanning the keyboard labels. After learning the total capability of the particular machine from a reading of the manual, the user should be able to know the relationship between keys, the function(s) each key initiates and, therefore, the operation of the machine, from his knowledge of the keyboard itself.

The limitations of miniature calculators are overcome in the preferred embodiment of the present invention by incorporating easily interpreted, coded legends on the surface of the keyboard immediately above certain keys to which an additional function is assigned. The coding, not only designates the additional function of each such key, but also refers to the prefix key which is depressed to activate the additional function of the key.

Some conventional business calculators have a day between date function, but do not check for improper date entries (e.g., 32nd of June) or compensate for the extra day in a leap year. The present invention automatically checks for improper date entries and compensates for the extra days in leap years in all number-of-days calculations between 1900 A.D. and 2100 A.D. Furthermore, the present invention has the unique capability of determining a future or past date with compensation for leap years given the number of days away.

Conventional business calculators have very complicated procedures for establishing a trend line from a set of periodic data points. In the existing prior art, the user enters the data points and is given the value of the y intercept and the slope of the straight line that best fits the data points (hereinafter referred to as the best fit line). In order to forecast future values, the user had to multiply the slope by the future time period and add the result to the value of the y intercept to get the desired future value.

The present calculator is capable of determining a best fit line from a set of data points and, without any intermediate steps or interpolation, of providing ordinal values on the best fit line at any point on the time axis. This calculator can also extrapolate either single or multiple time periods into the past or the future.

Hence, the user can either request the ordinal value at any time period up to 10 digits long (e.g., -2.5, 0, 7.53452) or utilize the feature which automatically calculates the ordinal value for single time period increments.

Conventional business calculators for bond price and bond yield calculations have a manual switch to initiate different bond price and bond yield algorithms for bonds maturing in less than 181 days (these are considered as notes rather than bonds). The present calculator has an automatic feature to check the maturity period and initiate the proper algorithm. The algorithms traditionally used in conventional business calculators to solve bond price and bond yield are very complex and require extensive hardware capability. This has made these calculators large, complicated, and expensive. In the present calculator, two new algorithms requiring significantly less hardware for bond price and bond yield calculations are used to allow the incorporation of these two bond calculations in a small, general purpose calculator at a fraction of the special bond calculator price.

Conventional business calculators designed to calculate accumulated mortgage interest and remaining principal on a mortgage can give the cumulative totals up to a given time period. However, it is often necessary to determine the accumulated mortgage interest and the cumulative principal paid during a specific time period. This is not possible with prior art calculators without making two separate calculations and then taking the difference. The present invention permits the user to find the amount of mortgage interest paid during any specified time period, as well as the remaining principal in one step. Thus, this calculator can for example automatically calculate either the interest paid during the past year or the interest paid during the sixth through the tenth years.

Conventional business calculators with discount cash flow capability discount the entire stream of expected cash flow and solve for the rate of return of the investment. This provides for a summary analysis of the cash producing life span of an asset, but does not provide any interim information on the repayment of the original investment. The present calculator has a discount cash flow capability that discounts each cash flow and maintains a running total of the outstanding amount of the original investment. Hence, when the outstanding balance becomes zero, or greater, the user will know the number of periods before the investment is recouped.

In the past discounted note calculations were made by employing discount note tables where the discounted interest rate is given at 0.05 percent increments, the discounted amount is given to six decimal places, and the equivalent annual interest rate is given only to four decimal places. Thus, one who wished to find the discounted amount or to convert the discounted interest rate to an equivalent annual interest rate is faced with two limitations: (1) the discrete interest value and, hence, the need for interpolation to obtain the discounted interest rate; and (2) the four decimal place accuracy of the equivalent annual interest rates. Both of these aforementioned limitations can have a sizable effect on large sum calculations.

In some money markets outside the United States, a 365 day year is used for interest bearing transactions. This makes it necessary for an international investor to

make extra calculations to convert a given interest rate from a 360 day basis to a 365 day basis or vice versa. Thus, it would be very helpful for an international investor to have the capability of finding the equivalent interest for either interest bearing period.

The present calculator replaces the discount note tables and performs the calculations related to discounted notes without being limited to discrete discounted interest rates and is accurate to ten digits. It automatically calculates the discounted amount and the equivalent annual interest rate to ten digit accuracy for both the 360 day year and the 365 day year which enables prompt evaluation of interest bearing instruments in different money markets.

Conventional business calculators with arithmetic mean and standard deviation capability are very restricted in their performance. In most cases, to find the standard deviation, the user must first find the variance from the sum of the squares of the input data and then take the square root of the variance to find the standard deviation. Also, once the data is entered and the arithmetic mean calculation is made, it is not possible to add or remove a data point to evaluate the effect on the arithmetic mean and the standard deviation without entering all the data points and performing the calculations over again.

The present calculator computes both the arithmetic mean and the standard deviation automatically from the input data. Furthermore, once the arithmetic mean and the standard deviation are calculated, the present invention permits the user to add or remove input data values from the original set of data and calculate a new arithmetic mean and standard deviation without re-entering all the input values. Hence, this machine is highly interactive and permits the user to measure the influence of hypothetical inputs on the existing data.

The present calculator also has the capability for determining a depreciation schedule based on the sum-of-the-digit depreciation method. Given the life of an asset and the depreciable amount, this calculator computes the depreciation for any requested period and the remaining book value to be depreciated. Furthermore, the user can find the same information for all subsequent periods to set up a depreciation schedule.

In order to implement the extended capability and flexibility provided by this calculator, new algorithms were developed requiring less hardware to solve complex problems. A new algorithm using internal transformations was developed to solve for the interest in the present value of an annuity and in the future value of an annuity. This same algorithm also solves for the effective annual percentage rate from the so-called "add-on" rate. Thus, this one algorithm has made it possible for a user to solve any of these three inherently different interest problems without having to identify the problem. The present calculator automatically identifies the type of interest problem from the prescribed left to right keying sequence of the relevant data, and the input data is then converted to a form acceptable to the new algorithm.

Another new algorithm was also developed to reduce the complexity of the bond price and bond yield problem (and the solution thereof) so as to make this problem solvable using only five registers. This new algorithm uses an explicit term that eliminates the need for a series of summations, which would otherwise require substantially more hardware.

The algorithms for performing the functions of this calculator are stored in a read-only memory circuit including seven serial-address in, serial-instruction out read-only memories regulated by a control and timing circuit. This control and timing circuit includes a microprogrammed controller, which receives status conditions from throughout the calculator and sequentially outputs signals to control the flow of data. The control and timing circuit also scans the keyboard to obtain a six-bit read-only memory address, which is generated at the keyboard each time a key is actuated as required to initiate one or more algorithms for performing the functions associated with the actuated key.

Information from the addressed read-only memory is transmitted serially to an arithmetic and register circuit where a serial, binary-coded decimal (BCD) adder/subtractor performs the basic computations. The results of the computations are transmitted to registers in this circuit where they are either stored temporarily or outputted via a seven-segment, 15 -digit, LED display.

DESCRIPTION OF THE DRAWINGS

FIG. 1 is a top view of a business calculator according to the preferred embodiment of the invention.

FIG. 2 is a block diagram of the calculator of FIG. 1.

FIG. 3 is a waveform diagram illustrating the timing sequence of the interconnecting busses and lines of FIG. 2.

FIG. 4 is a block diagram of the control and timing circuit of FIG. 2.

FIG. 5 is a more detailed block diagram of the keyboard scanning circuitry of FIG. 4.

FIG. 6 is a block diagram of one of ROM's 0-6 of FIG. 2.

FIG. 7 is a waveform diagram illustrating a typical address signal and a typical instruction signal.

FIG. 8 is a timing diagram illustrating the important timing points for a typical addressing sequence.

FIG. 9 is a waveform diagram illustrating the word select signals generated in the control and timing circuit of FIGS. 2 and 4 and in ROM's 0-6 of FIGS. 2 and 6.

FIG. 10 is a block diagram of the arithmetic and register circuit of FIG. 2.

FIG. 11 is a path diagram of the actual data paths for the registers A-F and M of FIG. 10.

FIG. 12 is a waveform diagram illustrating the output signals for the display decoder outputs A-E of FIGS. 2, 10, and 11.

FIG. 13 is a waveform diagram illustrating the actual signals on the display decoder outputs A-E of FIGS. 2, 10, and 11 when the digit 9 is decoded.

FIG. 14 is a waveform diagram illustrating the timing of the START signal generated by the display decoder of FIG. 10.

FIG. 15 is a schematic diagram of the clock driver of FIG. 2.

FIG. 16 is a waveform diagram illustrating the timing relationship between the input and output signals of the clock driver of FIG. 15.

FIG. 17 is a logic diagram of the anode driver of FIG. 2.

FIG. 18 is a waveform diagram illustrating the timing relationship between the input, output, and other signals of the anode driver of FIG. 17.

FIG. 19 is a schematic diagram of the basic inductive drive circuit for one of the LED's employed in the LED display of FIG. 2.

FIG. 20 is a waveform diagram illustrating the timing relationship between the decimal point drive signals for the LED display of FIG. 2.

FIG. 21 is a schematic diagram of the inductive drive circuit for one digit in the LED display of FIG. 2.

FIG. 22 is a logic diagram of the cathode driver of FIG. 2.

FIG. 23 is a top view of a metal strip employed in the keyboard input unit of FIGS. 1 and 2.

FIG. 24 is a side view of the metal strip of FIG. 23.

FIG. 25 is a force-deflection curve for a typical key in the keyboard input unit of FIGS. 1 and 2.

FIG. 26 is a schematic diagram of the LED display of FIGS. 1 and 2 and the inductive drive circuits therefor.

FIG. 27 is a schematic diagram of one segment of the LED display of FIG. 26.

FIG. 28 is an equivalent piecewise-linear model for the circuitry of FIG. 27.

FIG. 29 is a waveform diagram illustrating the inductor current and LED anode voltages in the circuitry of FIG. 27.

FIG. 30 is a path diagram illustrating the possible transfer paths between ROM's 0-6 of FIG. 2.

FIG. 31 is a flow diagram of the display wait loop employed in the calculator of FIGS. 1 and 2.

FIG. 32 is a flow diagram of an interest algorithm employed in the calculator of FIGS. 1 and 2.

FIG. 33 is a flow diagram of a price of a bond algorithm employed in the calculator of FIGS. 1 and 2.

FIG. 34 is a flow diagram of a yield to maturity of a bond algorithm employed in the calculator of FIGS. 1 and 2.

FIGS. 35A and B comprise a flow diagram of a date algorithm employed in the calculator of FIGS. 1 and 2.

DESCRIPTION OF THE PREFERRED EMBODIMENT

SYSTEM ARCHITECTURE

Referring to FIGS. 1 and 2, there is shown a pocket-size electronic calculator 10 including a keyboard input unit 12 for entering data and instructions into the calculator and a seven-segment LED output display unit 14 for displaying each data entry and the results of calculations performed by the calculator. As shown in FIG. 2, calculator 10 also includes an MOS control and timing circuit 16, an MOS read-only memory circuit 18 (including ROM's 0-6), an MOS arithmetic and register circuit 20, a bipolar clock driver 22, and a solid state power supply unit 24.

The three MOS circuits are two-phase dynamic MOS/LSI circuits with low thresholds allowing compatibility with TTL bipolar circuits and allowing extremely low-power operation (less than 100 milliwatts for all three circuits). They are organized to process fourteen-digit BCD words in a digit-serial, bit-serial manner. The maximum bit rate or clock frequency is 200 kilohertz, which gives a word time of 280 microseconds (permitting a floating point addition to be completed in 60 milliseconds).

Control and timing circuit 16, read-only memory circuit 18, and arithmetic and register circuit 20 are tied together by a synchronization (SYNC) buss 26, an instruction (I_a) buss 28, a word select (WS) buss 30, an instruction address (I_a) line 32, and a carry line 34. All

operations occur on a 56 bit (b_0 - b_{55}) word cycle (14 four-bit BCD digits). The timing sequence for the interconnecting busses and lines 26-34 is shown in FIG. 3.

The SYNC buss 26 carries synchronization signals from control and timing circuit 16 to ROM's 0-6 in read-only memory circuit 18 and to arithmetic and register circuit 20 to synchronize the calculator system. It provides one output each word time. This output also functions as a ten-bit wide window (b_{45} - b_{54}) during which I_a buss 28 is active.

The I_a buss 28 carries ten-bit instructions from the active ROM in the read-only memory circuit 18 to the other ROM's, control and timing circuit 16, and arithmetic and register circuit 20, each of which decodes the instructions locally and responds to or acts upon them if they pertain thereto and ignores them if they do not. For instance, the ADD instruction affects arithmetic and register circuit 20 but is ignored by control and timing circuit 16. Similarly, the SET STATUS BIT 5 instruction sets status flip-flop 5 in control and timing circuit 16 but is ignored by arithmetic and register circuit 20.

The actual implementation of an instruction is delayed one word time from its receipt. For example, an instruction may require the addition of digit 2 in two of the registers in arithmetic and register circuit 20. The ADD instruction would be received by arithmetic and register circuit 20 during bit times b_{45} - b_{54} of word time N and the addition would actually occur during bit times b_8 - b_{11} of word time N + 1. Thus, while one instruction is being executed the next instruction is being fetched.

The WS buss 30 carries an enable signal from control and timing circuit 16 or one of the ROM's in read-only memory circuit 18 to arithmetic and register circuit 20 to enable the instruction being executed thereby. Thus, in the example of the previous paragraph, addition occurs only during digit 2 since the adder in the arithmetic and register circuit 20 is enabled by WS buss 30 only during this portion of the word. When WS buss 30 is low, the contents of the registers in arithmetic and register circuit 20 are recirculated unchanged. Three examples of WS timing signals are shown in FIG. 3. In the first example, digit 2 is selected out of the entire word. In the second example, the last eleven digits are selected. This corresponds to the mantissa portion of a floating point word format. In the third example, the entire word is selected. Use of the word select feature allows selective addition, transfer, shifting or comparison of portions of the registers within arithmetic and register circuit 20 with only one basic ADD, TRANSFER, SHIFT, or COMPARE instruction. Some customization in the ROM word select fields is available via masking options.

The I_a line 32 serially carries the addresses of the instructions to be read from ROM's 0-6. These addresses originate from control and timing circuit 16, which contains an instruction address register that is incremented each word time unless a JUMP SUBROUTINE or a BRANCH instruction is being executed. Each address is transferred to ROM's 0-6 during bit times b_{19} - b_{26} and is stored in an address register of each ROM. However, only one ROM is active at a time, and only the active ROM responds to an address by outputting an instruction on the I_a line 28. Control is transferred between ROM's by a ROM SELECT instruction. This technique allows a single eight-bit address, plus

eight special instructions, to address up to eight ROM's of 256 words each.

The carry line 34 transmits the status of the carry output of the adder in arithmetic and register circuit 20 to control and timing circuit 16. The control and timing circuit uses this information to make conditional branches, dependent upon the numerical value of the contents of the registers in arithmetic and register circuit 20.

Control and timing circuit 16 is organized to scan a five-by-eight matrix of switches in search of an interconnection that designates actuation of a key. Any type of metal-to-metal contact may be used as a key. Bounce problems are overcome by programmed lock-outs in the key entry routine. Each key has an associated six-bit code.

A power on circuit 36 in power supply unit 24 supplies a signal forcing the calculator to start up in a known condition when power is supplied thereto. Power is supplied to the calculator when the on-off switch of keyboard input unit 12 (see FIG. 1) is moved to the on position.

The primary outputs of the calculator are five output lines 38 connected between a display decoder of arithmetic and register circuit 20 and an anode driver of output display unit 14. Data for a seven-segment display plus a decimal point is time-multiplexed onto these five output lines. A start line 40 connected between the display decoder of arithmetic and register circuit 20 and a cathode driver of output display unit 14 indicates when the digit 0 occurs.

CONTROL AND TIMING CIRCUIT

Referring now to FIG. 4, control and timing circuit 16 contains the master system counter 42, scans the keyboard 12, retains status information about the system or the condition of an algorithm, and generates the next ROM address. It also originates the subclass of Word Select signals which involve the pointer 44, a four-bit counter that points to one of the register digit positions.

The control unit of control and timing circuit 16 is a microprogrammed controller 46 comprising a 58 word (25 bits per word) control ROM, which receives qualifier or status conditions from throughout the calculator and sequentially outputs signals to control the flow of data. Each bit in this control ROM either corresponds to a single control line or is part of a group of N bits encoded into 2^N mutually exclusive control lines and decoded external to the control ROM. At each phase 2 clock, a word is read from the control ROM as determined by its present address. Part of the output is fed back to become the next address.

Several types of qualifiers are checked. Since most commands are issued only at certain bit times during the word cycle, timing qualifiers are necessary. This means the control ROM may sit in a wait loop until the appropriate timing qualifier comes true, then move to the next address to issue a command. Other qualifiers are the state of the pointer register, the PWO (power on) line, the CARRY flip flop, and the state of each of the 12 status bits.

Since the calculator is a serial system based on a 56 bit word, a six-bit system counter 42 is employed for counting to 56. Several decoders from system counter 42 are necessary. The SYNC signal is generated during bit times b_{45} - b_{54} and transmitted to all circuits in the

system (see FIG. 3). Other timing qualifiers are sent to the microprogrammed control ROM 46 as mentioned in the previous paragraph.

System counter 42 is also employed as a keyboard scanner as shown in FIG. 5. The most significant three bits of system counter 42 go to a one-of-eight decoder 48, which sequentially selects one of the keyboard row lines 50. The least significant three bits of the system counter count modulo seven and go to a one-of-eight multiplexor 52, which sequentially selects one of the keyboard column lines 54 (during sixteen clock times no key is scanned). The multiplexor output is called the key down signal. If a contact is made at any intersection point in the five-by-eight matrix (by depressing a key), the key down signal will become high for one state of system counter 42 (i.e., when the appropriate row and column lines are selected). The key down signal will cause that state of the system counter to be saved in key code buffer 56. This six-bit code is then transferred to the ROM address register 58 and becomes a starting address for the program which services the key that was down (two leading 0 bits are added by hardware so an eight-bit address exists). Thus, during each state of system counter 42, the decoder-multiplexor combination 48 and 52 is looking to see if a specific key is down. If it is, the state of the system counter becomes a starting address for execution of that key function (note that 16 of the 56 states are not used for key codes). By sharing the function of the system counter and using a keyboard scanning technique directly interfaced to the MOS circuitry, circuit complexity is reduced significantly.

A 28 bit shift register which circulates twice each 56 bit word time, is employed in control and timing circuit 16. These 28 bits are divided into three functional groups: the main ROM address register 58 (eight bits), the subroutine return address register 60 (eight bits), and the status register 62 (twelve bits).

The main ROM's 0-6 each contain 256 (ten bit) words, thereby requiring an eight-bit address. This address circulates through a serial adder/subtractor 64 and is incremented during bit times b_{47} - b_{54} (except in the case of branch and jump-subroutine instructions for which the eight bit address field of the ten-bit instruction is substituted for the current address). The next address is transmitted over the I_a line 32 to each of the main ROM's 0-6 during bit times b_{19} - b_{26} .

The Status register 62 contains twelve bits or flags which are used to keep track of the state of the calculator. Such information as whether the decimal point has been hit, the minus sign set, etc. must be retained in the status bits. In each case the calculator remembers past events by setting an appropriate status bit and asking later if it is set. A yes answer to a status interrogation will set the carry flip-flop 66 as indicated by control signal IST in FIG. 4. Any status bit can be set, reset, or interrogated while circulating through the adder 64 in response to the appropriate instruction.

The instruction set allows one level of subroutine call. The return address is stored in the eight-bit return address register 60. Execution of a JUMP subroutine stores the incremented present address into return address register 60. Execution of the RETURN instruction retrieves this address for transmission over I_a line 32. Gating is employed to interrupt the 28 bits circulating in the shift register 58-62 for insertion of addresses

at the proper time as indicated by the JSB control signal in FIG. 4.

An important feature of the calculator system is the capability to select and operate upon a single digit or a group of digits (such as the exponent field) from the fourteen digit registers. This feature is implemented through the use of a four-bit pointer 44 which points at the digit of interest. Instructions are available to set, increment, decrement, and interrogate pointer 44. The pointer is incremented or decremented by the same serial adder/subtractor 64 used for addresses. A yes answer to the instruction "is pointer N" will set the carry flip-flop 66 via control signal IPT in FIG. 4.

The word select feature was discussed above in connection with FIGS. 2 and 3. Some of the word select signals are generated in control and timing circuit 16, namely those dependent on pointer 44, and the remainder in the main ROM's 0-6. The pointer word select options are (1) pointer position only and (2) pointer position and all less significant digits. For instance, assume the mantissa signs of the numbers in the A and C registers of arithmetic and register circuit 20 were to be exchanged. The pointer would be set to position 13 (last position) and the A EXCHANGE C instruction with a "pointer position" word select field would be given. If all of the word except the mantissa signs were to be exchanged, the A EXCHANGE C instruction would be given with the pointer set at 12 and the word select field set to pointer and less significant digits. The control and timing circuit word-select output 30 is tied with the ROM word-select output 30 and transmitted to arithmetic and register circuit 20.

Any carry signal out of the adder in arithmetic and register circuit 20, with word select also high, will set carry flip-flop 66. This flip-flop is interrogated during the BRANCH instruction to determine if the existing address should be incremented (yes carry) or replaced by the branch address (no carry). The branch address is retained in an eight-bit address buffer 68 and gated to I_a line 32 by the BRH control signal.

The power-on signal is used to synchronize and pre-set the starting conditions of the calculator. It has two functions, one of which is to get the address of control ROM 46 set to a proper starting state and the other of which is to get the system counter 42 in control and timing circuit 16 synchronized with the counter in each main ROM 0-6. As the system power comes on, the PWO signal is held at logic 1 (0 volts in this system) for at least 20 milliseconds. This allows system counter 42 to make at least one pass through bit times b_{45} - b_{54} when SYNC is high thereby setting main ROM 0 active and the rest of the ROM's inactive. When PWO goes to logic 0 (+6 volts), the address of control ROM 46 is set to 000000 where proper operation can begin.

READ-ONLY MEMORY CIRCUIT

The ROM's 0-6 in read-only memory circuit 18 store the programs for executing the functions required of the system. Since each ROM contains 256 ten-bit words, there are 1,536 words or 15,360 bits of ROM. A block diagram of each of the ROM's 0-6 is shown in FIG. 6.

The basic operation of each ROM is a serial-address in, a serial-instruction out. During every 56 bit word time the address comes in, least significant bit first from bit b_{19} through bit b_{28} . Every ROM 0-6 in the system receives this same eight-bit address and from bit time b_{45}

through b_{54} tries to output onto I_s line 28. However, a ROM enable (ROE) flip-flop 70 in each ROM insures that no more than one ROM actually sends an instruction on I_s line 28 at the same time.

All output signals are inverted so that the steady-state power dissipation is reduced. The calculator circuits are P-channel MOS. Thus, the active signals that turn on a gate are the more negative. This is referred to as negative logic, since the more negative logic level is the logic 1. As mentioned above, logic 0 is +6 volts and logic 1 is 0 volts. The signals on I_a and I_s are normally at logic 0. However, when the output buffer circuits are left at logic 0 they consume more power. A decision was therefore made to invert the signals on the I_a and I_s outputs and re-invert the signals at all inputs. Thus, signals appear at the I_a and I_s outputs as positive logic. The oscilloscope pattern that would be seen for instruction 1101 110 011 from state 11 010 101 is shown in FIG. 8.

The serial nature of the calculator circuits requires careful synchronization. This synchronization is provided by the SYNC pulse, generated in control and timing circuit 16 and lasting for bit times b_{45} - b_{54} . Each ROM has its own 56 state counter 72, synchronized to the system counter 42 in control and timing circuit 16. Decoded signals from this state counter 72 open the input to the address register 74 at bit time b_{19} , clock I_s out at bit time b_{45} and provide other timing control signals.

As the system power comes on, the PWO signal is held at 0 volts (logic 1) for at least 20 milliseconds. The PWO signal is wired (via a masking option) to set ROM Output Enable (ROE) flip-flop 70 on main ROM 0 and reset it on all other ROM's. Thus when operation begins, ROM 0 will be the only active ROM. In addition, control and timing circuit 16 inhibits the address output during start-up so that the first ROM address will be zero. The first instruction must be a JUMP SUBROUTINE to get the address register 58 in control and timing circuit 16 loaded properly.

FIG. 7 shows the important timing points for a typical addressing sequence. During bit times b_{19} - b_{28} the address is received serially from control and timing circuit 16 and loaded into address register 74 via I_a line 32. This address is decoded and at bit time b_{44} the selected instruction is gated in parallel into the I_s register 76. During bit times b_{45} - b_{54} the instruction is read serially onto I_s buss 28 from the active ROM (i.e., the ROM with the ROM enable flip-flop set).

Control is transferred between ROM's by a ROM SELECT instruction. Effectively this instruction will turn off ROE flip-flop 70 on the active ROM and turn on ROE flip-flop 70 on the selected ROM. Implementation is dependent upon the ROE flip-flop being a master-slave flip-flop. In the active ROM, the ROM SELECT instruction is decoded by a ROM select decoder 78 at bit time 44, and the master portion of ROE flip-flop 70 is set. The slave portion of ROE flip-flop 70 is not set until the end of the bit time (b_{55}). In the inactive ROM's the instruction is read serially into the I_s register 76 during bit times b_{45} - b_{54} and then decoded, and the ROE flipflop 70 is set at bit time b_{55} in the selected ROM. A masking option on the decoding from the least significant three bits of the I_s register 76 allows each ROM to respond only to its own code.

The six secondary word-select signals are generated in the main ROM's 0-6. Only the two word-select sig-

nals dependent upon the POINTER come from control and timing circuit 16. The word select of the instruction is retained in the word select register 80 (also a master-slave). If the first two bits are 01, the instruction is of the arithmetic type for which the ROM must generate a word select gating signal. At bit time b_{55} the next three bits are gated to the slave and retained for the next word time to be decoded into one of six signals. The synchronization counter 72 provides timing information to the word select decoder 82. The output WS signal is gated by ROE flip-flop 70 so only the active ROM can output on WS line 30, which is OR-tied with all other ROM's and also control and timing circuit 16. As discussed above, the WS signal goes to arithmetic and register circuit 20 to control the portion of a word time an instruction is active.

The six ROM generated word select signals used in the calculator are shown in FIG. 9. ROM's 0-6 output a one bit-time pulse on I_s buss 28 at bit time b_{11} to denote the exponent minus sign time. This pulse is used in the display decoder of arithmetic and register circuit 20 to convert a 9 into a displayed minus sign. The time location of this pulse is a mask option on the ROM.

ARITHMETIC AND REGISTER CIRCUIT

Arithmetic and register circuit 20 shown in FIG. 10 provides the arithmetic and data storage function for the calculator. It is controlled by the WS, I_s , and SYNC lines 30, 28, and 26, respectively; receives instructions from ROM's 0-6 over the I_s line 28; sends information back to control and timing circuit 16 via the CARRY line 34; partially decodes the display information before transmitting it via output lines 38 to the anode driver of output display unit 14; and provides a START pulse to the cathode driver of output display unit 14 to synchronize the display.

Arithmetic and register circuit 16 contains seven, fourteen-digit (56 bit) dynamic registers A-F and M and a serial BCD adder/subtractor 84. Actual data paths, not shown in FIG. 10 due to their complexity, are discussed below and shown in FIG. 11. The power and flexibility of an instruction set is determined to a great extent by the variety of data paths available. One of the advantages of a serial structure is that additional data paths are not very costly (only one additional gate per path). The structure of arithmetic and register circuit 20 is optimized for the type of algorithms required by the calculator.

The seven registers A-F and M can be divided into three groups: the working registers A, B, and C with C also being the bottom register of a four-register stack; the next three registers D, E, and F in the stack; and a separate storage register M communicating with the other registers through register C only. In FIG. 11, which shows the data paths connecting all the registers A-F and M, each circle represents the fifty-six bit register designed by the letter in the circle. In the idle state (when no instruction is being executed in arithmetic and register circuit 20) each register continually circulates since with dynamic MOS registers information is represented by a charge on a parasitic capacitance and must be continually refreshed or lost. This is represented by the loop re-entering each register.

Registers A, B, and C can all be interchanged. Either register A or C is connected to one adder input, and either register B or C to the other. The adder output can be directed to either register A or C. Certain instruc-

tions can generate a carry via carry flip-flop 85 which is transmitted to control and timing circuit 16 to determine conditional branching. Register C always holds a normalized version of the displayed data.

In the stack formed by registers C, D, E, and F a ROLL DOWN instruction is executed by the following transfers: $F \rightarrow E \rightarrow D \rightarrow C \rightarrow F$. A STACK UP instruction is executed by the following transfers: $C \rightarrow D \rightarrow E \rightarrow F$. Thus, it is possible to transfer a register and also let it recirculate so that in the last example the contents of C are not lost. The structure and operation of a stack such as this are further described in copending U.S. Pat. application Ser. No. 257,606 entitled IMPROVED PORTABLE ELECTRONIC CALCULATOR, filed on May 30, 1972, by David S. Cochran et al, and issued as U.S. Pat. No. 3,781,820 on Dec. 25, 1973.

In serial decimal adder/subtractor 84 a correction (addition of 6) to a BCD sum must be made if the sum exceeds nine (a similar correction for subtraction is necessary). It is not known if a correction is needed until the first three bits of the sum have been generated. This is accomplished by adding a four-bit holding register 86 ($A_{60}-A_{57}$) and inserting the corrected sum into a portion 88 ($A_{56}-A_{53}$) of register A if a carry is generated. This holding register 86 is also required for the SHIFT A LEFT instruction. One of the characteristics of a decimal adder is that non-BCD codes (i.e. 1101) are not allowed. They will be modified if circulated through the adder. The adder logic is minimized to save circuit area. If four bit codes other than 0000-1001 are processed, they will be modified. This is no constraint for applications involving only numeric data (however, if ASC11 codes, for instance, are operated upon, incorrect results will be obtained).

Arithmetic and register circuit 20 receives the instruction during bit times $b_{45}-b_{54}$. Of the ten types of instructions hereinafter described, arithmetic and register circuit 20 must respond to only two types (namely, ARITHMETIC & REGISTER instructions and DATA ENTRY/DISPLAY instructions). ARITHMETIC & REGISTER instructions are coded by a 10 in the least significant two bits of I_s register 90. When this combination is detected, the most significant five bits are saved in I_s register 90 and decoded by instruction decoder 92 into one of 32 instructions.

The ARITHMETIC & REGISTER instructions are active or operative only when the Word Select signal (WS) generated in one of the ROM's 0-6 or in control and timing circuit 16 is at logic one. For instance, suppose the instruction " $A+C \rightarrow C$, mantissa with sign only" is called. Arithmetic and register circuit 20 decodes only $A+C \rightarrow C$. It sets up registers A and C at the inputs to adder 84 and, when WS is high, directs the adder output to register C. Actual addition takes place only during bit times b_{12} to b_{55} (digits 3-13) since for the first three digit times the exponent and exponent sign are circulating and are directed unchanged back to their original registers. Thus, the word select signal is an "instruction enable" in arithmetic and register circuit 20 (when it is at logic 1, instruction execution takes place, and when it is at logic 0, recirculation of all registers continues).

The DATA ENTRY/DISPLAY instructions, except for digit entry, affect an entire register (the word select signal generated in the active ROM is at logic 1 for the entire word cycle). Some of these instructions are: up

stack, down stack, memory exchange $M \leftrightarrow C$, and display on or toggle. A detailed description of their execution is given hereinafter.

For increased power savings display decoder 94 is partitioned to partially decode the BCD data into seven segments and a decimal point in arithmetic and register circuit 20 by using only five output lines (A-E) 38 with time as the other parameter. Information for seven segments (a-g) and a decimal point (dp) are time shared on the five output lines A-E. The output wave forms for output lines A-E are shown in FIG. 12. For example, output line D carries the segment e information during T_1 (the first bit time of each digit time) and the segment D information during T_2 (the second bit time of each digit time); and output E carries the segment G information during T_1 , the segment F information during T_2 , and the decimal point (dp) during T_4 . The actual signals which would appear if a digit 9 were decoded are shown in FIG. 13. The decoding is completed in the anode driver of output display unit 14 as explained hereinafter.

The registers in arithmetic and register circuit 20 hold 14 digits comprising ten mantissa digits, the mantissa sign, two exponent digits, and the exponent sign. Although the decimal point is not allocated a register position, it is given a full digit position in the output display. This apparent inconsistency is achieved by using both the A and B registers to hold display information. The A register is set up to hold the displayed number with the digits in the proper order. The B register is used as a masking register with digits 9 inserted for each display position that is to be blanked and a digit 2 at the decimal point location. When the anode driver of output display unit 14 detects a decimal point code during T_4 , it provides a signal to the cathode driver of the output display unit to move to the next digit position. One digit and the decimal point share one of the 14 digit times. The digit 9 mask in register B allows both trailing and leading zeros to be blanked (i.e., by programming 9's into the B register). Use of all three working registers for display (i.e., the C register to retain the number in normalized form, the A register to hold the number in the displayed form, and the B register as a mask) allows the calculator to have both a floating point and a scientific display format at the expense of only a few more ROM states.

The display blanking is handled as follows. At time T_4 the BCD digit is gated from register A into display buffer 96. If this digit is to be blanked, register B will contain a 9 (1001) so that at T_4 the end bit (B_{01}) of the B register will be a one (an eight would therefore also work). The input to display buffer 96 is OR-ED with B_{01} and will be set to 1111 if the digit is to be blanked. The decimal point is handled in a similar way. A 2 (0010) is placed in register B at the decimal point location. At time T_2 the decimal point buffer flip-flop is set by B_{01} . Any digit with a one in the second position will set the decimal point (i.e., 2, 3, 6, or 7).

Display decoder 94 also applies a START signal to line 40. This signal is a word synchronization pulse, which resets the digit scanner in the cathode driver of output display unit 14 to assure that the cathode driver will select digit 1 when the digit 1 information is on outputs A, B, C, D, and E. The timing for this signal is shown in FIG. 14.

One other special decoding feature is required. A minus sign is represented in tens complement notation

or sign and magnitude notation by the digit 9 in the sign location. However, the display must show only a minus sign (i.e., segment g). The digit 9 in register A in digit position 2 (exponent sign) or position 13 (mantissa sign) must be displayed as minus. The decoding circuitry uses the pulse on I_s buss 28 at bit time b_{11} (see FIG. 3) to know that the digit 9 in digit position 2 of register A should be a minus and uses the SYNC pulse to know that the digit 9 in digit position 13 of register A should also be a minus. The pulse on I_s buss 28 at bit time b_{11} can be set by a mask option, which allows the minus sign of the exponent to appear in other locations for other uses of the calculator circuits.

CLOCK DRIVER

The bipolar clock driver 22, one phase of which is shown in FIG. 15, requires less than 25 milliwatts and can drive loads up to 300 picroforads with a voltage swing of +7 to -14 volts. An ENABLE input 98 allows both outputs Q_1 and Q_2 to be held to V_{CC} , the MOS Logic 0 state. This is an effective means of strobing the clock. During dc operation, the transistor pair Q_1 - Q_2 allows only one of the output transistor pairs Q_5 - Q_6 or Q_7 - Q_8 to conduct. Diode D_3 prohibits conduction from transistor Q_6 to transistor Q_8 during transient operation. Thus, the only possible transient short circuit current must flow from transistor Q_5 to transistor Q_7 . However, the limited current handling capability of Q_5 (a lateral PNP) limits this current to less than 5 milliamps peak. The input signals for clock driver 22 are generated on the anode driver of output display unit 14, and the outputs of the clock driver go to each of the MOS circuits in the system. The timing relationships are shown in FIG. 16.

ANODE DRIVER

As discussed above, the display information is partially decoded in arithmetic and register circuit 20 and completely decoded into seven segment plus decimal point signals in the bipolar anode driver of output display unit 14. The anode driver also includes the basic clock generator for the system and a circuit for detecting low battery voltage to turn on all the decimal points. Such a circuit is shown and described in co-pending U.S. Pat. application Ser. No. 206,407 entitled LOW BATTERY VOLTAGE INDICATOR FOR A PORTABLE DIGITAL ELECTRONIC INSTRUMENT, filed on Dec. 9, 1971, by Thomas M. Whitney and now abandoned. A logic diagram of the anode driver is shown in FIG. 17.

The clock generator uses an external LC series circuit to set the oscillator frequency. The advantages of an LC series circuit to set the frequency are: (1) the components can be specified to up to 2 percent tolerance; and (2) a crystal can be connected to the same external pin to set the frequency to 0.001 percent for timing applications.

The square-wave oscillator frequency (all times in this section will be referred to an 800 KHz oscillator frequency, which translates to a 200 KHz clock for the calculator, the actual frequency being somewhat less) is divided by flip-flop BL to 400 KHz. Flip-flops B1 and B2 are clocked off alternate phases of flip-flop B1 to provide two 200 KHz square waves as shown in FIG. 18. Flip-flop B3 is clocked from flip-flop B2 and in turn clocks flip-flop B4 to provide further count-down of the basic clock frequency. The two-phase clock signals

Q_1 and Q_2 are generated from flip-flops BL and B1 and the 800 KHz oscillator 100. They are on for 625 nsec and separated by 625 μ sec as shown in FIG. 18. One other periodic signal is derived in the anode driver. Once each digit time a signal (counter-clock) is sent to the cathode driver of output display unit 14 (the trailing edge of this signal will step the display to the next digit).

The display consists of fifteen characters while the basic calculator word cycle consists of fourteen digits. The extra character is the decimal point. As explained above, a BCD two is placed in register B at the digit position of the decimal point. The display decoder 94 in arithmetic and register circuit 20 indicates this by a signal on outputs B and E during bit time T_4 (see FIG. 12). When this condition is decoded by the anode driver, the decimal point is excited and an extra counter clock signal is given to step the display to the next position (see FIGS. 18, 19, and 20). Therefore all remaining digits in register A are displaced one digit in the display.

FIGS. 19 and 20 show the simplified circuit and the timing relationship for the decimal point. The timing is critical since all the inductor current in segment b (the last to be excited) must be decayed before the counter clock signal is given to step to the next digit or the remaining current would be discharged through the wrong digit and a faint lighting of segment b on the same digit with the decimal point would occur. The decimal point insertion technique is the reason all other seven segments are excited during the first half of the digit time. The decimal point charging time is one-half that of the other segments. The decimal point segment gets the same current in one-half the time and is one-half as bright as the other segments.

An inductive circuit method of driving the light-emitting diodes is employed. Basically the method involves using the time it takes current to build up in an inductor to limit current, rather than using a resistor as is normally done with LED read-outs. This saves power since the only lossy components in the drive system are the parasitic inductor and transistor resistances. The drive circuit for one digit is shown in FIG. 21. Assuming the cathode transistor switch T_c is closed, an anode switch T_a is closed for 2.5 μ sec allowing the current to build up to a value I_p along a nearly triangular waveform (the early part of an exponential buildup). When anode switch T_a is opened, the current is dumped through the LED, decaying in about 5 μ sec. The anodes are strobed according to the sequence in FIG. 18. The primary reason for sequentially exciting the anodes is to reduce the peak cathode transistor current. Since the decay time is approximately twice the buildup time, it works out that the peak cathode current is about 2.5 times the peak current in any segment. The LED's are more efficient when excited at a low duty cycle. This means high currents for short periods (80 ma. anode current, 250 ma. cathode current). FIG. 18 also shows the relationship between the anode strobing sequence and the display output signals (A-E) from arithmetic and register circuit 20.

Since the anode driver operates from the battery voltage directly and drives the decimal point segment, a circuit is provided that senses when the voltage drops below a certain limit and thereupon turns on all decimal points. An external pin is provided to connect a

trimming resistor to set the voltage where the indication is to occur.

CATHODE DRIVER

The cathode driver of output display unit 14 comprises a fifteen position shift register for scanning the fifteen digit display once each word time. This scanning operation moves from digit to digit in response to counter clock signals from the anode driver. Once each word time a START signal arrives from arithmetic and register circuit 20 to restart the procedure. A block diagram is shown in FIG. 22.

KEYBOARD

The calculator employs a reliable, low-profile, low cost keyboard with tactile feedback such as that shown and described in copending U.S. Pat. application Ser. No. 173,754 entitled KEYBOARD HAVING SWITCHES WITH TACTILE FEEDBACK and filed on Aug. 23, 1971, by William W. Misson et al. The keyboard employs metal strips 102 with slots 104 etched or punched out as shown in FIG. 23, leaving an area which can be stretched to form small humps as shown in FIG. 24. The strips are spot welded to a printed circuit board such that orthogonal traces run under each hump. Pressing a key makes electrical contact between one of the horizontal strips and the corresponding vertical trace. The bounce is less than one millisecond (the calculator contains a wait loop to prevent double entries). Extensive life testing of the keyboard indicates more than a million cycles can be expected. Tolerances must be maintained carefully to prevent the possibility that a key is depressed but no contact is made and to insure uniformity.

One of the main advantages of the keyboard is the "overcenter" or "fall away" feel. FIG. 25 shows a force-deflection curve for a typical key. As can be seen a force of about 100 grams must be exceeded before the metal hump "breaks" through. After this critical value the operator cannot prevent contact from being made. Similarly when the key is released, contact is maintained until a critical value when the hump bounces back. Again, past a critical point the operator cannot prevent the key from releasing. This type of action prevents a condition known as "teasing" in which a key is nearly depressed and slight movement causes multiple entries. The point on the force deflection curve at which contact is made or released is most desirably on the negative slope portion. In the calculator it is either there or exactly at the bottom (point A in FIG. 25), but never on the final positive slope portion.

FIG. 1 shows the layout of keyboard input unit 12 which includes a plurality of function and numeric keys. Several of the function keys are capable of performing more than one function when used in conjunction with prefix key 15. For example, function key 17 carries one legend Y^x which refers to its direct function. Immediately above the key location, legend 19 indicates a second function \sqrt{x} . Legend 19 is color-coded to designate not only the second function \sqrt{x} , but also to refer the user to prefix key 15 which initializes that function when depressed prior to depressing key 17. The coloration of the body of prefix key 15 corresponds to the coloration of all legends, such as legend 19, for association with of the present invention initialized by prefix key

15 are YTM, INTR, BOND, Δ%, COMPUTE, DATE, √x, → Σ, CLEAR, and Σ-.

LED DISPLAY

As mentioned above, the inductive drive technique 5 employed for the LED display is inherently efficient because there are no dissipative components other than parasitic resistances and the forward voltage drop across saturated transistor switches. An inductive driver like that used in the calculator is shown and described in copending U.S. Pat. application Ser. No. 202,475 entitled LIGHT EMITTING DIODE DRIVER, filed on Nov. 26, 1971, by Donald K. Miller, and issued as U.S. Pat. No. 3,755,697 on Aug. 28, 1973.

The display circuitry used in the calculator is shown 15 in FIG. 26. It comprises an 8 × 15 array of LED's in which the eight rows are scanned by the anode driver and the fifteen columns by the cathode driver. The timing for this scanning was discussed above. A simplified circuit diagram for one segment is shown in FIG. 27. 20 The equivalent piecewise-linear circuit model is shown in FIG. 28. An analysis of this model shows the inductor current buildup and discharge to be nearly linear for the parameters used in the calculator. The discharge time to charge-time ratio is approximately:

$$t_{discharge}/t_{charge} = (V_s - V_{sat})/(V_d + V_{sat}) = (3.8 - 0.1)/(1.6 + 0.2) = 3.7/1.8 = 2.06$$

FIG. 29 shows the inductor current for a basic calcu-

lator clock frequency of 175 KHz. The average LED current can be calculated from the formula

$$\begin{aligned} \text{Ave } I_{LED} &= \text{pulse current} \times \text{duty cycle} \\ &= \left(\frac{1}{2} \times 80 \text{ ma}\right) \frac{5.88 \text{ sec}}{\frac{1}{175} \text{ KHz} \times 56} \\ &= \frac{(80)(5.88)(.175)}{(2)(56)} = .735 \text{ ma} \end{aligned}$$

The worst case display power (i.e., thirteen 8's and two minus signs) is about 110 milliwatts. FIG. 29 also shows the ringing inherent with inductive drive.

INSTRUCTION SET

Every function performed by the calculator is implemented by a sequence of one or more ten-bit instructions stored in ROM's 0-6 of read-only memory circuit 18. The serial nature of the MOS calculator circuits allows the instruction bits to be decoded from LSB to MSB (right to left) serially. If the first bit is a one, the instruction is either a subroutine jump or a conditional branch as selected by the second bit, with eight bits left for an address. The next largest set of instructions, the arithmetic set, starts with a zero followed by a one (right to left), leaving eight bits for encoded instructions. The ten different types of instructions, employed by the calculator are shown in the table below.

TABLE OF INSTRUCTION TYPES (X=DON'T CARE)

Type	Available Instructions	Name	Fields
1	256 (ADDRESSES)	JUMP SUBROUTINE	SUBROUTINE ADDRESS
	256 (ADDRESSES)	CONDITIONAL BRANCH	BRANCH ADDRESS
2	32×8=256	ARITHMETIC/REGISTER	OPERATION-CODE WORD-SELECT
3	64 (37 used)	STATUS OPERATIONS	N F 1 1 1 1
		SET BIT N INTERROGATE N RESET N CLEAR ALL	F=00 F=01 F=10 F=11 (N=0000)
4	64 (30 used)	POINTER OPERATIONS	P F 1 1 0 0
		SET POINTER TO P INTERROGATE P DECREMENT P INCREMENT P	F=00 F=10 F=01 F=11 P=XXXX
5	64 (20 used)	DATA ENTRY/DISPLAY	N F 1 0 0 0
		LOAD CONSTANT IS→A BCD INPUT TO C REG STACK INSTRUCTIONS AVAILABLE	F=01 F=1X (N=XX01) F=1X (N=XX11) F=10 N=(—0) F=00
6	32 (11 used)	ROM SELECT, MISC.	N F 1 0 0 0 0
		SELECT ROM "N" KEYBOARD ENTRY EXTERNAL ENTRY SUBROUTINE RETURN	F=00 F=10 (N=XX1) F=10 (N=XX0) F=01 (N=XXX)

TABLE OF INSTRUCTION TYPES (X=DON'T CARE) - Continued

Type	Available instructions	Name	Fields
7.....16.....	{ RESERVED FOR PROGRAM STORAGE MOS CIRCUIT }		4 X X X X 1 0 0 0 0 0
8.....8.....			3 X X X 1 0 0 0 0 0 0
9.....7.....	AVAILABLE		X X X 0 0 0 0 0 0 0
10.....1.....	NO OPERATION (NOP)		0 0 0 0 0 0 0 0 0

There are two type 1 instructions, jump subroutine and conditional branch. They are decoded only by control and timing circuit 16. No word select is generated and all registers in arithmetic and register circuit 20 merely recirculate. The object of the jump subroutine instruction is to move to a new address in ROM and to save the existing address (plus one) as a return address. The last instruction in a subroutine must be a RETURN to continue the program where it was left previously.

buffer register 68 onto the I_a buss 32 and loaded into the ROM address register 74 (see FIG. 6). Thus, the instruction is a BRANCH IF NO CARRY. There are three ways the carry flip-flop 66 can be set: (1) by a carry generated in the arithmetic and register circuit 20; (2) by a successful interrogation of the pointer position; and (3) by a successful interrogation of one of the twelve status bits. An example is given in the table below.

Example Conditional Branch Execution				
WORD	ADDRESS RECEIVED AT ROM	INSTRUCTION SENT BY ROM	INSTRUCTION EXECUTED	RESULT
N-1	P	INCREMENT SIGN DIGIT	—	—
N	P+1	CONDITIONAL BRANCH TO ADDRESS Q	INCREMENT SIGN DIGIT	CARRY GENERATED IF "A" REG. NEGATIVE
N+1	P+2	CONTENTS OF P+2	CONDITIONAL BRANCH	SEND P+2
	or Q	or CONTENTS OF Q		or SEND Q

As discussed above, control and timing circuit 16 contains a 20 eight-bit shift register 58-62 which holds the current eight-bit ROM address and also has eight bits of storage for one return address (see FIG. 4). During bit times b₄₇-b₅₄ the current ROM address flows through the adder 64 and is incremented by one. Normally, this address is updated each word time. However, if the first two bits of the instruction, which arrive at bit times b₄₅-b₄₆ are 10, the incremented current address is routed to the return address portion 60 of the 20 eight-bit shift register and the remaining eight bits of the instruction, which are the subroutine address, are inserted into the address portion 58. These data paths with the JSB control line are shown in FIG. 4. In this way the return address has been saved and the jump address is ready to be transmitted to ROM at bit times b₁₉-b₂₆ of the next word time.

The most frequently used instruction is the conditional branch, which based upon data or system status implements the decision-making capability of the calculator. In the calculator system described here this instruction also functions as an unconditional branch.

The format of the branch instruction, as shown in the instruction table above, is two ones followed by an eight-bit branch address. The instruction is received at bit times b₄₅-b₅₄. The last eight bits of the instruction are stored in the address buffer register 68 (see FIG. 4). During the next word time the carry flip-flop 66 is checked at bit time b₁₉. If the carry flip-flop was set during the previous word time, the current ROM address is transmitted to ROM's 0-6. If the carry flip-flop was not set, the branch address is read from the address

A typical test condition is to determine the sign of a number. Support at address P in the program a branch to location Q is desired if the sign of A is positive, while program execution is to continue if the sign if negative. In the example given in the table above, the instruction "increment the A register, word select of sign digit only" is given at location P. During word time N-1 the instruction is received by arithmetic and register circuit 20 and is executed at word time N (the same word time when the CONDITIONAL BRANCH instruction is received by control and timing circuit 16). If the sign of A is negative, there will be a nine in the sign digit. Incrementing this position will generate a carry and set the carry flip-flop 66 in control and timing circuit 16. Since the instruction is a branch if no carry is generated, the program execution will jump to location Q only if the sign is positive (i.e., was a zero), otherwise execution continues at P+2.

Note that during word time N+1 the calculator did nothing more than to select which of two addresses to send next (all registers merely recirculate). To perform a branch actually takes two word cycles to execute, one to ask a question and set the carry flip-flop 66 if the answer is YES, and the other to test if the CARRY flip-flop was set and transmit the proper address. In many cases the asking of the question is an arithmetic operation (i.e., A+B → A) which must be performed anyway. Then the branch takes only one extra instruction.

Contrary to most instruction sets, this set has no unconditional branch instruction. However, since an ordinary "jump" is one of the most used instructions, the conditional branch is also used as an unconditional

branch or jump by insuring that the carry flip-flop 66 is reset when an unconditional branch is desired. This is the reason the sense of the conditional branch is **BRANCH ON NO CARRY**. The carry flip-flop 66 is reset during execution of every instruction except arithmetic (type 2) and interrogation of pointer or status (types 3 and 4). Since only arithmetic and interrogation instructions can set the carry flip-flop 66, the constraint is not severe. The jump subroutine instruction can also be used as an unconditional branch if the previous return address does not have to be saved. In summary, conditional branch can be used as an unconditional branch provided the state of the carry flip-flop 66 is known to be reset (i.e., provided the conditional branch does not follow an arithmetic or an interrogation of pointer or status instruction).

Arithmetic and register (Type 2) instructions apply to the arithmetic and register circuit 20 only. There are 32 arithmetic and register instructions divided into eight classes encoded by the left-hand five bits of the instruction. Each of these instructions can be combined with any of eight word select signals to give a total capability of 256 instructions. The 32 arithmetic and register instructions are listed in the table below.

TABLE OF TYPE TWO INSTRUCTIONS
(in order of binary code)

CODE	INST	CODE	INST
0 0000	0→B	1 0000	A→B
0 0001	0→B	1 0001	B↔C
0 0010	A→C	1 0010	SHIFT C RIGHT
0 0011	C→1	1 0011	A→1
0 0100	B→C	1 0100	SHIFT B RIGHT
0 0101	0→C→C	1 0101	C+C→C
0 0110	0→C	1 0110	SHIFT A RIGHT
0 0111	0→C→1→C	1 0111	0→A
0 1000	SHIFT A LEFT	1 1000	A→B→A
0 1001	A→B	1 1001	A↔B
0 1010	A→C→C	1 1010	A→C→A
0 1011	C→1→C	1 1011	A→1→A
0 1100	C→A	1 1100	A+B→A
0 1101	0→C	1 1101	A↔C
0 1110	A+C→C	1 1110	A+C→A
0 1111	C+1→C	1 1111	A+1→A

KEY: A, B, C are registers; → means goes into; ↔ means interchange

The eight classes of arithmetic and register instructions are:

- (1) Clear (3);
- (2) Transfer/Exchange (6);
- (3) Add/Subtract (7);
- (4) Compare (6);
- (5) Complement (2);
- (6) Increment (2);
- (7) Decrement (2); and
- (8) Shift (4).

There are three clear instructions. These instructions are $0 \rightarrow A$, $0 \rightarrow B$, and $0 \rightarrow C$. They are implemented by simply disabling all the gates entering the designated register. Since these instructions can be combined with any of the eight word select options, it is possible to clear a portion of a register or a single digit.

There are six transfer/exchange instructions. These instructions are $A \rightarrow B$, $B \rightarrow C$, $C \rightarrow A$, $A \leftrightarrow B$, $B \leftrightarrow C$, and $C \leftrightarrow A$. This variety permits data in registers A, B, and C to be manipulated in many ways. Again, the power of the instruction must be viewed in conjunction with the word select option. Single digits can be exchanged or transferred.

There are seven add/subtract instructions which use the adder circuitry 84. They are $A+C \rightarrow C$, $A+B \rightarrow A$, $A+C \rightarrow A$, and $C+C \rightarrow C$. The last instruction can be

used to divide by five. This is accomplished by first adding the number to itself via $C+C \rightarrow C$, multiplying by two, then shifting right one digit, and dividing by ten. The result is a divide by five. This is used in the square root routine.

There are six compare instructions. These instructions are always followed by a conditional branch. They are used to check the value of a register or a single digit in a register and still not modify or transfer the contents. These instructions may easily be found in the type two instruction table above since there is no transfer arrow present. They are:

- (1) 0-B (Compare B to zero);
- (2) A-C (Compare A and C);
- (3) C-1 (Compare C to one);
- (4) 0-C (Compare C to zero);
- (5) A-B (Compare A and B); and
- (6) A-1 (Compare A to one).

If, for example, it is desired to branch if B is zero (or any digit or group of digits is zero as determined by WS), the 0-B instruction is followed by a conditional branch. If B was zero, no carry (or borrow) would be generated and the branch would occur. The instruction can be read: IF $U \geq V$ THEN BRANCH. Again it is easy to compare single digits or a portion of a register by appropriate word select options.

There are two complement instructions. The number representation system in the calculator is sign and magnitude notation for the mantissa, and tens complement notation in the exponent field. Before numbers can be subtracted, the subtrahend must be tens-complemented (i.e., $0-C \rightarrow C$). Other algorithms require the nines complement (i.e., $0-C-1 \rightarrow C$).

There are four increment/decrement instructions (two of each). They are $A \pm 1 \rightarrow A$ and $C \pm 1 \rightarrow C$.

There are four shift instructions. All three registers A, B, and C can be shifted right, while only A has a shift left capability. The arithmetic and register instruction set is summarized by class in the table below.

TABLE OF TYPE TWO INSTRUCTIONS
(divided by class)

Class	Instruction	Code
1) Clear	$0 \rightarrow A$	0111
	$0 \rightarrow B$	0001
	$0 \rightarrow C$	0110
2) Transfer/Exchange	$A \rightarrow B$	01001
	$B \rightarrow C$	00100
	$C \rightarrow A$	01100
	$A \leftrightarrow B$	11001
	$B \leftrightarrow C$	10001
	$C \leftrightarrow A$	11101
3) Add/Subtract	$A+C \rightarrow C$	01110
	$A-C \rightarrow C$	01010
	$A+B \rightarrow A$	11100
	$A-B \rightarrow A$	11000
	$A+C \rightarrow A$	11110
	$A-C \rightarrow A$	11010
	$C+C \rightarrow A$	10101
4) Compare	0-B	00000
	0-C	01101
	A-C	00010
	A-B	10000
	A-1	10011
	C-1	00011
5) Complement	$0-C \rightarrow C$	00101
	$0-C-1 \rightarrow C$	00111
6) Increment	$A+1 \rightarrow A$	11111
	$C+1 \rightarrow C$	01111
7) Decrement	$A-1 \rightarrow A$	11011
	$C-1 \rightarrow C$	01011
8) Shift	Sh A Right	10110
	Sh B Right	10100
	Sh C Right	10010
	Sh A Left	01000

The twenty eight-bit shift register **58-62** in control and timing circuit **16** contains twelve status bits or flags used to remember conditions of an algorithm or some past event (e.g., the decimal point key has already been depressed). These flags can be individually set, reset, or interrogated or all bits can be cleared (reset simultaneously). The format for the status operation (type **3**) instructions given in the instruction types table above is repeated below

TABLE OF STATUS INSTRUCTION DECODING

Bit No.	1 1 1 1 9 8 7 6	1 1 5 4	1 3	1 1 1 2 1 0
FIELD	N	F	0	1 0 0
F	INSTRUCTION			
0 0	SET FLAG N			
0 1	INTERROGATE FLAG N			
1 0	RESET FLAG N			
1 1	CLEAR ALL FLAGS (N=0000)			

If status bit N is one when the instruction "interrogate N" is executed, the CARRY flip-flop 66 in control and timing circuit 16 will be set. The status bit will remain set. Interrogate is always followed by a conditional branch instruction. The form of the interrogation is: "If status bit N=0, then branch," or "If status bit N \neq 1, then branch." The reason for this negative orientation is that all branches occur if the test is false (i.e., CARRY flip-flop=0), a result derived from using the conditional and unconditional branches as the same instruction.

Status bit 0 is set when a key is depressed. if cleared it will be set every word time as long as the key is down.

A four-bit counter **44** in control and timing circuit **16** acts as a pointer or marker to allow arithmetic instructions to operate on a portion of a register. Instructions are available to set and interrogate the pointed at one of fourteen locations or to increment or decrement the

present position. The pointer instruction decoding is given in the table below.

TABLE OF POINTER INSTRUCTION DECODING

BIT No.	9	8	7	6	5	4	3	2	1	0
FIELD	P				F		1	1	0	0
F	INSTRUCTION									
00	Set pointer to P									
10	Interrogate if pointer at P									
01	Decrement pointer									
11	Increment pointer									

P = XXXX
i.e. don't care

As with the status interrogate instruction, the CARRY flip-flop 66 is set if the pointer is at P when the "pointer at P?" instruction is executed (as with status interrogation, the actual question is in the negative form: IF $P \neq N$, THEN BRANCH or IF $P = \text{OTHER THAN } N$, THEN BRANCH). This instruction would be followed by a conditional branch. In a math routine the pointer allows progressive operation on a larger and larger portion of a word. After each iteration (cycle) through a loop, the pointer is decremented (or incremented) and then tested for completion to force another iteration or a jump out of the loop.

The data entry and display (type 5) instructions are used to enter data into arithmetic and register circuit 20, manipulate the stack and memory registers, and blank the display (sixteen instructions in this set are not recognized by any of the existing circuits and are therefore available for other external circuits that might be employed with other embodiments of the calculator). The table below contains a detailed coding of the data entry and display (type 5) instructions.

The first set of sixteen instructions ($I_5I_4 = 00$) in this table are not used by any of the main MOS circuits. They may be used by additional circuits or external circuitry listening to the I_5 line such as may be employed with other embodiments of the calculator.

The next instruction ($I_5I_4 = 01$) in this table is called the LOAD CONSTANT (LDC) or DIGIT ENTRY in-

TABLE OF TYPE 5 INSTRUCTION DECODING

(X = don't care, which in this context means the instruction does not depend on this bit; either a 1 or a 0 here will cause the same execution.)

	I ₉	I ₈	I ₇	I ₆	I ₅	I ₄	INSTRUCTION
	0000	→	1111	0	0		16 Available instructions
		^N					
10	0000	→	1001	0	1		Enters 4 bit code N into
							C Register at pointer position (LOAD POSITION)
	0	0	0	0	1	X	Display Toggle
	0	0	1	0	1	X	Exchange Memory, C → M → C
	0	1	0	0	1	X	Up Stack, C → C → D → E → F
8	0	1	1	0	1	X	Down Stack, F → F → E → D → A
	1	0	0	0	1	X	Display OFF
	1	0	1	0	1	X	Recall Memory, M → M → C
	1	1	0	0	1	X	Rotate Down, C → F → E → D → C
	1	1	1	0	1	X	Clear all registers 0 → A,B,C,D,E,F,M
1	X	X	0	1	1	X	I ₄ → A Register (56 bits)
1	X	X	1	1	1	X	BCD → C' register (56 bits)

struction. The four bits in $I_9 - I_6$ will be inserted into the C register at the location of the pointer, and the pointer will be decremented. This allows a constant, such as π (pi), to be stored in ROM and transferred to arithmetic and register circuit 20. To transfer a ten-digit constant requires only eleven instructions (one to preset the pointer). Several exclusions exist in the use of this instruction. When used with the pointer in position 13, it cannot be followed by an arithmetic and register instruction (i.e., by Type 2 or 5 instructions as there are problems in common use of the five-bit I_8 buffer 91 in arithmetic and register circuit 20). With $P=12$, LDC can be followed by another LDC but not by any other type 2 or 5 instruction. When used with the pointer in position 14, the instruction has no effect. However, when $P=12$ and LDC is followed by a type 2 or 5 instruction, position 13 in register C is modified. Loading non-digit codes (1010-1111) is not allowed since they will be modified passing through the adder. The next set of instructions ($I_6 I_5 I_4 = 01X$) in the type 5 instruction decoding table contains two display instructions and six stack or memory instructions. The display flip-flop in arithmetic and register circuit 20 controls blanking of all the LED's. When it is reset, the 1111 code is set into the display buffer 96, which is decoded so that no segments are on. There is one instruction to reset this flip-flop $I_9 I_8 I_7 = 100$ and another to toggle it (000). The toggle feature is convenient for blinking the display.

The remaining instructions in the type 5 instruction decoding table include two affecting memory (Exchange $C \leftrightarrow M$ and Recall $M \rightarrow C$), three affecting the stack (Up, Down, and Rotate Down), one general clear, one for loading register A from I_8 buss 28 (namely, $I_7 I_6 I_5 = 011$), and one for loading register C from BCD (111). Neither of the two last-mentioned instructions depends on bits I_9 , I_8 , or I_4 . The $I_8 \rightarrow A$ instruction is designed to allow a key code to be transmitted from a program storage circuit to arithmetic and register circuit 20 for display. The entire 56 bits are loaded although only two digits of information are of interest. The BCD $\rightarrow C$ instruction allows data input to arithmetic and register circuit 20 from a data storage circuit or other external source such as might be employed with other embodiments of the calculator.

The ROM select and other type six instructions are denoted by the pattern 10000 in instruction bits $I_4 - I_0$. The decoding table for these instructions is shown below.

The ROM SELECT instruction allows transfer of control from one ROM to another. Each ROM has a masking option which is programmed to decode bits $I_9 - I_7$. A Select ROM 3 instruction read from ROM 1 will reset the ROE flip-flop 70 in ROM 1 and set the ROE flip-flop 70 in ROM 3. The address is incremented in control and timing circuit 16 as usual. Thus, if Select ROM 3 is in location 197 in ROM 1, the first instruction read from ROM 3 will be location 198.

There are three ways to arrive at a desired address, on a different ROM as shown in FIG. 30. In path AA', a transfer (via an unconditional branch or a jump subroutine) to an address one before the desired address (L1) is executed in ROM N first. Then a ROM select M is given. In BB', the opposite order is shown (first ROM N select, then a transfer). Because the desired transfer location (L1 or L2) may already be occupied by an instruction, a third possibility may be used that is less efficient in states but does not depend on program locations. A transfer to L3 is made, then a ROM select, and then an additional transfer from L4 to the final desired location. With this method, L3 and L4 are overhead states.

Bits $I_6 I_5 = 01$ designate a subroutine return (RET). There are eight bits of storage in the twenty eight-bit shift register 58-62 of control and timing circuit 16 for retaining the return address when Jump Subroutine is executed. This address has already been incremented so execution of RET is simply a matter of outputting the address on I_4 line 32 at bit times $b_{19}-b_{26}$ and also inserting it into the ROM address portion 58 of the shift register. It is also still retained in the return address portion 60.

A key code is entered into control and timing circuit 16 by depressing a key on the keyboard. A key depression is detected by a positive interrogation of status bit 0. During a computation the keyboard is locked out because this status bit would ordinarily not be interrogated until return to the display loop. The actual key depression saves the state of the system counter (which is also the key code) in the key code buffer 56 (see FIG. 4) and also sets status bit 0. Execution of the KEYBOARD ENTRY instruction routes the key code (six bits) in the key code buffer 56 onto I_4 line 32 and into ROM address register 58 at bit times $b_{19}-b_{26}$. The most significant two bits b_{25} and b_{26} are set to zero so that a keyboard entry always jumps to one of the first 64 states.

TABLE OF TYPE SIX INSTRUCTION DECODING

Circuit Affected	$I_9 I_8 I_7 I_6 I_5 I_4 I_3 I_2 I_1 I_0$	Instruction
ROM	0 0 0 0 0 1 0 0 0 0	ROM select. One of 8 as specified in bits 19 - 17.
	↓ 0 0 1 0 0 0 0	
C&T	1 1 1 0 0 1 0 0 0 0	Subroutine return
	X X X 0 1 1 0 0 0 0	
	X X 0 1 0 1 0 0 0 0	External key code entry to C&T
	X X 1 1 0 1 0 0 0 0	Keyboard entry
DATA STORAGE	1 X 0 1 1 1 0 0 0 0	Send Address from C Register to Data Storage Circuit
	1 0 1 1 1 1 0 0 0 0	Send data from C Register into Data Storage Circuit

Two algorithms used in the calculator will now be discussed to further illustrate the instruction set. The first of these algorithms is a display wait loop, used after a key has been processed and while waiting for another key to be actuated. The second of these algorithms is a floating point multiply operation.

A flow diagram of the display wait loop is shown in FIG. 31. This loop is entered after a keystroke has been processed, register A has been properly loaded with the number to be displayed, and register B contains the display "mask" as discussed above. Two flags or status bits are required. Status bit 0 (SO) is hardwired in control and timing circuit 16 to automatically set whenever a key is down. Status bit 8 (S8) is used in this program to denote the fact that the key which is presently down has already been processed (since a routine may be finished before the key is released). In states D1S1 and D1S2 these two status bits are initialized. Then a loop is used as a time delay (about 14.4 ms.) to wait out any key bounce. In D1S4 status bit 8 (S8) is checked. The first time through the algorithm it must be 1 since it was set in D1S1 to denote the key has been processed. In state D1S5 the display is turned on (actually it is toggled since it must previously have been off; there is no DISPLAY ON instruction). At this time the answer appears to the user. In D1S6 status bit 0 (SO) is checked to see if a key is down. If not (i.e., SO=0), the previous key has been released, and status bit 8 (S8) is reset to 0 (D1S7). The machine is now ready to accept a new key since the previous key has been processed and released. The algorithm cycles through D1S6 and D1S7 waiting for a new key. This is the basic wait cycle of the calculator. If SO=1 in D1S6, the key which is down may be the old key (i.e., the one just processed) or a new key. This can be determined upon return to D1S4 where status bit 8 (S8) is checked. If a new key is down (S8=0), execution jumps to D1S8, the display is blanked, and a jump out is made to service the key. A listing of the algorithm is given in the table below.

TABLE OF WAIT LOOP ALGORITHM

LABEL	OPERATION	COMMENT
D1S1:	1 → S8	Set Status 8
D1S2:	0 → SO	Reset Status 0
D1S3:	P-1 → P	Decrement pointer, 48 word loop (3 × 16) to wait out key bounce
	IF P # 12	
	THEN GO TO D1S3	
	DISPLAY OFF	
D1S4:	IF S8 # 1	If key not processed, leave routine
	THEN GO TO D1S8:	
D1S5:	DISPLAY TOGGLE	Turn on display
D1S6:	IF SO # 1	If key up, reset S8 and wait
	THEN GO TO D1S7:	
	GO TO D1S2:	Key down. Check if same key
D1S7:	0 → S8	Indicate key not processed
	GO TO D1S6:	Back to wait for key
D1S8:		Blank display
D1S9:	KEYS → ROM ADDRESS	Jump to start of program to process key that was down.
	CONTINUE	

The floating point multiply algorithm multiplies x times y , where register C contains x in scientific notation and register D contains y (note that in the calculator register C corresponds to the user's X register and register D to the user's Y register). When the multiply key is depressed, the wait loop algorithm will jump instruction a ROM address corresponding to the first step of the multiply algorithm because of the way the instructions KEYS → ROM ADDRESS (state D1S9 in FIG. 31) is executed. The key code actually becomes the next ROM address. At this time the contents of registers A-D are indicated by the following:

- Register A floating point form of x
- Register B display mask for x
- Register C scientific form of x
- Register D scientific form of y

The algorithm for executing floating point multiply is given in the table below. The letters in parentheses indicate word select options as follows:

- P pointer position
- WP Up to a pointer position
- X Exponent field
- XS Exponent sign
- M mantissa field without sign
- MS mantissa with sign
- W entire word
- S mantissa sign only

TABLE OF FLOATING POINT MULTIPLY ALGORITHM

LABEL	OPERATION	COMMENT
MPY1:	STACK → A	Transfer y to A. Drop stack
MPY2:	A+C → C(X)	Add exponents to form exponent of answer
	A+C → C(S)	Add signs to form sign 1 of answer.
	IF NO CARRY GO TO MPY3	
MPY3:	0 → C(S)	Correct sign if both negative
	0 → B(W)	Clear B, then transfer mantissa of y. B(X) = 0.
	A → B(M)	Prepare A to accumulate product
	0 → A(W)	Set pointer to LSD (Least Significant Digit) Multiplier (Minus 1)
	2 → P	Increment to next digit.
MPY4:	P+1 → P	
MPY5:	A+B → A(W)	Add multiplier mantissa to partial product C(P) times. When C(P)=0, stop and go to next digit
	C-1 → C(P)	
	IF NO CARRY GO TO MPY5	
	SHIFT RIGHT A(W)	Shift partial product right.
	IF P # 12	Check if multiply is complete i.e. is pointer at MSD.
	THEN GO TO MPY4	
	IF A(P) > 1	Check if MSD = 0. If so must shift left and correct exp.
	THEN GO TO MPY6	
	SHIFT LEFT A(M)	Multiply by 10 and decrement exponent
	C-1 → C(X)	
MPY6:	C+1 → C(X)	Always do this to correct for factor of 10 too small
	A → B(XS)	Duplicate extra product digits
	A+B → A(XS)	add 11th digits
	IF NO CARRY GO TO MPY7	If sum less than 10, then done
	A+1 → A(M)	If sum more than 10, add 1
	IF NO CARRY GO TO MPY7	If answer was not all 9's, then done

TABLE OF FLOATING POINT MULTIPLY ALGORITHM - Continued

LABEL	OPERATION	COMMENT
MPY7	A+1 → A(P)	If answer was all 9's add 1 and increment exponent
	C+1 → C(X)	
	A EXCHANGE C(M)	Get answer mantissa into C
	GO TO MASK 1	Go to routine to position the answer in A and make the proper mask in B. Then to the DISPLAY program.

DETAILED LISTING OF ROUTINES AND SUBROUTINES OF INSTRUCTIONS

A complete listing of all of the routines and subroutines of instructions employed by the calculator and of all of the constants employed by these routines and subroutines is given below. All of these routines, subroutines, and constants are stored in ROM's 0-6, as indicated at the top of the first page associated with each ROM. Each line in each ROM is separately numbered in the first column from the left-hand side of the page. This facilitates reference to different parts of the list-

ting. Each address in ROM's 0-6 is represented in octal form by four digits in the second column from the left-hand side of the page. The first digit identifies which ROM, and the next three digits represent a nine-bit address (the L preceding these four digits is merely an address identifier). The instruction or constant stored in each address of ROM's 0-6 is represented in binary form in the third column from the left-hand side of the page. Branching addresses are represented in octal form by four digits in the fourth column from the left-hand side of the page. Explanatory comments are given in the remaining columns.

ROM 0

0	L0000:	1111.1.1	→	L0262	POWER1 :	JSB POWER2
1	L0001:	1111.1.11	→	L0134		GO TO ERZ
2	L0002:	1111.1.1.1			DIG3 :	A + 1 → A[X]
3	L0003:	1111.1.1.1			DIG2 :	A + 1 → A[X]
4	L0004:	1111.1.1.1			DIG1 :	A + 1 → A[X]
5	L0005:	11.1.1.11	→	L0044		IF NO CARRY GO TO DIG0
6	L0006:	11.1.1.1			MUL1 :	STACK → A
7	L0007:	11.1.1.1	→	L1010	*****	SELECT ROM 1
8	L0010:	11.1.1.1			MS1 :	DOWN ROTATE
9	L0011:	11.1.1.11	→	L0240		GO TO MS2
10	L0012:	11.1.1.1	→	L1013	*****	SELECT ROM 1
11	L0013:	1111.1.111	→	L0375	STOR1 :	GO TO STOR2
12	L0014:	11.1.1.1			RDOWN1 :	DOWN ROTATE
13	L0015:	11.1.1.111	→	L0017		GO TO OWFL3
14	L0016:	11.1.1.1			KEY :	STACK → A
15	L0017:	11.1.1.11	→	L0041		GO TO KEY1
16	L0020:	11.1.1.1				NO OPERATION
17	L0021:	11.1.1.1				NO OPERATION
18	L0022:	1111.1.1.1			DIG6 :	A + 1 → A[X]
19	L0023:	1111.1.1.1			DIG5 :	A + 1 → A[X]
20	L0024:	1111.1.1.1			DIG4 :	A + 1 → A[X]
21	L0025:	11.1.1.1	→	L0002		IF NO CARRY GO TO DIG3
22	L0026:	11.1.1.1			PLS1 :	STACK → A
23	L0027:	11.1.1.1	→	L1030	*****	SELECT ROM 1
24	L0030:	11.1.1.1			FV1 :	1 → S11
25	L0031:	1111.1.11	→	L0036		GO TO N1
26	L0032:	1111.1.11	→	L0036		GO TO N1
27	L0033:	11.1.1.1			PMT1 :	1 → S10
28	L0034:	1111.1.11	→	L0036		GO TO N1
29	L0035:	11.1.1.1			ROR1 :	RETURN
30	L0036:	1111.1.1.1			N1 :	IF S7 # 1
31	L0037:	11.1.1.11	→	L0060		THEN GO TO N2
32	L0040:	11.1.1.1	→	L3041	*****	SELECT ROM 3
33	L0041:	11.1.1.1	→	L1042	*****	SELECT ROM 1
34	L0042:	11.1.1.1	→	L1043	*****	SELECT ROM 1
35	L0043:	11.1.1.1			DPI :	1 → S5
36	L0044:	11.1.1.1	→	L1045	*****	SELECT ROM 1
37	L0045:	11.1.1.1			DIG0 :	NO OPERATION
38	L0046:	11.1.1.1			DIV1 :	STACK → A
39	L0047:	11.1.1.1	→	L1051	*****	SELECT ROM 1
40	L0050:	11.1.1.1	→	L6051	*****	SELECT ROM 6
41	L0051:	11.1.1.1			DATE :	NO OPERATION
42	L0052:	11.1.1.1	→	L4053	*****	SELECT ROM 4
43	L0053:	11.1.1.1	→	L4054	*****	SELECT ROM 4
44	L0054:	11.1.1.1			PRC1 :	STACK → A
45	L0055:	11.1.1.1	→	L1056	*****	SELECT ROM 1
46	L0056:	11.1.1.1			PRE :	1 → S7
47	L0057:	11.1.1.11	→	L0102		GO TO STOR3
48	L0060:	11.1.1.1			N2 :	1 → S7
49	L0061:	11.1.1.111	→	L0117		GO TO OWFL3
50	L0062:	1111.1.1.1			DIG9 :	A + 1 → A[X]
51	L0063:	1111.1.1.1			DIG8 :	A + 1 → A[X]
52	L0064:	1111.1.1.1			DIG7 :	A + 1 → A[X]
53	L0065:	11.1.1.11	→	L0022		IF NO CARRY GO TO DIG6
54	L0066:	1111.1.1.1			MIN1 :	0 - C - 1 + C[S]
55	L0067:	11.1.1.1	→	L0026		JSB PLS1
56	L0070:	11.1.1.1			CLR1 :	0 → C[W]
57	L0071:	11.1.1.1	→	L0104		JSB CLR2
58	L0072:	1111.1.1.11	→	L0366		GO TO CHS2
59	L0073:	11.1.1.1			RCL1 :	IF S8 # 1
60	L0074:	11.1.1.11	→	L0114		THEN GO TO RCL3
61	L0075:	11.1.1.111	→	L0115		GO TO RCL4
62	L0076:	11.1.1.1			ENTER1 :	C → STACK
63	L0077:	11.1.1.1				0 → S11
64	L0100:	11.1.1.1				0 → S10
65	L0101:	1111.1.1				0 → S7

ROM O—Continued

66	L0102:	1.1.1.1.		STOR3	:	0 → S4
67	L0103:	1.1.1.1.1	→	GO TO OWFL1		
68	L0104:	1.1.1.1.1.		CLR2	:	IF S7 # 1
69	L0105:	1.1.1.1.1.1	→			THEN GO TO CLR3
70	L0106:	1.1.1.1.				IF S4 # 1
71	L0107:	1.1.1.1.1.1	→			THEN GO TO CLR4
72	L0110:	1.1.1.1.1.1	→			GO TO CLR3
73	L0111:	1.1.1.1.		CLR4	:	C → STACK
74	L0112:	1.1.1.1.				C → STACK
75	L0113:	1.1.1.1.1.1	→			GO TO ENTER1
76	L0114:	1.1.1.1.		RCL3	:	C → STACK
77	L0115:	1.1.1.1.		RCL4	:	M → C
78	L0116:	1.1.1.1.		OWFL7	:	0 → S7
79	L0117:	1.1.1.1.		OWFL3	:	1 → S4
80	L0120:	1.1.1.1.		OWFL1	:	0 → S8
81	L0121:	1.1.1.1.1.		OWFL4	:	C → A[W]
82	L0122:	1.1.1.1.1.		OWFL5	:	A → B[W]
83	L0123:	1.1.1.1.				12 → P
84	L0124:	1.1.1.1.1.				0 → C[MS]
85	L0125:	1.1.1.1.1.				C + 1 → C[P]
86	L0126:	1.1.1.1.1.				C + 1 → C[P]
87	L0127:	1.1.1.1.1.				IF C[X] = 0
88	L0130:	1.1.1.1.1.1.1	→			THEN GO TO MSK1
89	L0131:	1.1.1.1.1.				0 - C - 1 → C[X]
90	L0132:	1.1.1.1.1.				IF C[X] = 0
91	L0133:	1.1.1.1.1.1	→			THEN GO TO MSK2
92	L0134:	1.1.1.1.		ERZ	:	1 → S5
93	L0135:	1.1.1.1.1.				A + B → A[X]
94	L0136:	1.1.1.1.1.1	→			IF NO CARRY GO TO OWFL2
95	L0137:	1.1.1.1.1.				0 → C[W]
96	L0140:	1.1.1.1.	→			JSB OWFL1
97	L0141:	1.1.1.1.1.		OWFL2	:	0 → C[W]
98	L0142:	1.1.1.1.1.				C - 1 → C[WP]
99	L0143:	1.1.1.1.1.				0 → C[X]
100	L0144:	1.1.1.1.1.1.				A EXCHANGE C[S]
101	L0145:	1.1.1.1.1.	→			GO TO OWFL1
102	L0146:	1.1.1.1.		DIS4	:	0 → S4
103	L0147:	1.1.1.1.		DIS3	:	0 → S9
104	L0150:	1.1.1.1.				12 → P
105	L0151:	1.1.1.1.		DIS5	:	0 → S0
106	L0152:	1.1.1.1.		DIS6	:	P - 1 → P
107	L0153:	1.1.1.1.1.				IF P # 12
108	L0154:	1.1.1.1.1.1	→			THEN GO TO DIS6
109	L0155:	1.1.1.1.				DISPLAY OFF
110	L0156:	1.1.1.1.				IF S9 # 1
111	L0157:	1.1.1.1.1.1	→			THEN GO TO DIS7
112	L0160:	1.1.1.1.				0 → S5
113	L0161:	1.1.1.1.				SHIFT LEFT A[X]
114	L0162:	1.1.1.1.		TKR	:	KEYS → ROM ADDRESS
115	L0163:	1.1.1.1.		DIS9	:	1 → S9
116	L0164:	1.1.1.1.				1 → P
117	L0165:	1.1.1.1.				IF S5 # 1
118	L0166:	1.1.1.1.1.1	→			THEN GO TO DIS10
119	L0167:	1.1.1.1.1.				C + 1 → C[WP]
120	L0170:	1.1.1.1.1.1	→			IF NO CARRY GO TO DIS9
121	L0171:	1.1.1.1.		DIS7	:	DISPLAY TOGGLE
122	L0172:	1.1.1.1.		DIS10	:	IF S0 # 1
123	L0173:	1.1.1.1.1.1	→			THEN GO TO DIS9
124	L0174:	1.1.1.1.1.1	→			GO TO DIS5
125	L0175:	1.1.1.1.1.		SCINT9	:	A → B[X]
126	L0176:	1.1.1.1.1.1	→			GO TO SCINT4
127	L0177:	1.1.1.1.1.	→	MSK1	:	JSB SROUND
128	L0200:	1.1.1.1.1.				0 → A[X]
129	L0201:	1.1.1.1.1.1.				A + 1 → 0 A[X]
130	L0202:	1.1.1.1.				SHIFT LEFT A[X]
131	L0203:	1.1.1.1.1.				A EXCHANGE B[W]
132	L0204:	1.1.1.1.				IF A ≤ B[X]
133	L0205:	1.1.1.1.1.1	→			THEN GO TO SCINT9
134	L0206:	1.1.1.1.				A → B[X]
135	L0207:	1.1.1.1.1.		MSK11	:	A - 1 → A[X]
136	L0210:	1.1.1.1.1.1	→			IF NO CARRY GO TO MSK12
137	L0211:	1.1.1.1.		MSK14	:	IF P # 3
138	L0212:	1.1.1.1.1.1	→			THEN GO TO MSK13
139	L0213:	1.1.1.1.	→	MSK15	:	SELECT ROM 1
140	L0214:	1.1.1.1.		MSK28	:	P - 1 → P
141	L0215:	1.1.1.1.1.1	→			GO TO MSK14
142	L0216:	1.1.1.1.		CLR3	:	0 → S7
143	L0217:	1.1.1.1.				0 → S4
144	L0220:	1.1.1.1.1.1	→			GO TO OWFL4
145	L0221:	1.1.1.1.		MSK2	:	C → A[X]
146	L0222:	1.1.1.1.1.	→			JSB SROUND
147	L0223:	1.1.1.1.1.1.				A + 1 → A[X]
148	L0224:	1.1.1.1.1.		MSK21	:	A - 1 → A[X]
149	L0225:	1.1.1.1.1.1	→			IF NO CARRY GO TO MSK22
150	L0226:	1.1.1.1.1.		MSK13	:	C - 1 → C[X]
151	L0227:	1.1.1.1.1.	→			IF NO CARRY GO TO MSK28
152	L0230:	1.1.1.1.1.1	→			GO TO MSK15
153	L0231:	1.1.1.1.1.		MSK22	:	SHIFT RIGHT A[M]
154	L0232:	1.1.1.1.1.	→			JSB MSK21
155	L0233:	1.1.1.1.		MSK12	:	P - 1 → P
156	L0234:	1.1.1.1.1.				SHIFT RIGHT C[M]
157	L0235:	1.1.1.1.1.	→			JSB MSK11
158	L0236:	1.1.1.1.	→	*****	SROUND	: SELECT ROM 1
159	L0237:	1.1.1.1.				RETURN

ROM O—Continued

160	L0240:	..1.1....	→	L1241	*****	MS2	:	SELECT ROM 1
161	L0241:	1....1.1.				SCINT3	:	IF A >= B[X]
162	L0242:	1.1..1.1.1	→	L0245				THEN GO TO SCINT4
163	L0243:	1.1.1.1.1.				SCINT2	:	A + 1 → A[X]
164	L0244:	1.1.1.1.1.						A - 1 → A[X]
165	L0245:	..1.1.1..				SCINT4	:	0 → C[X]
166	L0246:	..1.1.1..						3 → P
167	L0247:	1.1..1.1..				SCINT7	:	IF P ≠ 12
168	L0250:	1.1.1.1.1.1	→	L0361				THEN GO TO SCINT5
169	L0251:	1....1.1.1.				SCINT6	:	B EXCHANGE C[W]
170	L0252:	..1.1.1..				DIS1	:	0 → S6
171	L0253:	..1.1.1.1.	→	L0147				JSB DIS3
172	L0254:	1.1.1.1.1.						0 → A[MS]
173	L0255:	..1.1.1..				DENT1	:	IF S4 ≠ 1
174	L0256:	1.1....1.1	→	L0260				THEN GO TO DENT2
175	L0257:	..1.1.1..						C → STACK
176	L0260:	..1.1.1..				DENT2	:	0 → S6
177	L0261:	1.1.1.1.1.1	→	L0351				GO TO DENT3
178	L0262:	1.1.1.1.1..				POWER2	:	CLEAR REGISTERS
179	L0263:1.1.1..						CLEAR STATUS
180	L0264:	..1.1.1..						1 → S2
181	L0265:	..1.1.1.1.1	→	L0111				GO TO CLR4
182	L0266:	1.1.1.1.1.				DENT8	:	A + 1 → A[X]
183	L0267:1.1..						IF B[M] = 0
184	L0270:	1.1..1.1.1.1	→	L0346				THEN GO TO DENT18
185	L0271:	..1.1.1.1..						IF P ≠ 3
186	L0272:	1.1.1.1.1.1.1	→	L0277				THEN GO TO DENT5
187	L0273:	..1.1.1.1..						0 → C[W]
188	L0274:	1.1.1.1.1.1.1						C + 1 → C[S]
189	L0275:	1.1.1.1.1.1.1						C + 1 → C[S]
190	L0276:	1.1.1.1.1..						13 → P
191	L0277:	1.1.1.1.1.				DENT5	:	SHIFT RIGHT C[WP]
192	L0300:1.1..				DENT19	:	IF B[M] = 0
193	L0301:	1.1..1.1.1.1.1	→	L0313				THEN GO TO DENT10
194	L0302:	1.1....1.1..						12 → P
195	L0303:	1.1.1.1.1.						IF A[P] >= 1
196	L0304:	1.1..1.1.1.1.1	→	L0313				THEN GO TO DENT10
197	L0305:	1.1.1.1.1.1.						0 → A[X]
198	L0306:	1.1.1.1.1.1.				DENT11	:	A - 1 → A[X]
199	L0307:	1.1.1.1.1.						IF A[P] >= 1
200	L0310:	1.1..1.1.1.1.1	→	L0313				THEN GO TO DENT10
201	L0311:	..1.1.1.1..						SHIFT LEFT A[M]
202	L0312:	1.1....1.1.1	→	L0306				GO TO DENT11
203	L0313:	..1.1.1.1..				DENT10	:	A → B[X]
204	L0314:	1.1..1.1.1.						B EXCHANGE C[W]
205	L0315:	1.1.1.1.1.1.						A EXCHANGE C[W]
206	L0316:	..1.1.1..						0 → S5
207	L0317:	..1.1.1.1.	→	L0146				JSB DIS4
208	L0320:	1.1..1.1.1.						B EXCHANGE C[W]
209	L0321:	..1.1.1..						0 → C[X]
210	L0322:	1.1.1.1.1.1.						C + 1 → C[MS]
211	L0323:1.1..				DENT12	:	0 → P
212	L0324:	..1.1.1.1.						C - 1 → C[P]
213	L0325:1.1.1..				DENT4	:	P + 1 → P
214	L0326:	..1.1.1.1..						C - 1 → C[P]
215	L0327:	1.1.1.1.1.1.1	→	L0332				IF NO CARRY GO TO DENT6
216	L0330:	..1.1.1..						SHIFT LEFT A[WP]
217	L0331:	1.1.1.1.1.1.1	→	L0325				GO TO DENT4
218	L0332:	1.1..1.1.1.				DENT6	:	A EXCHANGE B[X]
219	L0333:	..1.1.1.1.						A → B[W]
220	L0334:	..1.1.1.1.						IF S5 ≠ 1
221	L0335:	1.1.1.1.1.1	→	L0340				THEN GO TO DENT7
222	L0336:	..1.1.1..						1 → S6
223	L0337:	1.1..1.1.1.1.1	→	L0313				GO TO DENT10
224	L0340:	..1.1.1..				DENT7	:	IF S6 ≠ 1
225	L0341:	1.1.1.1.1.1	→	L0266				THEN GO TO DENT8
226	L0342:1.1.1..						P - 1 → P
227	L0343:	..1.1.1.1.						IF P ≠ 2
228	L0344:	1.1.1.1.1.1.1.1.1	→	L0277				THEN GO TO DENT5
229	L0345:	1.1....1.1	→	L0300				GO TO DENT19
230	L0346:	1.1..1.1.1.				DENT18	:	SHIFT RIGHT C[W]
231	L0347:	1.1..1.1.1.						B EXCHANGE C[W]
232	L0350:	1.1..1.1.1.	→	L0146				JSB DIS4
233	L0351:	..1.1.1.1..				DENT3	:	0 → C[W]
234	L0352:1.1.1..						0 → B[W]
235	L0353:	1.1.1.1.1..						13 → P
236	L0354:	..1.1.1..						LOAD CONSTANT 3
237	L0355:	1....1.1.1.						0 → S8
238	L0356:	..1.1.1.1..						C - 1 → C[X]
239	L0357:	1....1.1.1.						B EXCHANGE C[X]
240	L0360:	1.1.1.1.1.1.1	→	L0323				GO TO DENT12
241	L0361:	1.1..1.1.1.				SCINT5	:	IF A[P] >= 1
242	L0362:	1.1.1.1.1.1	→	L0251		L0363:		THEN GO TO SCINT6
243	L0363:	..1.1.1.1.						C - 1 → C[P]
244	L0364:1.1.1..						P + 1 → P
245	L0365:	1.1..1.1.1.1.1	→	L0247				GO TO SCINT7
246	L0366:	1.1.1.1.1.1.1.				CHS2	:	0 - C - 1 → C[S]
247	L0367:	..1.1.1.1..						C → A[S]
248	L0370:	..1.1.1.1.						C → A[X]
249	L0371:	1.1..1.1.1.1	→	L0147				GO TO DIS3
250	L0372:1.1.1..						NO OPERATION
251	L0373:1.1.1..						NO OPERATION
252	L0374:1.1.1..						NO OPERATION
253	L0375:	..1.1.1..				STOR2	:	C EXCHANGE M
254	L0376:	1.1.1.1.1..						M → C
255	L0377:	..1.1.1.1.	→	L0102				GO TO STOR3

0	L1000:			NO OPERATION
1	L1001:			NO OPERATION
2	L1002:	...1...1..		R3 :	1 → S1
3	L1003:	...1111.11	→ L1066	R2 :	GO TO R12
4	L1004:	...1...1..		R1 :	1 → S1
5	L1005:	...1111.111	→ L1115		GO TO R13
6	L1006:	...1111.1..		XTY :	0 → S9
7	L1007:	...1...1..	→ L2010	*****	SELECT ROM 2
8	L1010:	...1111.1..1	→ L1346	SMUL11 :	JSB MPY
9	L1011:	...1111.111	→ L1115		GO TO R13
10	L1012:	...1...1..	→ L2013	*****	SELECT ROM 2
11	L1013:	...1...1..		SQR1 :	1 → S8
12	L1014:	...1111.1..		XTY11 :	IF S7 # 1
13	L1015:	...1111.11	→ L1026		THEN GO TO XTY12
14	L1016:	...1111.1..			0 → S7
15	L1017:	...1111.1..1	→ L1374		JSB SQR
16	L1020:	...1111.111	→ L1115		GO TO R13
17	L1021:	...1...1..	→ L4022	*****	SELECT ROM 4
18	L1022:	...1111.111	→ L1065	RETR4 :	GO TO R11
19	L1023:	...1...1..		R6 :	1 → S1
20	L1024:	...1...1..		R5 :	1 → S1
21	L1025:	...1111.111	→ L1115	R4 :	1 → S3
22	L1026:	...1111.1..	→ L1006		GO TO R13
23	L1027:	...1111.1..	→ L1114	XTY12 :	JSB XTY
24	L1030:	...1111.1..			GO TO R14
25	L1031:	...1111.1..1	→ L1154		0 → S7
26	L1032:	...1111.111	→ L1115	ADD11 :	JSB ADD
27	L1033:	...1111.1..			GO TO R13
28	L1034:	...1...1..	→ L0035	DIG10 :	0 → S7
29	L1035:	...1...1..		*****	SELECT ROM 0
30	L1036:	...1...1..		RET11 :	RETURN
31	L1037:	...1111.11	→ L1060	DIG11 :	IF S4 # 1
32	L1040:	...1111.111	→ L1033		THEN GO TO DIG14
33	L1041:	...1...1..			GO TO DIG10
34	L1042:	...1111.11111	→ L1317		NO OPERATION
35	L1043:	...1111.111	→ L1133	XEY :	GO TO MS17
36	L1044:	...1111.111	→ L1115	SUM11 :	GO TO SUM12
37	L1045:	...1111.1..		R0 :	GO TO R13
38	L1046:	...1111.111	→ L1033		IF S7 # 1
39	L1047:	...1111.11	→ L1036		THEN GO TO DIG10
40	L1050:	...1111.1..		SDIV11 :	GO TO DIG11
41	L1051:	...1111.111	→ L1076		0 → S7
42	L1052:	...1...1..			GO TO DI
43	L1053:	...1111.1..		PRC2 :	C → STACK
44	L1054:	...1...1..	→ L1100		IF S7 # 1
45	L1055:	...1111.11	→ L1070		THEN GO TO PRC4
46	L1056:	...1111.111		PRC11 :	GO TO PRC3
47	L1057:	...1111.11	→ L1052		A EXCHANGE C[W]
48	L1060:	...1111.1..			GO TO PRC2
49	L1061:	...1111.1..		DIG14 :	CLEAR STATUS
50	L1062:	...1...1..		TKRR1 :	KEYS → ROM ADDRESS
51	L1063:	...1...1..		R9 :	NO OPERATION
52	L1064:	...1...1..		R8 :	NO OPERATION
53	L1065:	...1111.1..		R7 :	1 → S1
54	L1066:	...1...1..		R11 :	1 → S3
55	L1067:	...1111.111	→ L1115	R12 :	1 → S2
56	L1070:	...1111.1..			GO TO R13
57	L1071:	...1111.11.1	→ L1153	PRC3 :	0 → S7
58	L1072:	...1111.1..			JSB SUB
59	L1073:	...1111.1..			DOWN ROTATE
60	L1074:	...1111.1..			C → STACK
61	L1075:	...1111.1..			C - 1 → C[X]
62	L1076:	...1111.1..1	→ L1344		C - 1 → C[X]
63	L1077:	...1111.1..	→ L1115	DI :	JSB DIV
64	L1100:	...1111.1..1	→ L1346		JSB R13
65	L1101:	...1111.1..		PRC4 :	JSB MPY
66	L1102:	...1111.1..			C - 1 → C[X]
67	L1103:	...1111.1..	→ L1115		C - 1 → C[X]
68	L1104:	...1111.1..			JSB R13
69	L1105:	...1111.1..		R100 :	DOWN ROTATE
70	L1106:	...1111.1..			C - 1 → C[X]
71	L1107:	...1111.111		ONE :	C - 1 → C[X]
72	L1110:	...1111.111			0 → A[W]
73	L1111:	...1111.111			A + 1 → A[S]
74	L1112:	...1111.111			SHIFT RIGHT A[W]
75	L1113:	...1111.111	→ L1337		A EXCHANGE C[W]
76	L1114:	...1111.1..			GO TO RTN16
77	L1115:	...1111.1..	→ L0116	*****	STACK → A
78	L1116:	...1111.1..		R14 :	SELECT ROM 0
79	L1117:	...1111.1..		R13 :	0 → C[S]
80	L1120:	...1111.1..		MSK20 :	0 → C[XS]
81	L1121:	...1111.1..	→ L1130		C + C → C[P]
82	L1122:	...1111.1..			IF NO CARRY GO TO
83	L1123:	...1111.111			MSK16
84	L1124:	...1111.111			B → C[WP]
85	L1125:	...1111.111	→ L1130		C + 1 → C[W]
86	L1126:	...1111.1..			IF C[S] = 0
87	L1127:	...1111.1..			THEN GO TO MSK16
88	L1130:	...1111.111		MSK16 :	SHIFT RIGHT C[MS]
89	L1131:	...1111.111			SHIFT RIGHT B[MS]
90	L1132:	...1111.1..1	→ L1250		A EXCHANGE C[W]
91	L1133:	...1111.1..			C → A[S]
92	L1134:	...1111.111	→ L1255	SUM12 :	GO TO MSKR0
					IF S7 # 1
					THEN GO TO SUM13

ROM 1—Continued

93	L1135:	1.1.1.1.1.				IF S4 # 1
94	L1136:	1.1.1.1.1.1	→ L1252			THEN GO TO SUM14
95	L1137:	1.1.1.1.1.	→ L3140	*****		SELECT ROM 3
96	L1140:	1.1.1.1.1.		R	:	0 → C[X]
97	L1141:	1.1.1.1.1.				IF S3 # 1
98	L1142:	1.1.1.1.1.1	→ L1227			THEN GO TO RND1
99	L1143:	1.1.1.1.1.1				C + 1 → C[X]
100	L1144:	1.1.1.1.1.1				C + C → C[X]
101	L1145:	1.1.1.1.1.1				C + C → C[X]
102	L1146:	1.1.1.1.1.		RN3	:	IF S2 # 1
103	L1147:	1.1.1.1.1.1	→ L1227			THEN GO TO RND1
104	L1150:	1.1.1.1.1.				IF S1 # 1
105	L1151:	1.1.1.1.1.1	→ L1234			THEN GO TO RND3
106	L1152:	1.1.1.1.1.1	→ L1240			GO TO SCIN
107	L1153:	1.1.1.1.1.1		SUB	:	0 - C - 1 → C[S]
108	L1154:	1.1.1.1.1.1		ADD	:	A → B[W]
109	L1155:	1.1.1.1.1.1		ADD1	:	A + 1 → A[XS]
110	L1156:	1.1.1.1.1.1				A + 1 → A[XS]
111	L1157:	1.1.1.1.1.1				C + 1 → C[XS]
112	L1160:	1.1.1.1.1.1				C + 1 → C[XS]
113	L1161:	1.1.1.1.1.				IF A >= C[X]
114	L1162:	1.1.1.1.1.1	→ L1164			THEN GO TO ADD4
115	L1163:	1.1.1.1.1.1				A EXCHANGE C[W]
116	L1164:	1.1.1.1.1.1		ADD4	:	IF A[M] >= 1
117	L1165:	1.1.1.1.1.1.1	→ L1367			THEN GO TO ADD2
118	L1166:	1.1.1.1.1.1	→ L1172			GO TO ADD7
119	L1167:	1.1.1.1.1.1		ADD3	:	SHIFT RIGHT A[M]
120	L1170:	1.1.1.1.1.				IF A[M] >= 1
121	L1171:	1.1.1.1.1.1	→ L1370			THEN GO TO ADD5
122	L1172:	1.1.1.1.1.1		ADD7	:	C - 1 → C[XS]
123	L1173:	1.1.1.1.1.1				C - 1 → C[XS]
124	L1174:	1.1.1.1.1.1				0 → A[X]
125	L1175:	1.1.1.1.1.1.1				A EXCHANGE C[S]
126	L1176:	1.1.1.1.1.1.1				A - C → A[S]
127	L1177:	1.1.1.1.1.1				IF A[S] >= 1
128	L1200:	1.1.1.1.1.1	→ L1204			THEN GO TO ADD8
129	L1201:	1.1.1.1.1.1				A + C → A[MS]
130	L1202:	1.1.1.1.1.1				A - C → A[S]
131	L1203:	1.1.1.1.1.	→ L2204	*****		SELECT ROM 2
132	L1204:	1.1.1.1.1.		ADD8	:	A - C → C[M]
133	L1205:	1.1.1.1.1.1	→ L1207			IF NO CARRY GO TO ADD9
134	L1206:	1.1.1.1.1.				0 - C → C[MS]
135	L1207:	1.1.1.1.1.		ADD9	:	C → A[M]
136	L1210:	1.1.1.1.1.	→ L2211	*****	ADD10	SELECT ROM 2
137	L1211:	1.1.1.1.1.				NO OPERATION
138	L1212:	1.1.1.1.1.				NO OPERATION
139	L1213:	1.1.1.1.1.				NO OPERATION
140	L1214:	1.1.1.1.1.1				C + 1 → C[P]
141	L1215:	1.1.1.1.1.				0 → C[X]
142	L1216:	1.1.1.1.1.				C - 1 → C[WP]
143	L1217:	1.1.1.1.1.1				B EXCHANGE C[W]
144	L1220:	1.1.1.1.1.1				A EXCHANGE C[W]
145	L1221:	1.1.1.1.1.				P - 1 → P
146	L1222:	1.1.1.1.1.1	→ L1116			GO TO MSK20
147	L1223:	1.1.1.1.1.				NO OPERATION
148	L1224:	1.1.1.1.1.				NO OPERATION
149	L1225:	1.1.1.1.1.				NO OPERATION
150	L1226:	1.1.1.1.1.				NO OPERATION
151	L1227:	1.1.1.1.1.		RND1	:	IF S1 # 1
152	L1230:	1.1.1.1.1.1	→ L1232			THEN GO TO RND2
153	L1231:	1.1.1.1.1.1				C + 1 → C[X]
154	L1232:	1.1.1.1.1.		RND2	:	IF S2 # 1
155	L1233:	1.1.1.1.1.1	→ L1236			THEN GO TO RND4
156	L1234:	1.1.1.1.1.1		RND3	:	C + 1 → C[X]
157	L1235:	1.1.1.1.1.1				C + 1 → C[X]
158	L1236:	1.1.1.1.1.	→ L0237	*****	RND4	SELECT ROM 0
159	L1237:	1.1.1.1.1	→ L1140			GO TO R
160	L1240:	1.1.1.1.1.	→ L0241	*****	SCIN	SELECT ROM 0
161	L1241:	1.1.1.1.1.		MS11	:	DOWN ROTATE
162	L1242:	1.1.1.1.1.				IF S7 # 1
163	L1243:	1.1.1.1.1.1	→ L1351			THEN GO TO MS12
164	L1244:	1.1.1.1.1.				0 → S7
165	L1245:	1.1.1.1.1.1	→ L1331			JSB ROT1
166	L1246:	1.1.1.1.1.				C → STACK
167	L1247:	1.1.1.1.1	→ L1010			GO TO SMULL11
168	L1250:	1.1.1.1.1.		MSKR0	:	0 → S6
169	L1251:	1.1.1.1.1.	→ L0252	*****		SELECT ROM 0
170	L1252:	1.1.1.1.1.1		SUM14	:	0 - C - 1 → C[S]
171	L1253:	1.1.1.1.1.				0 → S7
172	L1254:	1.1.1.1.1.				1 → S4
173	L1255:	1.1.1.1.1.		SUM13	:	STACK → A
174	L1256:	1.1.1.1.1.				0 → S8
175	L1257:	1.1.1.1.1.				C → STACK
176	L1260:	1.1.1.1.1.1	→ L1154			JSB ADD
177	L1261:	1.1.1.1.1.				DOWN ROTATE
178	L1262:	1.1.1.1.1.				C → A[W]
179	L1263:	1.1.1.1.1.1	→ L1346			JSB MPY
180	L1264:	1.1.1.1.1.				IF S4 # 1
181	L1265:	1.1.1.1.1.1.1	→ L1267			THEN GO TO SUM16
182	L1266:	1.1.1.1.1.1.				0 - C - 1 → C[S]
183	L1267:	1.1.1.1.1.		SUM16	:	STACK → A
184	L1270:	1.1.1.1.1.				C → STACK
185	L1271:	1.1.1.1.1.1				A EXCHANGE C[W]
186	L1272:	1.1.1.1.1.	→ L1107			JSB ONE

ROM 1—Continued

187	L1273:	1...1...1			IF S4 # 1
188	L1274:	111111.11	→	L1276	THEN GO TO SUM15
189	L1275:	1.111111.			C - 1 → C[S]
190	L1276:	11.11...1	→	L1154	SUM15 : JSB ADD
191	L1277:	11.11...			DOWN ROTATE
192	L1300:	11.11...			STACK → A
193	L1301:	11.11...1	→	L1154	JSB ADD
194	L1302:	1.1.1...			C → STACK
195	L1303:	11.11...			DOWN ROTATE
196	L1304:	11.11...			DOWN ROTATE
197	L1305:	1.11.111	→	L1115	GO TO R13
198	L1306:	1.111.1	→	L1107	MS14 : JSB ONE
199	L1307:	11.11.11.1	→	L1153	JSB SUB
200	L1310:	1...1.111.			B EXCHANGE C[W]
201	L1311:	11.11...			DOWN ROTATE
202	L1312:	111.1111.			A EXCHANGE C[W]
203	L1313:	111.1...1	→	L1344	JSB DIV
204	L1314:	1...1...			I → S8
205	L1315:	111111...1	→	L1374	JSB SQR
206	L1316:	11.11...			STACK → A
207	L1317:	1.1.1...			MS17 : C → STACK
208	L1320:	111.1111.			A EXCHANGE C[W]
209	L1321:	1.11.111	→	L1115	GO TO R13
210	L1322:	11.11...			CSN : DOWN ROTATE
211	L1323:	11.11...			DOWN ROTATE
212	L1324:	11111111.			0 - C - 1 → C[S]
213	L1325:	11.11...			DOWN ROTATE
214	L1326:	11.11...			DOWN ROTATE
215	L1327:	11.1111111	→	L1337	GO TO RTN16
216	L1330:			NO OPERATION
217	L1331:	1...1.111.			ROT1 : B EXCHANGE C[W]
218	L1332:	11.11...			DOWN ROTATE
219	L1333:	11.11...			STACK → A
220	L1334:	1...1.111.			B EXCHANGE C[W]
221	L1335:	1.1.1...			C → STACK
222	L1336:	1.1.111.			B → C[W]
223	L1337:	1.1.1...			RTN16 : DISPLAY OFF
224	L1340:	1111.1...11	→	L1364	GO TO RTN9
225	L1341:	1.1.1...			RTN8 : IF S4 # 1
226	L1342:	...1.111	→	L1021	
227	L1343:	11.1...	→	L3344	***** RETR3 : SELECT ROM 3
228	L1344:	111.11.11.			DIV : A EXCHANGE C[MS]
229	L1345:	111.11111	→	L1347	GO TO DIV12
230	L1346:	1.1.1...	→	L2347	***** MPY : SELECT ROM 2
231	L1347:	1.1.1.1.			DIV12 : A - C → C[X]
232	L1350:	1.1.1...	→	L2351	***** : SELECT ROM 2
233	L1351:	11.11...			MS12 : STACK → A
234	L1352:	11.11.1.1	→	L1331	JSB ROT1
235	L1353:	1.1.1...			C → STACK
236	L1354:	111.1111.			A EXCHANGE C[W]
237	L1355:	111.1...1	→	L1344	JSB DIV
238	L1356:	11.11...			DOWN ROTATE
239	L1357:	111.11.1	→	L1346	JSB MPY
240	L1360:	11.1.1...			STACK → A
241	L1361:	11.11.11.1	→	L1153	JSB SUB
242	L1362:	11.11.1.1	→	L1331	JSB ROTL
243	L1363:	11.11.11	→	L1306	GO TO MS14
244	L1364:	111.1.1.			RTN9 : IF S7 # 1
245	L1365:	...111.111	→	L1035	THEN GO TO RET11
246	L1366:	111...111	→	L1341	GO TO RTN8
247	L1367:	111.1111.			ADD2 : A EXCHANGE C[W]
248	L1370:	...1.1.1.			ADD5 : IF A >= C[X]
249	L1371:	1111.11.11	→	L1172	THEN GO TO ADD7
250	L1372:	11111.1.1.			A + 1 → A[X]
251	L1373:	111.11111	→	L1167	IF NO CARRY GO TO ADD3
252	L1374:	...11.11.			SQR : IF C[M] >= 1
253	L1375:	...1.1.11	→	L1012	THEN GO TO SQR1
254	L1376:	...11...			RETURN

ROM 2

0	L2000:	...1...	→	L0001	*****	ERR21 : SELECT ROM 0
1	L2001:	111...111.				PMU23 : A + B → A[W]
2	L2002:	1.111111.				PMU24 : C - 1 → C[S]
3	L2003:111	→	L2001		IF NO CARRY GO TO PMU23
4	L2004:	111.1111.				A EXCHANGE C[W]
5	L2005:	1...1.11.				SHIFT LEFT A[MS]
6	L2006:	111.1111.				A EXCHANGE C[W]
7	L2007:	1111.11	→	L2074		GO TO PQO23
8	L2010:	11.1.1...				XTY21 : STACK → A
9	L2011:	1.1.1...				C → STACK
10	L2012:	111.1111.				A EXCHANGE C[W]
11	L2013:	1.111.111.				LN22 : 0 → A[W]
12	L2014:	11...1.				1 → S6
13	L2015:	11.1...1.				A - C → A[M]
14	L2016:11	→	L2000		IF NO CARRY GO TO ERR21
15	L2017:	1.11.111.				SHIFT RIGHT A[W]
16	L2020:	1.111111.				C - 1 → C[S]
17	L2021:11	→	L2000		IF NO CARRY GO TO ERR21
18	L2022:	11111111.				LN25 : C + 1 → C[S]
19	L2023:	1.1.111.				LN26 : A → B[W]

20	L2024:	...1.11.11	→	L2026		GO TO ECA22
21	L2025:	1.11.1.1.			ECA21 :	SHIFT RIGHT A[WP]
22	L2026:	1.1.11.11.			ECA22 :	A - 1 → A[S]
23	L2027:	...1.1.111	→	L2025		IF NO CARRY GO TO ECA21
24	L2030:	1.11.11.11.				0 → A[S]
25	L2031:	111...111.				A + B → A[W]
26	L2032:	1.11.1.1.				IF S6 # 1
27	L2033:	1.11.1.11	→	L2132		THEN GO TO EXP29
28	L2034:	1.11.1.1.				A - 1 → A[P]
29	L2035:	...1.1.11	→	L2022		IF NO CARRY GO TO LN25
30	L2036:	1.1.1.1.				A EXCHANGE B[WP]
31	L2037:	1.1.1.1.				SHIFT LEFT A[WP]
32	L2040:	111...111.				A + B → A[S]
33	L2041:	111...111	→	L2341		IF NO CARRY GO TO LN24
34	L2042:	1.11.1.1.				7 → P
35	L2043:	...1111.11	→	L2074		GO TO PQO23
36	L2044:	1.1111.1.1.			PRE23 :	A + 1 → A[X]
37	L2045:	1.1.11.1.			PRE29 :	IF A[XS] >= 1
38	L2046:	1.1.11.11.11	→	L2337		THEN GO TO PRE27
39	L2047:	1.1...1.1.			PRE24 :	A - B → A[MS]
40	L2050:	...1.1.11	→	L2044		IF NO CARRY GO TO PRE23
41	L2051:	1.11.1.11.				A + B → A[MS]
42	L2052:	1...111.				SHIFT LEFT A[W]
43	L2053:	1.11.1.1.				C - 1 → C[X]
44	L2054:	...1.1.111	→	L2045		IF NO CARRY GO TO PRE29
45	L2055:	1.11.1.11.			PRE25 :	SHIFT RIGHT A[W]
46	L2056:	...1.1.1.				0 → C[WP]
47	L2057:	1.11.1.1.1.			PRE26 :	A EXCHANGE C[X]
48	L2060:	1.1.11.11.				IF C[S] = 0
49	L2061:	...1.1.111	→	L2065		THEN GO TO PRE28
50	L2062:	1.1.1.11.				A EXCHANGE B[W]
51	L2063:	1.1...111.				A - B → A[W]
52	L2064:	...111.111.				0 - C - 1 → C[W]
53	L2065:	1.11.1.11.			PRE28 :	SHIFT RIGHT A[W]
54	L2066:	1...1.1.				0 → S8
55	L2067:	...1111.11	→	L2074		GO TO PQO23
56	L2070:	1.11.11.11.			PQO15 :	C + 1 → C[S]
57	L2071:	1.1...111.			PQO16 :	A - B → A[W]
58	L2072:	...111...11	→	L2070		IF NO CARRY GO TO PQO15
59	L2073:	1.11...111.				A + B → A[W]
60	L2074:	1...1.111.			PQO23 :	B EXCHANGE C[W]
61	L2075:	...11.111.				0 → C[W]
62	L2076:	1.11...11.				C - 1 → C[M]
63	L2077:	1...11...				LOAD CONSTANT 4
64	L2100:	1.1111.11.			PQO24 :	C + 1 → C[M]
65	L2101:	1.1.1.11.				SHIFT RIGHT C[W]
66	L2102:	1.11.11...				IF P # 5
67	L2103:	1.11...111	→	L2321		THEN GO TO EXP35
68	L2104:	1.11...11.				6 → P
69	L2105:	1.1111.1.				0 → A[WP]
70	L2106:	1.1.1.11.				13 → P
71	L2107:	1.1.1.11.				B EXCHANGE C[W]
72	L2110:	1.11.1.11.				A EXCHANGE C[W]
73	L2111:	1.1.11...				LOAD CONSTANT 6
74	L2112:	1.11...111	→	L2141		GO TO EXP23
75	L2113:	1.11.1.11.			EXP32 :	IF P # 11
76	L2114:	1.1.1.1.11	→	L2246		THEN GO TO EXP31
77	L2115:	1.11.11...			LNC2 :	LOAD CONSTANT 6
78	L2116:	1.1.1.11...				LOAD CONSTANT 9
79	L2117:	1.11.11...				LOAD CONSTANT 3
80	L2120:	...1.11...				LOAD CONSTANT 1
81	L2121:	1...1.11...				LOAD CONSTANT 4
82	L2122:	1.11.11...				LOAD CONSTANT 7
83	L2123:	...1.11...				LOAD CONSTANT 1
84	L2124:	1...1.11...				LOAD CONSTANT 8
85	L2125:	...11...				LOAD CONSTANT 0
86	L2126:	1.1.1.11...				LOAD CONSTANT 5
87	L2127:	1.1.1.11...				LOAD CONSTANT 6
88	L2130:	1.11.1.11.				11 → P
89	L2131:	1.11.1.11.11	→	L2327		GO TO LN35
90	L2132:	1.1111.1.			EXP29 :	A + 1 → A[P]
91	L2133:	1.1.1.111.			EXP32 :	A → B[W]
92	L2134:	1.11.11.11.				C - 1 → C[S]
93	L2135:	...1.1.1.11	→	L2026		IF NO CARRY GO TO ECA22
94	L2136:	1.11.1.1.				SHIFT RIGHT A[WP]
95	L2137:	1.11.1.11.				A EXCHANGE C[W]
96	L2140:	1...1.11.			EXP23 :	SHIFT LEFT A[MS]
97	L2141:	1.11.1.11.				A EXCHANGE C[W]
98	L2142:	1.1.11.11.				A - 1 → A[S]
99	L2143:	1.11.1.11.11	→	L2133		IF NO CARRY GO TO EXP22
100	L2144:	1.1.1.11.				A EXCHANGE B[W]
101	L2145:	1.1111...1.				A + 1 → A[P]
102	L2146:	1...1.111	→	L2211		IF NO CARRY GO TO NRM21
103	L2147:	1.1.1.1.1.			PQO21 :	SHIFT RIGHT C[MS]
104	L2150:	1...111.				SHIFT LEFT A[W]
105	L2151:	...111.111	→	L2071		GO TO PQO16
106	L2152:	1.1.1.11.			EXP34 :	IF P # 9
107	L2153:	1.11.1.1.11	→	L2164		THEN GO TO EXP33
108	L2154:	1.11.1.11.			LNCD2 :	7 → P
109	L2155:	1.11.11...				LOAD CONSTANT 3
110	L2156:	1.11.11...				LOAD CONSTANT 3
111	L2157:	...11...				LOAD CONSTANT 0
112	L2160:	1...11...				LOAD CONSTANT 8
113	L2161:	1.1.1.11...				LOAD CONSTANT 5

114	L2162:	1.1.1.11..			9 → P
115	L2163:	1.1.1.1111	→	L2327	GO TO LN35
116	L2164:	1.1.1.11..		EXP33 :	IF P # 10
117	L2165:	1.1.1.1111	→	L2113	THEN GO TO EXP32
118	L2166:	1.1.1.11..		LNCD1 :	9 → P
119	L2167:	1.1.1.11..			LOAD CONSTANT 3
120	L2170:	1.1.1.11..			LOAD CONSTANT 1
121	L2171:	1.1.1.11..			LOAD CONSTANT 0
122	L2172:	1.1.1.11..			LOAD CONSTANT 1
123	L2173:	1.1.1.11..			LOAD CONSTANT 7
124	L2174:	1.1.1.11..			LOAD CONSTANT 9
125	L2175:	1.1.1.11..			LOAD CONSTANT 8
126	L2176:	1.1.1.11..			LOAD CONSTANT 1
127	L2177:	1.1.1.11..			10 → P
128	L2200:	1.1.1.1111	→	L2327	GO TO LN35
129	L2201:	1.1.1.11..		MPY26 :	A + B → A[MS]
130	L2202:	1.1.1.11..		MPY27 :	C - 1 → C[P]
131	L2203:	1.1.1.11..	→	L2201	IF NO CARRY GO TO MPY26
132	L2204:	1.1.1.11..		MPY28 :	SHIFT RIGHT A[W]
133	L2205:	1.1.1.11..			P + 1 → P
134	L2206:	1.1.1.11..			IF P # 13
135	L2207:	1.1.1.11..	→	L2202	THEN GO TO MPY27
136	L2210:	1.1.1.1.1.			C + 1 → C[X]
137	L2211:	1.1.1.1111		NRM21 :	0 → A[S]
138	L2212:	1.1.1.11..			12 → P
139	L2213:	1.1.1.11..		NRM23 :	IF A[P] >= 1
140	L2214:	1.1.1.11..	→	L2222	THEN GO TO NRM24
141	L2215:	1.1.1.11..			SHIFT LEFT A[W]
142	L2216:	1.1.1.1.1.			C - 1 → C[X]
143	L2217:	1.1.1.11..			IF A[W] >= 1
144	L2220:	1.1.1.11..	→	L2213	THEN GO TO NRM23
145	L2221:	1.1.1.11..			0 → C[W]
146	L2222:	1.1.1.1.1.		NRM24 :	A EXCHANGE C[X]
147	L2223:	1.1.1.11.1.			C + C → C[XS]
148	L2224:	1.1.1.11.1.	→	L2226	IF NO CARRY GO TO NRM29
149	L2225:	1.1.1.11.1.			A + 1 → A[MS]
150	L2226:	1.1.1.1.1.		NRM29 :	A EXCHANGE C[X]
151	L2227:	1.1.1.11.1.			IF A[S] >= 1
152	L2230:	1.1.1.11..	→	L2204	THEN GO TO MPY28
153	L2231:	1.1.1.11..			A EXCHANGE C[M]
154	L2232:	1.1.1.11..		NRM25 :	C → A[W]
155	L2233:	1.1.1.11..			IF S8 # 1
156	L2234:	1.1.1.11.1.	→	L2336	THEN GO TO NRM26
157	L2235:	1.1.1.11..			IF S6 # 1
158	L2236:	1.1.1.11.1.	→	L2246	THEN GO TO EXP31
159	L2237:	1.1.1.11..			0 → S6
160	L2240:	1.1.1.11..			IF S9 # 1
161	L2241:	1.1.1.11.1.	→	L2345	THEN GO TO XTY32
162	L2242:	1.1.1.11..			0 → C[W]
163	L2243:	1.1.1.11..			P - 1 → P
164	L2244:	1.1.1.11..			LOAD CONSTANT 5
165	L2245:	1.1.1.11.1.	→	L2347	GO TO MPY21
166	L2246:	1.1.1.11..		EXP31 :	IF P # 12
167	L2247:	1.1.1.11.1.	→	L2327	THEN GO TO LN35
168	L2250:	1.1.1.11..			0 → C[W]
169	L2251:	1.1.1.11..		LNC10 :	LOAD CONSTANT 2
170	L2252:	1.1.1.11..			LOAD CONSTANT 3
171	L2253:	1.1.1.11..			LOAD CONSTANT 0
172	L2254:	1.1.1.11..			LOAD CONSTANT 2
173	L2255:	1.1.1.11..			LOAD CONSTANT 5
174	L2256:	1.1.1.11..			LOAD CONSTANT 8
175	L2257:	1.1.1.11..			LOAD CONSTANT 5
176	L2260:	1.1.1.11..			LOAD CONSTANT 0
177	L2261:	1.1.1.11..			LOAD CONSTANT 9
178	L2262:	1.1.1.11..			LOAD CONSTANT 3
179	L2263:	1.1.1.11..			12 → P
180	L2264:	1.1.1.11.1.			A EXCHANGE C[W]
181	L2265:	1.1.1.11..			IF S6 # 1
182	L2266:	1.1.1.11..	→	L2310	THEN GO TO PRE21
183	L2267:	1.1.1.11..			A - C → C[W]
184	L2270:	1.1.1.11..			IF B[XS] = 0
185	L2271:	1.1.1.11.1.	→	L2273	THEN GO TO LN27
186	L2272:	1.1.1.11..			A - C → C[W]
187	L2273:	1.1.1.11.1.		LN27 :	A EXCHANGE B[W]
188	L2274:	1.1.1.11..		LN28 :	P - 1 → P
189	L2275:	1.1.1.11..			SHIFT LEFT A[W]
190	L2276:	1.1.1.11..			IF P # 1
191	L2277:	1.1.1.11.1.	→	L2274	THEN GO TO LN28
192	L2300:	1.1.1.11.1.			A EXCHANGE C[W]
193	L2301:	1.1.1.11.1.			IF C[S] = 0
194	L2302:	1.1.1.11..	→	L2304	THEN GO TO LN29
195	L2303:	1.1.1.11..			0 - C - 1 → C[M]
196	L2304:	1.1.1.1.1.		LN29 :	C + 1 → C[X]
197	L2305:	1.1.1.11..			DISPLAY TOGGLE
198	L2306:	1.1.1.11..			11 → P
199	L2307:	1.1.1.11..	→	L2202	GO TO MPY27
200	L2310:	1.1.1.11..		PRE21 :	A → B[W]
201	L2311:	1.1.1.11..			C → A[M]
202	L2312:	1.1.1.11.1.			C + C → C[XS]
203	L2313:	1.1.1.11.1.	→	L2047	IF NO CARRY GO TO PRE24
204	L2314:	1.1.1.11.1.			C + 1 → C[XS]
205	L2315:	1.1.1.11..		PRE22 :	SHIFT RIGHT A[W]
206	L2316:	1.1.1.1.1.			C + 1 → C[X]
207	L2317:	1.1.1.11.1.	→	L2315	IF NO CARRY GO TO PRE22

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208	L2320:	..11....11	→	L2060		GO TO PRE26
209	L2321:	1....11..			EXP35 :	IF P # 8
210	L2322:	..11.1.11	→	L2152		THEN GO TO EXP34
211	L2323:	..11.1.1..			LNCD3 :	5 → P
212	L2324:	..11.11...				LOAD CONSTANT 3
213	L2325:	..11.11...				LOAD CONSTANT 3
214	L2326:	1....11..				8 → P
215	L2327:	1....1.11.			LN35 :	B EXCHANGE C[W]
216	L2330:	..11.1.1..				IF S6 # 1
217	L2331:	..11.11111	→	L2147		THEN GO TO PQO21
218	L2332:	..11.1.11.				SHIFT RIGHT A[W]
219	L2333:1111..				P + 1 → P
220	L2334:1111..				P + 1 → P
221	L2335:1.11	→	L2002		GO TO PMU24
222	L2336:	..1.1....	→	L1337	*****	SELECT ROM 1
223	L2337:	11111.11.			NRM26 :	A + 1 → A[M]
224	L2340:	..11.1.11	→	L2055		IF NO CARRY GO TO PRE25
225	L2341:	11.11111.			LN24 :	A EXCHANGE B[S]
226	L2342:	1.1.1.11.				SHIFT RIGHT C[MS]
227	L2343:	11111111.				A + 1 → A[S]
228	L2344:1.1111	→	L2023		IF NO CARRY GO TO LN26
229	L2345:	11.1.1.1..			XTY32 :	DOWN ROTATE
230	L2346:	..1.1.1..				C → STACK
231	L2347:	..11.1.1..			MPY21 :	3 → P
232	L2350:	..111.1.1.			MPY22 :	A + C → C[X]
233	L2351:	..1.1.111.			DIV21 :	A - C → C[S]
234	L2352:	111.11.11	→	L2354		IF NO CARRY GO TO DIV22
235	L2353:	..1.11111.				0 - C → C[S]
236	L2354:	11.1.111.			DIV22 :	A EXCHANGE B[W]
237	L2355:	1.111.111.				0 → A[W]
238	L2356:	..1.11111.				A → B[S]
239	L2357:	11.1.1.1..				IF P # 12
240	L2360:	1....1.11	→	L2202		THEN GO TO MPY27
241	L2361:11..				IF B[M] = 0
242	L2362:11..	→	L2000		THEN GO TO ERR21
243	L2363:	111.11.1.			DIV23 :	A EXCHANGE C[WP]
244	L2364:	1111.11.11	→	L2366		GO TO DIV15
245	L2365:	..1111....			DIV14 :	C + 1 → C[P]
246	L2366:	11.1.1.1.			DIV15 :	A - B → A[MS]
247	L2367:	1111.1.11	→	L2365		IF NO CARRY GO TO DIV14
248	L2370:	111.1.11.				A + B → A[MS]
249	L2371:	..1.1.1.1.				SHIFT LEFT A[MS]
250	L2372:111..			DIV16 :	P - 1 → P
251	L2373:1.11..				IF P # 0
252	L2374:	1111.11.11	→	L2366		THEN GO TO DIV15
253	L2375:	..11.1111.				C → A[S]
254	L2376:	111.1.111.				A EXCHANGE C[W]
255	L2377:	1....1.11	→	L2211		GO TO NRM21

ROM 3

0	L3000:	1....11.1	→	L3106	FVR47	JSB ONE
1	L3001:	111.1.111.				A EXCHANGE C[W]
2	L3002:	111....11.	→	L3343		JSB DIV
3	L3003:1.1.11	→	L3024		GO TO FVR48
4	L3004:	1....1..			XTY :	1 → S8
5	L3005:	..1.1....	→	L1006	*****	SELECT ROM 1
6	L3006:	1....1.1.			LN :	0 → S8
7	L3007:1.111	→	L3011		GO TO SQR1
8	L3010:	1....1..				1 → S8
9	L3011:	..1.1....	→	L1012	*****	SELECT ROM 1
10	L3012:	111.1.1.1	→	L3345	SQR1 :	JSB MPY
11	L3013:	1....11.1	→	L3106	N46 :	JSB ONE
12	L3014:	1.11.1.1.				IF S11 # 1
13	L3015:	..1111.111	→	L3075		THEN GO TO N42
14	L3016:	..11.1.11.	→	L3153		JSB ADD
15	L3017:11.1.	→	L3006	N44 :	JSB LN
16	L3020:	11.1.1.1..				DOWN ROTATE
17	L3021:11.1.	→	L3006		JSB LN
18	L3022:	..1.1.1.1	→	L3106		JSB ONE
19	L3023:	..111.111	→	L3131		GO TO N48
20	L3024:	..11.1.1..			FVR48 :	STACK → A
21	L3025:	..1.1.1..				C → STACK
22	L3026:	11.11.1.1	→	L3331		JSB S12
23	L3027:	..11111.11	→	L3136		GO TO P47
24	L3030:	1.111.1..			FV40 :	0 → S11
25	L3031:	11.11.1.1	→	L3314		GO TO FV46
26	L3032:	..1.1.1.1	→	L3104	PV40 :	GO TO PV41
27	L3033:	11.1.1.11	→	L3312	PMT40 :	GO TO PMT42
28	L3034:	..11.11.11	→	L3155	ROR40 :	GO TO FVR
29	L3035:				NO OPERATION
30	L3036:	..11.111.			N40 :	C → A[W]
31	L3037:	11.1.1.1..				DOWN ROTATE
32	L3040:	1111.1.111	→	L3365		GO TO N41
33	L3041:	..1.1.1..			N5 :	IF S4 # 1
34	L3042:	1111.1.11	→	L3364		THEN GO TO SELR4
35	L3043:	..11.1..			TKRR3 :	KEYS → ROM ADDRESS
36	L3044:				NO OPERATION
37	L3045:	11.1.1.1..			CASH1 :	DOWN ROTATE
38	L3046:	..1.1.1..				C EXCHANGE M
39	L3047:	11.1.1.1..				DOWN ROTATE

ROM 3—Continued

40	L3050:	1..1.1...				C → STACK
41	L3051:	1..1.1.1	→ L3103			JSB R100
42	L3052:	1.1.1.1.1	→ L3153			JSB ADD
43	L3053:	1.1.1.1...				STACK → A
44	L3054:	1..1.1...				C → STACK
45	L3055:	1.1.1.1.1.1				A EXCHANGE C[W]
46	L3056:	1..1.1.1	→ L3106			JSB ONE
47	L3057:	1.1.1.1.1	→ L3153			JSB ADD
48	L3060:1..1	→ L3004			JSB XTY
49	L3061:	1.1.1.1...				DOWN ROTATE
50	L3062:	1..1.1...				C → STACK
51	L3063:	1.1.1.1...				DOWN ROTATE
52	L3064:	1.1.1.1...				DOWN ROTATE
53	L3065:	1.1.1.1.1				A EXCHANGE C[W]
54	L3066:	1.1.1.1.1	→ L3343			JSB DIV
55	L3067:	1.1.1.1...				M → C
56	L3070:	1.1.1.1.1	→ L3153			JSB ADD
57	L3071:	1.1.1.1...		R13	:	0 → S10
58	L3072:	1.1.1.1...				0 → S11
59	L3073:	1.1.1.1.1	→ L3115			GO TO R14
60	L3074:				NO OPERATION
61	L3075:	1.1.1.1.1.1		N42	:	A EXCHANGE C[W]
62	L3076:	1.1.1.1.1	→ L3152			JSB SUB
63	L3077:	1.1.1.1.1.1				A EXCHANGE B[W]
64	L3100:	1.1.1.1.1	→ L3343			JSB DIV
65	L3101:1.1.1.1	→ L3017			GO TO N44
66	L3102:				NO OPERATION
67	L3103:	1..1.1...	→ L1104	*****	R100	SELECT ROM 1
68	L3104:	1.1.1.1...			PV41	0 → S11
69	L3105:	1..1.1.1	→ L3110			GO TO PV42
70	L3106:	1..1.1...	→ L1107	*****	ONE	SELECT ROM 1
71	L3107:	1.1.1.1...			PV46	1 → S11
72	L3110:	1.1.1.1.1	→ L3321		PV42	JSB CSN
73	L3111:	1..1.1.1	→ L3103		PV49	JSB R100
74	L3112:	1.1.1.1.1	→ L3153			JSB ADD
75	L3113:	1.1.1.1...	→ L3330			JSB ROT1
76	L3114:	1..1.1.1.1	→ L3116			GO TO PV48
77	L3115:1...	→ L0116	*****	R14	SELECT ROM 0
78	L3116:1..1	→ L3004		PV48	JSB XTY
79	L3117:	1..1.1.1	→ L3106			JSB ONE
80	L3120:	1.1.1.1.1.1				A EXCHANGE C[W]
81	L3121:	1.1.1.1...			PV43	IF S10 # 1
82	L3122:	1.1.1.1.1.1	→ L3134			THEN GO TO PV44
83	L3123:	1.1.1.1.1	→ L3152			JSB SUB
84	L3124:	1.1.1.1...				DOWN ROTATE
85	L3125:	1.1.1.1...				DOWN ROTATE
86	L3126:	1.1.1.1...				DOWN ROTATE
87	L3127:	1.1.1.1.1...				IF S11 # 1
88	L3130:	1.1.1.1.1.1	→ L3132			THEN GO TO PV45
89	L3131:	1.1.1.1.1.1		N48	:	A EXCHANGE C[W]
90	L3132:	1.1.1.1.1	→ L3343		PV45	JSB DIV
91	L3133:	1.1.1.1...				STACK → A
92	L3134:	1.1.1.1...			PV44	STACK → A
93	L3135:	1.1.1.1...				STACK → A
94	L3136:	1.1.1.1.1	→ L3345		P47	JSB MPY
95	L3137:	1.1.1.1.1.1	→ L3071			GO TO R13
96	L3140:	1..1.1.1...		CASH	:	IF S10 # 1
97	L3141:	1..1.1.1.1	→ L3045			THEN GO TO CASH1
98	L3142:	1..1.1.1...				0 → S4
99	L3143:	1..1.1...	→ L4144	*****		SELECT ROM 4
100	L3144:	1.1.1.1.1		FVR49	:	C + 1 → C[X]
101	L3145:	1.1.1.1.1	→ L3343			JSB DIV
102	L3146:	1..1.1.1	→ L3106			JSB ONE
103	L3147:	1.1.1.1.1	→ L3153			JSB ADD
104	L3150:	1.1.1.1...				STACK → A
105	L3151:	1.1.1.1.1.1.1	→ L3167			GO TO FVR43
106	L3152:	1..1.1...	→ L1153	*****	SUB	SELECT ROM 1
107	L3153:	1..1.1...	→ L1154	*****	ADD	SELECT ROM 1
108	L3154:	1..1.1...	→ L1155	*****	ADD1	SELECT ROM 1
109	L3155:	1.1.1.1.1...		FVR	:	IF S11 # 1
110	L3156:	1.1.1.1.1.1	→ L3340			THEN GO TO FV42
111	L3157:	1.1.1.1.1...				0 → S11
112	L3160:	1.1.1.1.1...				IF S10 # 1
113	L3161:	1.1.1.1.1.1.1	→ L3273			THEN GO TO FVR1
114	L3162:	1.1.1.1.1	→ L3321	FVR3	:	JSB CSN
115	L3163:	1.1.1.1.1.1.1				0 - C - 1 → C[S]
116	L3164:	1..1.1.1...		FVR2	:	C → STACK
117	L3165:	1.1.1.1...				STACK → A
118	L3166:	1.1.1.1...				DOWN ROTATE
119	L3167:	1.1.1.1.1	→ L3343	FVR43	:	JSB DIV
120	L3170:	1..1.1.1...				C EXCHANGE M
121	L3171:	1.1.1.1...				DOWN ROTATE
122	L3172:	1..1.1.1...				C → STACK
123	L3173:	1..1.1.1	→ L3106			JSB ONE
124	L3174:	1.1.1.1.1.1	→ L3153			JSB ADD
125	L3175:	1.1.1.1.1.1				A EXCHANGE B[W]
126	L3176:	1.1.1.1.1.1	→ L3345			JSB MPY
127	L3177:	1..1.1.1...				C EXCHANGE M
128	L3200:	1.1.1.1...				STACK → A
129	L3201:	1..1.1.1...				C → STACK
130	L3202:	1.1.1.1.1	→ L3321			JSB CSN
131	L3203:	1.1.1.1...				DOWN ROTATE
132	L3204:	1.1.1.1.1	→ L3152			JSB SUB
133	L3205:	1.1.1.1.1.1	→ L3153			JSB ADD

ROM 3—Continued

134	L3206:	..1.1.1...			C EXCHANGE M
135	L3207:	111...11.1	→L3343		JSB DIV
136	L3210:	..1.1.1...		FVR44 :	C EXCHANGE M
137	L3211:	11.1.1.1...			DOWN ROTATE
138	L3212:	..1.1.1...			C EXCHANGE M
139	L3213:	..1.1.1...			STACK → A
140	L3214:	..1...11.1	→L3106		JSB ONE
141	L3215:	..11.1.11.1	→L3153		JSB ADD
142	L3216:	..1.1.1...			C → STACK
143	L3217:	..1.1.1...			C → STACK
144	L3220:	..1...111.			B → C[W]
145	L3221:	..1.1.1...			C EXCHANGE M
146	L3222:1	→L3004		JSB XTY
147	L3223:	11.1.1.1...			DOWN ROTATE
148	L3224:	..1.1.1...			C → STACK
149	L3225:	111...11.1	→L3345		JSB MPY
150	L3226:	..1...111.			B → C[W]
151	L3227:	11.1.1.1...			DOWN ROTATE
152	L3230:	11.1.1.1...			DOWN ROTATE
153	L3231:	111...11.1	→L3343		JSB DIV
154	L3232:	11.1.1.1...			DOWN ROTATE
155	L3233:	11.1.1.1...			DOWN ROTATE
156	L3234:	..1...11.1	→L3106		JSB ONE
157	L3235:	..11.1.1.1	→L3152		JSB SUB
158	L3236:	1.1.1.1...			M → C
159	L3237:	111...11.1	→L3343		JSB DIV
160	L3240:	11.1.1.1...			DOWN ROTATE
161	L3241:	11.1.1.1...			DOWN ROTATE
162	L3242:	..11.1.1.1	→L3152		JSB SUB
163	L3243:	11.1.1.1...			STACK → A
164	L3244:	..1.1.1...			C → STACK
165	L3245:	1...1.111.			B EXCHANGE C[W]
166	L3246:	..11.1.11.1	→L3153		JSB ADD
167	L3247:	11.1.1.1...			DOWN ROTATE
168	L3250:	1...1.111.			B EXCHANGE C[W]
169	L3251:	11.1.1.1...			DOWN ROTATE
170	L3252:	1...1.111.			B EXCHANGE C[W]
171	L3253:	111...11.1	→L3343		JSB DIV
172	L3254:	1.1.1.1...			M → C
173	L3255:	111...11.1	→L3345		JSB MPY
174	L3256:	..1.1.1...			C EXCHANGE M
175	L3257:	..11.1.11.1	→L3153		JSB ADD
176	L3260:	..1.1.1...			C EXCHANGE M
177	L3261:	..11...111.			C → A[W]
178	L3262:	111.111...1	→L3354		JSB TEN6
179	L3263:	1.1111.1...			IF A[X] ≥ 1
180	L3264:	11.1.1.11	→L3322		THEN GO TO FVR46
181	L3265:	1...1.111	→L3211		GO TO FVR44
182	L3266:	111...11.1	→L3345	FVR49 :	JSB MPY
183	L3267:	11.11...1.1	→L3331		JSB S12
184	L3270:	..111.1.1.1			C + 1 → C[X]
185	L3271:	1.11...1...			1 → S11
186	L3272:	..11.1.1.1	→L3144		GO TO FVR49
187	L3273:	..1.1.1...		FVR1 :	STACK → A
188	L3274:	111.1.111.			A EXCHANGE C[W]
189	L3275:	111...11.1	→L3343		JSB DIV
190	L3276:	..1.1.1...			STACK → A
191	L3277:	..1.1.1...			C → STACK
192	L3300:	111.1.111.			A EXCHANGE C[W]
193	L3301:	..1...1.1	→L3106		JSB ONE
194	L3302:	111.1.111.			A EXCHANGE C[W]
195	L3303:	111...11.1	→L3343		JSB DIV
196	L3304:1	→L3304		JSB XTY
197	L3305:	..1...1.1	→L3016		JSB ONE
198	L3306:	..11.1.1.1	→L3152		JSB SUB
199	L3307:	..1111.1.1			C + 1 → C[X]
200	L3310:	..1111.1.1			C + 1 → C[X]
201	L3311:	..111.1.1	→L3071		JSB R13
202	L3312:	1.11.1.1...		PNT42 :	IF S11 ≠ 1
203	L3313:	..1...11111	→L3107		THEN GO TO PV46
204	L3314:	1...1.1.1		FV46 :	IF S10 ≠ 1
205	L3315:	..1.1.1.11	→L3111		THEN GO TO PV49
206	L3316:	..1111111.			0 - C - 1 → C[S]
207	L3317:	..1.1.1.1	→L3111		JSB PV49
208	L3320:			NO OPERATION
209	L3321:	..1.1.1...	→L3322	***** CSN :	SELECT ROM 1
210	L3322:	1.11.1.1...		FVR46 :	IF S11 ≠ 1
211	L3323:	..111...111	→L3071		THEN GO TO R13
212	L3324:	11.1.1.1...			DOWN ROTATE
213	L3325:	11.1.1.1...			DOWN ROTATE
214	L3326:	11.1.1.1...			DOWN ROTATE
215	L3327:11	→L3000		GO TO FVR47
216	L3330:	..1.1.1...	→L1331	***** ROT1 :	SELECT ROM 1
217	L3331:	..11.111.		S12 :	0 → C[W]
218	L3332:	..111...1.			C + 1 → C[P]
219	L3333:	..1111111.			C + 1 → C[S]
220	L3334:	1.1.1.1.1			C + C → C[WP]
221	L3335:	1.1.1.1.1			SHIFT RIGHT C[MS]
222	L3336:	..111.1.1			C + 1 → C[X]
223	L3337:11			RETURN
224	L3340:	1.1.1.1.1		FV42 :	IF S10 ≠ 1
225	L3341:	111...11.1	→L3346		THEN GO TO FVR4
226	L3342:	..111.1.1	→L3164		GO TO FVR2
227	L3343:	..1.1.1...	→L1344	***** DIV :	SELECT ROM 1

ROM 3—Continued

228	L3344:11....		TRN16	:	RETURN
229	L3345:	...1.1....	→L1346	*****	MPY	: SELECT ROM 1
230	L3346:	111.1111.		FVR4	:	A EXCHANGE C[W]
231	L3347:	11.1.1....				DOWN ROTATE
232	L3350:	...1.1.1..				C → STACK
233	L3351:	...1.1.1..				C → STACK
234	L3352:	...1.1.1..				C → STACK
235	L3353:	1.11.11.11	→L3266			GO TO FVR9
236	L3354:	1.1.1.1....		TEN6	:	M → C
237	L3355:	1111.1.1.				C + 1 → C[X]
238	L3356:	1111.1.1.				C + 1 → C[X]
239	L3357:	1111.1.1.				A + 1 → A[X]
240	L3360:	1111.1.1.				A + 1 → A[X]
241	L3361:	1111.1.1.				A + 1 → A[X]
242	L3362:	1111.1.1.				A + 1 → A[X]
243	L3363:11....				RETURN
244	L3364:	1....1....	→L4365	*****	SELR4	: SELECT ROM 4
245	L3365:	11...11.1	→L3343		N41	: JSB DIV
246	L3366:	...1.1.1.	→L3103			JSB R1000
247	L3367:	11.1.11.1	→L3153			JSB ADD
248	L3370:	11.1.1....				DOWN ROTATE
249	L3371:	11.1.1....				STACK → A
250	L3372:	11.1.1....				STACK → A
251	L3373:	1...1.111.				B EXCHANGE C[W]
252	L3374:	111.1.111.				A EXCHANGE C[W]
253	L3375:	1.1.1.1..				IF S10 # 1
254	L3376:	...111111	→L3017			THEN GO TO N44
255	L3377:	...1.1.11	→L3012			GO TO N46

ROM 4

0	L4000:1....	→L0001	*****	ERROR	:	SELECT ROM 0
1	L4001:					NO OPERATION
2	L4002:					NO OPERATION
3	L4003:	11.1.1....	→L3004	*****	XTY	:	SELECT ROM 3
4	L4004:11....			RETUR1	:	RETURN
5	L4005:111.1	→L4016		SOD2	:	JSB DOWN3
6	L4006:	111.1.1..					1 → S7
7	L4007:	1.1.1.1..					0 → S4
8	L4010:	1...11.1.	→L4106				JSB ONE
9	L4011:	11.1.11.1	→L4153				JSB ADD
10	L4012:	...1...111	→L4041				GO TO SOD3
11	L4013:	11.1.1....			STA1	:	STACK → A
12	L4014:	1.1.1.1..					C → STACK
13	L4015:11....					RETURN
14	L4016:	11.1.1....			DOWN3	:	DOWN ROTATE
15	L4017:	11.1.1....			DOWN2	:	DOWN ROTATE
16	L4020:	11.1.1....					DOWN ROTATE
17	L4021:11....					RETURN
18	L4022:	1.1.1.1..					IF S5 # 1
19	L4023:1.1.1	→L4004				THEN GO TO RETUR1
20	L4024:	1.1.1....	→L5025	*****			SELECT ROM 5
21	L4025:	111.1.1..			SOD	:	IF S7 # 1
22	L4026:1.111	→L4005				THEN GO TO SOD2
23	L4027:	...1111111	→L4037				GO TO SOD6
24	L4030:11	→L4000		FV	:	GO TO ERROR
25	L4031:					NO OPERATION
26	L4032:	1.1.1....	→L5033	*****	PV	:	SELECT ROM 5
27	L4033:	1111.11.11	→L4366		PMT	:	GO TO DNOTE1
28	L4034:	1.1.1.1..			R	:	1 → S5
29	L4035:	1.1.1....	→L5036	*****			SELECT ROM 5
30	L4036:11	→L4000		N	:	GO TO ERROR
31	L4037:	1...1.1..			SOD6	:	IF S4 # 1
32	L4040:	111.11111	→L4347				THEN GO TO SOD1
33	L4041:	11.1.1....			SOD3	:	DOWN ROTATE
34	L4042:	1.1.1.1..			SOD5	:	0 → S4
35	L4043:	11.1.1....					STACK → A
36	L4044:	11.1.1....					DOWN ROTATE
37	L4045:	11...111.					C → A[W]
38	L4046:	...1111.1	→L4017				JSB DOWN2
39	L4047:	11.1.1.1.	→L4152				JSB SUB
40	L4050:	...111.1.	→L4013				JSB STA1
41	L4051:	1...11.1.	→L4106				JSB ONE
42	L4052:	1.1.1.111	→L4132				GO TO SOD4
43	L4053:	...1.1.111	→L4025		DEPR	:	GO TO SOD
44	L4054:	111.1.1..			TRND1	:	IF S7 # 1
45	L4055:	...11.1.11	→L4062				THEN GO TO TRND5
46	L4056:	1...1.1..					IF S4 # 1
47	L4057:	11.11.111	→L4155				THEN GO TO TRND3
48	L4060:	1.1.1.1..					0 → S4
49	L4061:	1.11.111	→L4231				GO TO TRND8
50	L4062:	111...1..			TRND5	:	1 → S7
51	L4063:	1...1.1..					IF S4 # 1
52	L4064:	1...11111	→L4107				THEN GO TO TRND4
53	L4065:	1.1.1.1..					0 → S4
54	L4066:	11.1.1....					DOWN ROTATE
55	L4067:	1.1.1.111	→L4225				GO TO TRND2
56	L4070:	11.1....	→L3071	*****	R13	:	SELECT ROM 3
57	L4071:	111.1.111			L360	:	A EXCHANGE C[W]
58	L4072:	11.1.111.					0 → C[W]
59	L4073:	11.1.1....					LOAD CONSTANT 3
60	L4074:	11.1.1....					LOAD CONSTANT 6

ROM 4—Continued

61	L4075:	.1111.1.1.				C + 1 → C[X]
62	L4076:	.1111.1.1.				C + 1 → C[X]
63	L4077:11....				RETURN
64	L4100:	.11.1.1.1	→ L4152	DN0TE4 :		JSB SUB
65	L4101:11.1.1	→ L4013			JSB STA1
66	L4102:	.1.1.1....	→ L5103	*****		SELECT ROM 5
67	L4103:	.1.1.1....	→ L1104	*****	R100 :	SELECT ROM 1
68	L4104:	.11.1.1....			TRND6 :	DOWN ROTATE
69	L4105:	.111....11	→ L4070			GO TO R13
70	L4106:	.1.1.1....	→ L1107	*****	ONE :	SELECT ROM 1
71	L4107:	.111....1.			TRND4 :	1 → S7
72	L4110:11.1.1	→ L4013			JSB STA1
73	L4111:	.11.1.111.				A EXCHANGE C[W]
74	L4112:	.1....11.1	→ L4106			JSB ONE
75	L4113:	.11.1.11.1	→ L4153			JSB ADD
76	L4114:	.11.1.1....				DOWN ROTATE
77	L4115:	.1.1.1....				C → STACK
78	L4116:	.111....1.1	→ L4345			JSB MPY
79	L4117:	.11.1.1....				DOWN ROTATE
80	L4120:	.1....11.1.				B EXCHANGE C[W]
81	L4121:	.11.1.1....				DOWN ROTATE
82	L4122:	.11.1.111.		TRND9 :		A EXCHANGE B[W]
83	L4123:	.11.1.11.1	→ L4153			JSB ADD
84	L4124:	.11.1.1....				DOWN ROTATE
85	L4125:	.11.1.1....				STACK → A
86	L4126:	.11.1.11.1	→ L4153			JSB ADD
87	L4127:	.1.1.1....				C → STACK
88	L4130:1111.1	→ L4017			JSB DOWN2
89	L4131:	.111....11	→ L4070			GO TO R13
90	L4132:	.11.1.11.1	→ L4153	SOD4 :		JSB ADD
91	L4133:	.1.1.1....				M → C
92	L4134:	.11.1.11.1	→ L4345			JSB MPY
93	L4135:	.11.1.11.1	→ L4153			JSB ADD
94	L4136:11.1.1	→ L4013			JSB STA1
95	L4137:	.1.1.111.				B → C[W]
96	L4140:	.111....1.1	→ L4345			JSB MPY
97	L4141:11.1.1	→ L4013			JSB STA1
98	L4142:	.11.1.111.				A EXCHANGE C[W]
99	L4143:	.111....11	→ L4070			GO TO R13
100	L4144:	.1.1.1....		INTER :		C EXCHANGE M
101	L4145:	.1....11.1	→ L4103			JSB R100
102	L4146:	.1.1.1.1....				M → C
103	L4147:	.11.1.111.				A EXCHANGE C[W]
104	L4150:	.1.1.1....				C EXCHANGE M
105	L4151:	.1.1....1.1	→ L4242	*****	SUB :	GO TO INTER1
106	L4152:	.1.1.1....	→ L1153	*****	ADD :	SELECT ROM 1
107	L4153:	.1.1.1....	→ L1154	*****	ADD1 :	SELECT ROM 1
108	L4154:	.1.1.1....	→ L1155	*****	TRND3 :	SELECT ROM 1
109	L4155:11.1.1	→ L4013			JSB STA1
110	L4156:	.1.1.1....				C → STACK
111	L4157:	.111....1.1	→ L4343			JSB DIV
112	L4160:	.11.1.1....				DOWN ROTATE
113	L4161:	.1....11.1	→ L4106			JSB ONE
114	L4162:	.11.1.11.1	→ L4153			JSB ADD
115	L4163:	.1.1.111.				B → C[W]
116	L4164:	.11.1.1....				DOWN ROTATE
117	L4165:	.111....1.1	→ L4345			JSB MPY
118	L4166:	.11.1.1....				STACK → A
119	L4167:	.111....1.1	→ L4343			JSB DIV
120	L4170:	.11.1.11.1	→ L4153			JSB ADD
121	L4171:	.11.1.1....				DOWN ROTATE
122	L4172:	.1.1.1....				C → STACK
123	L4173:	.11.1.1.1	→ L4152			JSB SUB
124	L4174:	.11.1.11.1	→ L4153			JSB ADD
125	L4175:	.11.1.11.				A EXCHANGE B[W]
126	L4176:	.11.1.11.1	→ L4153			JSB ADD
127	L4177:111.1	→ L4016			JSB DOWN3
128	L4200:	.1....11.1	→ L4106			JSB ONE
129	L4201:	.11.1.1.1	→ L4152			JSB SUB
130	L4202:	.11.1.1....				STACK → A
131	L4203:	.111....1.1	→ L4343			JSB DIV
132	L4204:	.11.1.11.1	→ L4153			JSB ADD
133	L4205:	.11.1.1....				DOWN ROTATE
134	L4206:	.11.1.111.				A EXCHANGE B[W]
135	L4207:	.11.1.1.1	→ L4152			JSB SUB
136	L4210:	.11.1.1....				STACK → A
137	L4211:11.1.1	→ L4013			JSB STA1
138	L4212:	.1....111.				B → C[W]
139	L4213:	.111....1.1	→ L4345			JSB MPY
140	L4214:	.11.1.1....				STACK → A
141	L4215:	.11.1.11.1	→ L4153			JSB ADD
142	L4216:	.11111111.				0 - C - 1 → C[S]
143	L4217:	.1.1.1....				C EXCHANGE M
144	L4220:	.111....111.				0 → C[W]
145	L4221:	.1.1.1....				C → STACK
146	L4222:	.1.1.1.1....				M → C
147	L4223:	.1.1.1....				C → STACK
148	L4224:	.1....11.1	→ L4104			GO TO TRND6
149	L4225:	.1....11.1	→ L4106	TRND2 :		JSB ONE
150	L4226:	.11.1.11.1	→ L4153			JSB ADD
151	L4227:	.1.1.1....				C → STACK
152	L4230:	.1.1.1....				C → STACK
153	L4231:	.11.1.1....		TRND8 :		STACK → A
154	L4232:	.11.1.1....				STACK → A

ROM 4—Continued

155	L4233:	..11...111.				C → A[W]
156	L4234:1111.1	→ L4017			JSB DOWN2
157	L4235:	111...1.1.1	→ L4345			JSB MPY
158	L4236:	1.1.1.1...				M → C
159	L4237:	..11.1.1.1	→ L4153			JSB ADD
160	L4240:1.1.1.1	→ L4013			JSB STA1
161	L4241:	1...1.1.1	→ L4104			GO TO TRND6
162	L4242:	1.1.1.1...		INTER1 :		M → C
163	L4243:	111...1.1.1	→ L4343			JSB DIV
164	L4244:	11...1.1...				DOWN ROTATE
165	L4245:	..1.1.1...				STACK → A
166	L4246:	..1.1.1.1	→ L4152			JSB SUB
167	L4247:1.1.1.1	→ L4013			JSB STA1
168	L4250:	1...1.1.1				B → C[W]
169	L4251:	..1.1.1.1	→ L4152			JSB SUB
170	L4252:	..1.1.1...				C EXCHANGE M
171	L4253:	1...1.1.1	→ L4106			JSB ONE
172	L4254:	..1.1.1.1	→ L4153			JSB ADD
173	L4255:	1...1.1...				C → STACK
174	L4256:	..1.1.1...				C EXCHANGE M
175	L4257:1.1.1	→ L4003			JSB XTY
176	L4260:	..1.1.1...				C EXCHANGE M
177	L4261:	1...1.1.1	→ L4106			JSB ONE
178	L4262:	..1.1.1.1	→ L4152			JSB SUB
179	L4263:1.1.1.1	→ L4013			JSB STA1
180	L4264:	111...1.1.1	→ L4345			JSB MPY
181	L4265:	11...1.1...				DOWN ROTATE
182	L4266:	1...1.1.1	→ L4106			JSB ONE
183	L4267:	..1.1.1.1	→ L4153			JSB ADD
184	L4270:1.1.1.1	→ L4013			JSB STA1
185	L4271:	111.1.1.1.				A EXCHANGE C[W]
186	L4272:1.1.1	→ L4003			JSB XTY
187	L4273:	..1.1.1...				STACK → A
188	L4274:	..1.1.1...				C EXCHANGE M
189	L4275:	1...1.1.1	→ L4106			JSB ONE
190	L4276:	..1.1.1.1	→ L4152			JSB SUB
191	L4277:	1.1.1.1...				M → C
192	L4300:	111...1.1.1	→ L4345			JSB MPY
193	L4301:	..1.1.1...				C EXCHANGE M
194	L4302:	1...1.1.1	→ L4106			JSB ONE
195	L4303:	..1.1.1.1	→ L4152			JSB SUB
196	L4304:	11...1.1...				DOWN ROTATE
197	L4305:	..1111.1.1.				0 - C - 1 → C[S]
198	L4306:	1...1.1...				C → STACK
199	L4307:	111...1.1.1	→ L4345			JSB MPY
200	L4310:	11...1.1...				DOWN ROTATE
201	L4311:	..1.1.1...				STACK → A
202	L4312:	..1.1.1...				STACK → A
203	L4313:	..1.1.1...				C EXCHANGE M
204	L4314:	..1.1.1.1	→ L4152			JSB SUB
205	L4315:	1.1.1.1...				M → C
206	L4316:	111...1.1.1	→ L4345			JSB MPY
207	L4317:	..1.1.1.1	→ L4070			GO TO R13
208	L4320:	111...1.1.1	→ L4345	DNOTE3 :		JSB MPY
209	L4321:	..1.1.1...				C EXCHANGE M
210	L4322:	..1.1.1...				STACK → A
211	L4323:1.1.1.1	→ L4013			JSB STA1
212	L4324:	..1.1.1.1	→ L4071			JSB L360
213	L4325:	11...1.1...				12 → P
214	L4326:	111...1.1.1	→ L4343			JSB DIV
215	L4327:	11...1.1...				DOWN ROTATE
216	L4330:	..1.1.1.1	→ L4071			JSB L360
217	L4331:	..1.1.1...				LOAD CONSTANT 5
218	L4332:	11...1.1...				12 → P
219	L4333:	111...1.1.1	→ L4343			JSB DIV
220	L4334:1.1.1.1	→ L4013			JSB STA1
221	L4335:	..1.1.1.1	→ L4152			JSB SUB
222	L4336:1111.1	→ L4017			JSB DOWN2
223	L4337:	..1.1.1...				STACK → A
224	L4340:	111.1.1.1.				A EXCHANGE C[W]
225	L4341:	1...1.1...				C → STACK
226	L4342:	1...1.1	→ L4100			GO TO DNOTE4
227	L4343:	1...1...	→ L1344	***** DIV :		SELECT ROM 1
228	L4344:				NO OPERATION
229	L4345:	1...1...	→ L1346	***** MPY :		SELECT ROM 1
130	L4346:	1...1...	→ L1347	***** TEN6 :		SELECT ROM 1
231	L4347:	1...1.1...		SOD1 :		0 → S4
232	L4350:	..1.1.1...				C EXCHANGE M
233	L4351:	11...1.1...				DOWN ROTATE
234	L4352:	1...1.1...				C → STACK
235	L4353:	1...1.1...				C → STACK
236	L4354:	1...1.1.1	→ L4106			JSB ONE
237	L4355:	..1.1.1.1	→ L4153			JSB ADD
238	L4356:	1...1.1.1				B EXCHANGE C[W]
239	L4357:	111...1.1.1	→ L4345			JSB MPY
240	L4360:	..1.1.1...				C EXCHANGE M
241	L4361:	111.1.1.1.				A EXCHANGE C[W]
242	L4362:	111...1.1.1	→ L4343			JSB DIV
243	L4363:	..1.1.1...				C EXCHANGE M
244	L4364:	..1.1.1.1	→ L4042			GO TO SOD5
245	L4365:	..1.1.1...		SEL4 :		KEYS → ROM ADDRESS
246	L4366:	1...1.1...		DNOTE1 :		0 → S4
247	L4367:	..1.1.1...				C EXCHANGE M
248	L4370:	1...1.1.1	→ L4103			JSB R100

ROM 4—Continued

249	L4371:	111.1.111.			A EXCHANGE C[W]
250	L4372:	1.1.1.1...			C → STACK
251	L4373:	111.1.1.1	→ L4343		JSB DIV
252	L4374:	1.1.1.1...			STACK → A
253	L4375:	11.1.1.1...			DOWN ROTATE
254	L4376:	1.1.1.1...			C EXCHANGE M
255	L4377:	11.1.1.11	→ L4320		GO TO DNOTE

ROM 5

0	L5000:			NO OPERATION
1	L5001:			NO OPERATION
2	L5002:	1.1.1.1...	→ L4003	*****	XTY : SELECT ROM 4
3	L5003:	11.1.1.1...			S182 : DOWN ROTATE
4	L5004:	1.1.1.1...			C → STACK
5	L5005:	111.1.111.			A EXCHANGE C[W]
6	L5006:	1.1.1.1...			C → STACK
7	L5007:	11.1.1.1...			0 → C[W]
8	L5010:	1.1.1.1...			LOAD CONSTANT 1
9	L5011:	1.1.1.1...			LOAD CONSTANT 8
10	L5012:	1.1.1.1...			LOAD CONSTANT 2
11	L5013:	1.1.1.1...			LOAD CONSTANT 5
12	L5014:	1.1.1.1...			LOAD CONSTANT 0
13	L5015:	1111.1.1.			S185 : C + 1 → C[X]
14	L5016:	1111.1.1.			C + 1 → C[X]
15	L5017:	11.1.1.1...			12 → P
16	L5020:	1.1.1.1...			RETR5 : RETURN
17	L5021:	11.1.1.1...			S180 : 0 → C[W]
18	L5022:	1.1.1.1...			LOAD CONSTANT 1
19	L5023:	1.1.1.1...			LOAD CONSTANT 8
20	L5024:	1.1.1.1...	→ L5015		GO TO S185
21	L5025:	1.1.1.1...			IF S10 # 1
22	L5026:	1.1.1.1...	→ L5030		THEN GO TO SEL6
23	L5027:	1.1.1.1...			RETURN
24	L5030:	11.1.1.1...	→ L6031	*****	SEL6 : SELECT ROM 6
25	L5031:	1.1.1.1...			NO OPERATION
26	L5032:	1.1.1.1...			NO OPERATION
27	L5033:	1.1.1.1...			BOND1 : 1 → S5
28	L5034:	1.1.1.1...	→ L5105		JSB ONE
29	L5035:	11.1.1.111	→ L5155		GO TO BOND3
30	L5036:	1.1111.1	→ L5076		BONDR1 : JSB STA1
31	L5037:	1.1.1.1...			1 → S11
32	L5040:	1.1.1.111.			A → B[W]
33	L5041:	11.1.1.1...			C → A[W]
34	L5042:	1.1.1.1...	→ L5154		JSB ADD1
35	L5043:	1.1.1.111.			A EXCHANGE B[W]
36	L5044:	11.1.1.1...	→ L5343		JSB DIV
37	L5045:	1.1.1.1...			C EXCHANGE M
38	L5046:	1.1.1.1...	→ L5102		JSB R100
39	L5047:	111.1.111.			A EXCHANGE C[W]
40	L5050:	11.1.1.1...	→ L5343		JSB DIV
41	L5051:	1.1.1.1...	→ L5003		JSB S182
42	L5052:	11.1.1.1...	→ L5343		JSB DIV
43	L5053:	1.1111.1			IF C[X5] >= 1
44	L5054:	1.1111.1	→ L5134		THEN GO TO BONDR2
45	L5055:	1.1111.1	→ L5072		JSB DOWN2
46	L5056:	1.1.1.1...			M → C
47	L5057:	11.1.1.1...			BON2 : DOWN ROTATE
48	L5060:	1.1.1.1...			BONDR3 : M → C
49	L5061:	1.1.1.1...	→ L5105		JSB ONE
50	L5062:	11.1.1.1...	→ L5153		JSB ADD
51	L5063:	1.1111.1	→ L5076		JSB STA1
52	L5064:	111.1.111.			A EXCHANGE C[W]
53	L5065:	1.1.1.111	→ L5261		GO TO BONDR7
54	L5066:	1.1.1.1...			STACK → A
55	L5067:	1.1.1.1...			0 → S5
56	L5070:	1.1.1.1...	→ L3071	*****	SELECT ROM 3
57	L5071:	11.1.1.1...			DOWN3 : DOWN ROTATE
58	L5072:	11.1.1.1...			DOWN2 : DOWN ROTATE
59	L5073:	11.1.1.1...			DOWN ROTATE
60	L5074:	1.1.1.1...			RETURN
61	L5075:	11.1.1.1...			STA2 : DOWN ROTATE
62	L5076:	1.1.1.1...			STA1 : STACK → A
63	L5077:	1.1.1.1...			C → STACK
64	L5100:	1.1.1.1...			RETURN
65	L5101:	1.1.1.1...			NO OPERATION
66	L5102:	1.1.1.1...	→ L4103	*****	R100 : SELECT ROM 4
67	L5103:	1.1.1.1...			DNOTE5 : 1 → S5
68	L5104:	1.1.1.11	→ L5106		GO TO DNOTE6
69	L5105:	1.1.1.1...	→ L4106	*****	ONE : SELECT ROM 4
70	L5106:	1.1.1.1...			DNOTE6 : M → C
71	L5107:	111.1.1.1	→ L5345		JSB MPY
72	L5110:	1.1.1.1...	→ L5102		JSB R100
73	L5111:	111.1.111.			A EXCHANGE C[W]
74	L5112:	11.1.1.1...	→ L5343		JSB DIV
75	L5113:	1.1111.1	→ L5075		JSB STA2
76	L5114:	1.1.1.1...			M → C
77	L5115:	11.1.1.1...	→ L5345		JSB MPY
78	L5116:	1.1.1.1...	→ L5102		JSB R100
79	L5117:	111.1.111.			A EXCHANGE C[W]
80	L5120:	11.1.1.1...	→ L5343		JSB DIV
81	L5121:	1.1.1.1...			DOWN ROTATE
82	L5122:	1.1.1.1...			IF S11 # 1

ROM 5—Continued

83	L5123:	..11.11.11	→ L5066			THEN GO TO DNOTE2
84	L5124:	..11.11.11	→ L5067			GO TO R13
85	L5125:	11.1.1.1		IT1 :		DOWN ROTATE
86	L5126:	11.1.1.1				C → A[W]
87	L5127:	1.111.1.1		IT2 :		IF A[X] ≥ 1
88	L5130:	...1.1.1	→ L5020			THEN GO TO RETRS
89	L5131:	11.11.1.1				A - 1 → A[X]
90	L5132:	1.1.1.1				SHIFT LEFT A[M]
91	L5133:	1.1.11.11	→ L5127			GO TO IT2
92	L5134:	11.1.1.1		BONDR2 :		STACK → A
93	L5135:	1.1.1.1				M → C
94	L5136:	1.1.1.1.1	→ L5153			JSB ADD
95	L5137:	1.111.1.1	→ L5076			JSB STA1
96	L5140:	...1.1.1	→ L5021			JSB S180
97	L5141:	111.1.1.1	→ L5343			JSB DIV
98	L5142:	1.1.1.1.1	→ L5105			JSB ONE
99	L5143:	1.1.1.1.1	→ L5152			JSB SUB
100	L5144:	1.1.1.1.1				B EXCHANGE C[W]
101	L5145:	1.1.1.1				C EXCHANGE M
102	L5146:	1.111.1.1				0 - C - 1 → C[S]
103	L5147:	111.1.1.1	→ L5345			JSB MPY
104	L5150:	1.1.1.1.1	→ L5105			JSB ONE
105	L5151:	11.1.1.11	→ L5331			GO TO BONDR8
106	L5152:	1.1.1.1	→ L1153 *****	SUB :		SELECT ROM 1
107	L5153:	1.1.1.1	→ L1154 *****	ADD :		SELECT ROM 1
108	L5154:	1.1.1.1	→ L1155 *****	ADD1 :		SELECT ROM 1
109	L5155:	1.1.1.11.1		BOND 3 :		C + C → C[W]
110	L5156:	111.1.1.1	→ L5343			JSB DIV
111	L5157:	1.1.1.1				C EXCHANGE M
112	L5160:	1.1.1.1	→ L5102			JSB R100
113	L5161:	1.1.1.1	→ L5003			JSB S182
114	L5162:	111.1.1.1	→ L5343			JSB DIV
115	L5163:	1.111.1.1				IF C[X] ≥ 1
116	L5164:	1.1.1.1.1	→ L5232			THEN GO TO BOND2
117	L5165:	1.111.1.1	→ L5075			JSB STA2
118	L5166:	1.1.1.1.1	→ L5105			JSB ONE
119	L5167:	1.1.1.1.1				C + C → C[W]
120	L5170:	111.1.1.1				A EXCHANGE C[W]
121	L5171:	1.1.1.1.1	→ L5153			JSB ADD
122	L5172:	1.1.1.1				B → C[W]
123	L5173:	111.1.1.1	→ L5343			JSB DIV
124	L5174:	1.1.1.1				DOWN ROTATE
125	L5175:	1.1.1.1				STACK → A
126	L5176:	1.1.1.1				DOWN ROTATE
127	L5177:	1.111.1.1				0 - C - 1 → C[S]
128	L5200:	1.1.1.1	→ L5002			JSB XTY
129	L5201:	1.1.1.1.1	→ L5125			JSB IT1
130	L5202:	111.1.1.1				A EXCHANGE C[W]
131	L5203:	1.1.1.1.1	→ L5105			JSB ONE
132	L5204:	1.1.1.1.1	→ L5153			JSB ADD
133	L5205:	1.1.1.1	→ L5002			JSB XTY
134	L5206:	1.11.1.1.1	→ L5071			JSB DOWN3
135	L5207:	1.111.1.1	→ L5076			JSB STA1
136	L5210:	1.1.1.1.1	→ L5152			JSB SUB
137	L5211:	1.1.1.1.1	→ L5153			JSB ADD
138	L5212:	1.1.1.1				M → C
139	L5213:	11.1.1.1.1	→ L5345			JSB MPY
140	L5214:	1.1.1.1				DOWN ROTATE
141	L5215:	1.1.1.1	→ L5015			JSB S185
142	L5216:	1.1.1.1				DOWN ROTATE
143	L5217:	1.1.1.1.1				C → A[W]
144	L5220:	1.1.1.1				M → C
145	L5221:	111.1.1.1	→ L5345			JSB MPY
146	L5222:	1.1.1.1				DOWN ROTATE
147	L5223:	1.1.1.1				STACK → A
148	L5224:	111.1.1.1	→ L5343			JSB DIV
149	L5225:	1.1.1.1				STACK → A
150	L5226:	1.1.1.1.1	→ L5153			JSB ADD
151	L5227:	1.1.1.1				DOWN ROTATE
152	L5230:	1.1.1.1.1	→ L5152			JSB SUB
153	L5231:	1.1.1.1.1	→ L5067			GO TO R13
154	L5232:	1.111.1.1	→ L5075	BOND2 :		JSB STA2
155	L5233:	1.1.1.1	→ L5021			JSB S180
156	L5234:	111.1.1.1	→ L5343			JSB DIV
157	L5235:	1.1.1.1				DOWN ROTATE
158	L5236:	1.1.1.1.1	→ L5345			JSB MPY
159	L5237:	1.1.1.1.1	→ L5105			JSB ONE
160	L5240:	1.1.1.1.1				C + C → C[W]
161	L5241:	1.1.1.1.1	→ L5153			JSB ADD
162	L5242:	1.1.1.1				C EXCHANGE M
163	L5243:	1.1.1.1.1	→ L5105			JSB ONE
164	L5244:	1.1.1.1	→ L5015			JSB S185
165	L5245:	1.1.1.1.1	→ L5153			JSB ADD
166	L5246:	1.1.1.1				B → C[W]
167	L5247:	1.1.1.1				C EXCHANGE M
168	L5250:	111.1.1.1	→ L5343			JSB DIV
169	L5251:	1.1.1.1.1	→ L5153			JSB ADD
170	L5252:	1.1.1.1.1	→ L5071			JSB DOWN3
171	L5253:	1.1.1.1.1	→ L5105			JSB ONE
172	L5254:	1.1.1.1.1	→ L5152			JSB SUB
173	L5255:	1.1.1.1				C EXCHANGE M
174	L5256:	1.1.1.1.1	→ L5345			JSB MPY
175	L5257:	1.1.1.1				DOWN ROTATE
176	L5260:	1.111.1.1	→ L5376			GO TO BONDR6

ROM 5—Continued

177	L5261:1..1	→ L5002	BONDR7 :	JSB XTY
178	L5262:	..1...1..1	→ L5105		JSB ONE
179	L5263:	..11..1..1	→ L5152		JSB SUB
180	L5264:	..11..1...1	→ L5330		JSB ROT1
181	L5265:	...111..1..1	→ L5071		JSB DOWN3
182	L5266:	..11..1..1	→ L5152		JSB SUB
183	L5267:	..11..1...1	→ L5330		JSB ROT1
184	L5270:	..11..1...			DOWN ROTATE
185	L5271:	..111..111.			A EXCHANGE C[W]
186	L5272:	..111...11.1	→ L5343		JSB DIV
187	L5273:	..11..1...			M → C
188	L5274:	..111..1..1	→ L5345		JSB MPY
189	L5275:	...11111..1	→ L5076		JSB STA1
190	L5276:	..11..111.1	→ L5153		JSB ADD
191	L5277:	..11..1...			STACK → A
192	L5300:	...111..1..1	→ L5072		JSB DOWN2
193	L5301:	...1...111.			B → C[W]
194	L5302:	...1111..1..1	→ L5075		JSB STA2
195	L5303:	..11..1...			M → C
196	L5304:	..11..111.1	→ L5153		JSB ADD
197	L5305:	...11..1...			C EXCHANGE M
198	L5306:	...1...111.			B → C[W]
199	L5307:	...11..1..1	→ L5015		JSB S185
200	L5310:	...11..1..1	→ L5015		JSB S185
201	L5311:	...11..1..1	→ L5015		JSB S185
202	L5312:	..11..1..1.			IF C[M] = 0
203	L5313:	..11..111.1	→ L5316		THEN GO TO BON1
204	L5314:	..11..111.1			IF C[X5] = 0
205	L5315:	...11...11.	→ L5060		THEN GO TO BONDR3
206	L5316:	..11..1..1.		BON1 :	IF S11 # 1
207	L5317:	..1111..1111	→ L5373		THEN GO TO BONDR4
208	L5320:	..111..1..1			0 → S11
209	L5321:	..11..1...			DOWN ROTATE
210	L5322:	..11..1...			C → STACK
211	L5323:	..11..111.1	→ L5125		JSB IT1
212	L5324:	..111..111.			A EXCHANGE C[W]
213	L5325:	...1...1..1	→ L5105		JSB ONE
214	L5326:	..11..1..1	→ L5152		JSB SUB
215	L5327:	..111..11.1	→ L5346		GO TO BONDR9
216	L5330:	...1...1...	→ L1331	*****	SELECT ROM 1
217	L5331:	..11..111.1	→ L5153	ROT1 :	JSB ADD
218	L5332:	..11..1...		BONDR8 :	STACK → A
219	L5333:	..111...11.1	→ L5343		JSB DIV
220	L5334:	...1...1..1	→ L5105		JSB ONE
221	L5335:	..11..1..1	→ L5152		JSB SUB
222	L5336:	..11..1...			M → C
223	L5337:	..111...11.1	→ L5343		JSB DIV
224	L5340:	..11111..11	→ L5374		GO TO BONDR5
225	L5341:			NO OPERATION
226	L5342:			NO OPERATION
227	L5343:	...1...1...	→ L1344	*****	SELECT ROM 1
228	L5344:		DIV :	NO OPERATION
229	L5345:	...1...1...	→ L1346	*****	SELECT ROM 1
230	L5346:	...1...111.		MPY :	B → C[W]
231	L5347:	..111..1..1	→ L5345	BONDR9 :	JSB MPY
232	L5350:	..11..1...			M → C
233	L5351:	..111..1..1	→ L5345		JSB MPY
234	L5352:	...111..1..1	→ L5072		JSB DOWN2
235	L5353:	...11...111.			C → A[W]
236	L5354:	...111..1..1	→ L5072		JSB DOWN2
237	L5355:	..111..1..1	→ L5345		JSB MPY
238	L5356:	...1...1..1	→ L5105		JSB ONE
239	L5357:	..111..111.			C + C → C[W]
240	L5360:	..111..111.			A EXCHANGE C[W]
241	L5361:	..11..111.1	→ L5153		JSB ADD
242	L5362:	...1...111.			B → C[W]
243	L5363:	..111...11.1	→ L5343		JSB DIV
244	L5364:	...11111..1	→ L5076		JSB STA1
245	L5365:	..111..1..1	→ L5345		JSB MPY
246	L5366:	..11..1...			DOWN ROTATE
247	L5367:	..11..1...			STACK → A
248	L5370:	..111..1..1	→ L5345		JSB MPY
249	L5371:	...1...1...			C → STACK
250	L5372:	...11111111	→ L5057		GO TO BON2
251	L5373:	..11..1...		BONDR4 :	M → C
252	L5374:	...11..1..1	→ L5015	BONDR5 :	JSB S185
253	L5375:	..11...111.			C → A[W]
254	L5376:	..11..111.1	→ L5153	BONDR6 :	JSB ADD
255	L5377:	...11111111	→ L5067		GO TO R13

ROM 6

0	L6000:1...	→ L0001	*****	ERR71 :	SELECT ROM 0
1	L6001:	..1.....1..			DA8 :	1 → S4
2	L6002:	..111..11.11	→ L6354			GO TO DA12
3	L6003:	..111..11.			DM6 :	IF P # 9
4	L6004:	...111..11	→ L6214			THEN GO TO DM7
5	L6005:	...11....				RETURN
6	L6006:	..11..1111.			DM3 :	A EXCHANGE B[W]
7	L6007:	...11..1...			DM1 :	IF P # 2
8	L6010:	...111..111	→ L6211			THEN GO TO DM4
9	L6011:	..1111..11.				A + 1 → A[M]

ROM 6—Continued

10	L6012:	.111..11.			C - 1 → C[M]
11	L6013:	.11.1.1.1.			IF C[X]=0
12	L6014:111111	→ L6017		THEN GO TO DM2
13	L6015:	11111..11.			A + 1 → A[M]
14	L6016:	.1.1.1.1.			C - 1 → C[M]
15	L6017:11....		DM2	RETURN
16	L6020:11.11...		YC1	LOAD CONSTANT 3
17	L6021:11.11...			LOAD CONSTANT 6
18	L6022:11.11...			LOAD CONSTANT 5
19	L6023:1.1.1.1.			IF S4 # 1
20	L6024:1.111111	→ L6027		THEN GO TO YC2
21	L6025:	11111.1.1.			A + 1 → A[X]
22	L6026:11....			RETURN
23	L6027:1.1.11...		YC2	LOAD CONSTANT 2
24	L6030:1.1.11...			LOAD CONSTANT 5
25	L6031:11....			RETURN
26	L6032:			NO OPERATION
27	L6033:			NO OPERATION
28	L6034:1.1.1	→ L6006	DD6	JSB DM3
29	L6035:	11.1.1.111.			A EXCHANGE B[W]
30	L6036:	.111..11.			A + C → C[M]
31	L6037:11.1.		DD8	P - 1 → P
32	L6040:1.1.1.			IF P # 0
33	L6041:111..11	→ L6034		THEN GO TO DD6
34	L6042:11.1.1.			0 → C[X]
35	L6043:	.11....111.			C → A[W]
36	L6044:	.111.1.1.			IF S7 # 1
37	L6045:	1.1.1.1.11	→ L6222		THEN GO TO DA4
38	L6046:	.11.1.1....		DN1	STACK → A
39	L6047:	111.1.111.			A EXCHANGE C[W]
40	L6050:	1.111..11	→ L6234		GO TO DN3
41	L6051:	.111.1.1.			IF S7 # 1
42	L6052:	.1.1.1.111	→ L6055		THEN GO TO DA1
43	L6053:	.1..1.1....			IF S4 # 1
44	L6054:	.11....11	→ L6060		THEN GO TO DA3
45	L6055:	.1111..1..		DA1	0 → S7
46	L6056:	.1.1.1....			C → STACK
47	L6057:	.1.1.1....			C → STACK
48	L6060:	.11.1.1....		DA3	STACK → A
49	L6061:	.111.1.1.			IF S7 # 1
50	L6062:	.11.1.1.11	→ L6065		THEN GO TO DA13
51	L6063:	.1.1.1....			C → STACK
52	L6064:	111.1.111.		DA13	A EXCHANGE C[W]
53	L6065:	111.1.1111	→ L6353		GO TO DA5
54	L6066:	.11.1.1....		DM5	IF P # 6
55	L6067:1111	→ L6003		THEN GO TO DM6
56	L6070:11....			RETURN
57	L6071:			NO OPERATION
58	L6072:	.1.1.1.11.		DA7	A - C → C[MS]
59	L6073:	.1111.111	→ L6075		IF NO CARRY GO TO DA11
60	L6074:	.1.1.1.1.			0 - C → C[MS]
61	L6075:	1..1.1.11.		DA11	B EXCHANGE C[W]
62	L6076:	.11.1.1....			STACK → A
63	L6077:	11.1.1....			DOWN ROTATE
64	L6100:	.11111111.			0 - C - 1 → C[S]
65	L6101:	.11.1.1.1	→ L6151		JSB ADD63
66	L6102:	1.1111111.			0 → A[S]
67	L6103:	.11.1.11.			0 → C[W]
68	L6104:	.1.1.11..			5 → P
69	L6105:1.1.1.			LOAD CONSTANT 1
70	L6106:	1..11....			LOAD CONSTANT 8
71	L6107:	111.1.111.			A EXCHANGE C[W]
72	L6110:	1..1.1.1.			IF A >= B[MS]
73	L6111:	.11.1.111	→ L6145		THEN GO TO DA10
74	L6112:	.1.1.1....		DA9	C → STACK
75	L6113:	.1..1.11.			B → C[W]
76	L6114:	.1.1.1.1	→ L6150	ADD62	JSB ADD61
77	L6115:	.1.1.1.1.			0 → S5
78	L6116:	.1111..1..			0 → S7
79	L6117:1....	→ L0120	*****	SELECT ROM 0
80	L6120:1111..		DD4	P + 1 → P
81	L6121:	.1.1.1.1.1.			C - 1 → C[X]
82	L6122:	1....1.111	→ L6205		IF NO CARRY GO TO DD3
83	L6123:	111.1.111.			A EXCHANGE C[W]
84	L6124:11.			IF B[M] = 0
85	L6125:11.	→ L6000		THEN GO TO ERR71
86	L6126:11.1.	→ L6006		JSB DM3
87	L6127:	.111..11.			A + C → C[M]
88	L6130:	11.1.1.11.			A EXCHANGE B[W]
89	L6131:	.1..1.1....			IF S4 # 1
90	L6132:	.1.111..11	→ L6134		THEN GO TO DD5
91	L6133:	.11111111	→ L6037		GO TO DD8
92	L6134:	1....11.		DD5	IF A >= B[M]
93	L6135:	.11111111	→ L6037		THEN GO TO DD8
94	L6136:11.	→ L6000		GO TO ERR71
95	L6137:	1.111.111.		DY1	0 → A[W]
96	L6140:	1..1.1.1.			B EXCHANGE C[WP]
97	L6141:1	→ L6020		JSB YC1
98	L6142:	1....11..			8 → P
99	L6143:	1..11.1.			B EXCHANGE C[WP]
100	L6144:	.111....11	→ L6160		GO TO MU3
101	L6145:	.1..1.11.		DA10	B → C[W]
102	L6146:	.11.1....1	→ L6150		JSB ADD61

ROM 6—Continued

103	L6147:	1.1.1.11	→ L6112		GO TO DA9
104	L6150:	1.111.111.		ADD61 :	0 → A[W]
105	L6151:	1.1.1.1.		ADD63 :	1 → S5
106	L6152:	1.11.1.1.			1 → S7
107	L6153:	1.1.1.1.			12 → P
108	L6154:	1.1.1.1.	→ L1155	*****	SELECT ROM 1
109	L6155:	1.11.1.11.		MU1 :	A + B → A[W]
110	L6156:	1.1.1.1.		MU2 :	C - 1 → C[P]
111	L6157:	1.1.1.1.11	→ L6155		IF NO CARRY GO TO MU1
112	L6160:	1.1.1.11.		MU3 :	SHIFT RIGHT B[W]
113	L6161:	1.1.1.1.			P - 1 → P
114	L6162:	1.1.1.1.			IF P # 3
115	L6163:	1.1.1.1.11	→ L6156		THEN GO TO MU2
116	L6164:	1.1.1.1.11.			A - 1 → A[W]
117	L6165:	1.1.1.1.111	→ L6167		IF NO CARRY GO TO DY2
118	L6166:	1.111.1.11.			A + 1 → A[W]
119	L6167:	1.111.1.1.		DY2 :	A + 1 → A[X]
120	L6170:	1.1.1.1.		DD1 :	P + 1 → P
121	L6171:	1.1.1.1.			SHIFT RIGHT C[W]
122	L6172:	1.1.1.1.			IF P # 9
123	L6173:	1.111.1.11	→ L6170		THEN GO TO DD1
124	L6174:	1.1.1.1.			LOAD CONSTANT 3
125	L6175:	1.1.1.1.			4 → P
126	L6176:	1.1.1.1.			B EXCHANGE C[W]
127	L6177:	1.1.1.1.		DD2 :	SHIFT RIGHT C[W]
128	L6200:	1.1.1.1.			P - 1 → P
129	6201:	1.111.1.1.			IF P # 14
130	L6202:	1.111.1.1.11	→ L6177		THEN GO TO DD2
131	L6203:	1.1.1.1.			IF C[X] = 0
132	L6204:	1.1.1.1.	→ L6000		THEN GO TO ERR71
133	L6205:	1.1.1.1.		DD3 :	IF P # 12
134	L6206:	1.1.1.1.	→ L6120		THEN GO TO DD4
135	L6207:	1.1.1.1.	→ L6000		GO TO ERR71
136	L6210:	1.1.1.1.			NO OPERATION
137	L6211:	1.1.1.1.		DM4 :	IF P # 4
138	L6212:	1.1.1.1.11	→ L6066		THEN GO TO DM5
139	L6213:	1.1.1.1.			RETURN
140	L6214:	1.1.1.1.		DM7 :	IF P # 11
141	L6215:	1.1.1.1.11	→ L6217		THEN GO TO DM8
142	L6216:	1.1.1.1.			RETURN
143	L6217:	1.1.1.1.		DM8 :	A - 1 → A[M]
144	L6220:	1.111.1.1.			C + 1 → C[M]
145	L6221:	1.1.1.1.			RETURN
146	L6222:	1.1.1.1.		DA4 :	C - 1 → C[P]
147	L6223:	1.1.1.1.			DOWN ROTATE
148	L6224:	1.1.1.1.			IF S9 # 1
149	L6225:	1.1.1.1.11	→ L6351		THEN GO TO DA6
150	L6226:	1.1.1.1.			IF S4 # 1
151	L6227:	1.1.1.1.	→ L6001		THEN GO TO DA8
152	L6230:	1.1.1.1.			0 → S9
153	L6231:	1.1.1.1.11	→ L6354		GO TO DA12
154	L6232:	1.1.1.1.		DN2 :	0 → P
155	L6233:	1.1.1.1.			SHIFT RIGHT C[M]
156	L6234:	1.1.1.1.		DN3 :	C + 1 → C[P]
157	L6235:	1.1.1.1.11	→ L6232		IF NO CARRY GO TO DN2
158	L6236:	1.1.1.1.			IF C[X] >= 1
159	L6237:	1.1.1.1.	→ L6000		THEN GO TO ERR71
160	L6240:	1.1.1.1.11.			IF C[S] = 0
161	L6241:	1.1.1.1.11	→ L6243		THEN GO TO DN4
162	L6242:	1.1.1.1.			0 - C → C[M]
163	L6243:	1.111.1.11.		DN4 :	A + C → A[MS]
164	L6244:	1.1.1.1.			7 → P
165	L6245:	1.1.1.1.			0 → C[W]
166	L6246:	1.1.1.1.			LOAD CONSTANT 7
167	L6247:	1.1.1.1.			LOAD CONSTANT 3
168	L6250:	1.1.1.1.			LOAD CONSTANT 0
169	L6251:	1.1.1.1.			LOAD CONSTANT 5
170	L6252:	1.1.1.1.			IF A >= C[M]
171	L6253:	1.1.1.1.	→ L6000		THEN GO TO ERR71
172	L6254:	1.1.1.1.			8 → P
173	L6255:	1.1.1.1.	→ L6020		JSB YC1
174	L6256:	1.1.1.1.			0 → B[W]
175	L6257:	1.1.1.1.11.			B EXCHANGE C[W]
176	L6260:	1.1.1.1.			IF A[M] >= 1
177	L6261:	1.1.1.1.11	→ L6263		THEN GO TO DN11
178	L6262:	1.1.1.1.	→ L6000		GO TO ERR71
179	L6263:	1.111.1.11.		DN11 :	A + 1 → A[M]
180	L6264:	1.1.1.1.		DN15 :	SHIFT RIGHT B[W]
181	L6265:	1.1.1.1.			P - 1 → P
182	L6266:	1.1.1.1.11	→ L6270		GO TO DN6
183	L6267:	1.1.1.1.		DN5 :	C + 1 → C[P]
184	L6270:	1.1.1.1.		DN6 :	A - B → A[W]
185	L6271:	1.1.1.1.11	→ L6267		IF NO CARRY GO TO DN5
186	L6272:	1.1.1.1.			A + B → A[W]
187	L6273:	1.1.1.1.			IF P # 0
188	L6274:	1.1.1.1.11	→ L6264		THEN GO TO DN15
189	L6275:	1.1.1.1.			IF A[M] >= 1
190	L6276:	1.1.1.1.	→ L6301		THEN GO TO DN12
191	L6277:	1.1.1.1.			A + B → A[W]
192	L6300:	1.1.1.1.			C - 1 → C[X]
193	L6301:	1.1.1.1.		DN12 :	IF C[X] >= 1
194	L6302:	1.1.1.1.	→ L6304		THEN GO TO DN7
195	L6303:	1.1.1.1.			A - 1 → A[W]
196	L6304:	1.1.1.1.		DN7 :	4 → P

ROM 6—Continued

197	L6305:	..11.11..			LOAD CONSTANT 3
198	L6306:11..			0 → P
199	L6307:	1...1.111.			B EXCHANGE C[W]
200	L6310:	..11..111.			0 → C[W]
201	L6311:	111.1.1.1.			A EXCHANGE C[X]
202	L6312:1111..			P + 1 → P
203	L6313:	11111.1.1.		DN8	A + 1 → A[X]
204	L6314:	..11...11.			0 → C[M]
205	L6315:111.1	→ L6007		JSB DM1
206	L6316:	11.....11.			A - B → A[M]
207	L6317:	11.1...111	→ L6321		IF NO CARRY GO TO DN13
208	L6320:	11.1..1111	→ L6323		GO TO DN14
209	L6321:	1.1.1..11.		DN13	IF A[M] >= 1
210	L6322:	11.1.1.1.1	→ L6312		THEN GO TO DN8
211	L6323:	111....11.		DN14	A + B → A[M]
212	L6324:	1111...11.			A + C → A[M]
213	L6325:	..1....11.			SHIFT LEFT A[M]
214	L6326:	1111..11.			A + 1 → A[M]
215	L6327:	11.1.1.1.			A EXCHANGE B[X]
216	L6330:	..11..111.			0 → C[W]
217	L6331:	..11111.1.			C - 1 → C[XS]
218	L6332:	1111..111.			A + C → A[W]
219	L6333:	11.1.111.			A EXCHANGE B[W]
220	L6334:	11.1..11..			13 → P
221	L6335:111..		DN9	P - 1 → P
222	L6336:	..1...111.			SHIFT LEFT A[W]
223	L6337:	..1111.11..			IF P ≠ 7
224	L6340:	11.111.111.	→ L6335		THEN GO TO DN9
225	L6341:	111...111.			A + B → A[W]
226	L6342:1111..		DN10	P + 1 → P
227	L6343:	..1...111.			SHIFT LEFT A[W]
228	L6344:	11.1..11..			IF P ≠ 12
229	L6345:	111...1.11	→ L6342		THEN GO TO DN10
230	L6346:	11111.1.1.			A + 1 → A[X]
231	L6347:	111.1.111.			A EXCHANGE C[W]
232	L6350:	..1.11..11	→ L6114		GO TO ADD62
233	L6351:	..1...1.1.		DA6	IF S4 ≠ 1
234	L6352:	..111.1.11	→ L6072		THEN GO TO DA7
235	L6353:	..1..1..1.		DA5	0 → S4
236	L6354:	..1.11.1.1.		DA12	C - 1 → C[X]
237	L6355:	1111....11	→ L6360		IF NO CARRY GO TO DA2
238	L6356:	1.1..1.1.			SHIFT RIGHT C[M]
239	L6357:	..11..1.1.			0 → C[]
240	L6360:	..11.1.1.		DA2	IF C[X] >= 1
241	L6361:11	→ L6000		THEN GO TO ERR71
242	L6362:1.111.			0 → B[W]
243	L6363:	..11...111.			C → A[W]
244	L6364:	1....11..			8 → P
245	L6365:	..11.1.1.			0 → C[WP]
246	L6366:	..1..11..			LOAD CONSTANT 2
247	L6367:	...1.11..			LOAD CONSTANT 1
248	L6370:	1....11..			8 → P
249	L6371:	...1.1.1.			IF A >= C[WP]
250	L6372:11	→ L6000		THEN GO TO ERR71
251	L6373:	...1.11..			LOAD CONSTANT 1
252	L6374:	1.1.11..			LOAD CONSTANT 9
253	L6375:	1...11..			8 → P
254	L6376:	..1.1.1.1.			A - C → C[WP]
255	L6377:	..1.111111	→ L6137		IF NO CARRY GO TO DY1

FUNCTIONS

All of the functions performed by the calculator are given in the table below. The notes referred to in this table are given at the end of the table.

TABLE OF FUNCTIONS INCORPORATED INTO THE BUSINESS CALCULATOR

1. SIMPLE COMPOUNDING
 - 1.1 $FV = PV(1+i)^n$
 - 1.2 $PV = FV/(1+i)^n$
 - 1.3 $= (FV/PV)^{1/n} - 1$
 - 1.4 $n = (\text{Log}(FV))/\text{Log}(PV(1+i))$
2. ANNUITY
 - 2.1 $FV = PMT(1+i)^n - 1/i$
 - 2.2 $PV = PMT(1+i)^n - 1/i(1+i)^n$
 - 2.3 Solve For i In (see Note 1):
 $PV - PMT[(1+i)^n - 1]/i(1+i)^n = 0$
 - 2.4 Solve For i In (see Note 1):
 $FV - PMT[(1+i)^n - 1]/i = 0$
 - 2.5 $n = \log(PMT/(PV - PMT))/\log(1+i)$
- 2.6 $n = \log(FV \times i/PMT + 1)/\log(1+i)$
- 2.7 $PMT = PV i(1+i)^n/[(1+i)^n - 1]$
- 2.8 $PMT = FV i/[(1+i)^n - 1]$
3. ADD ON TO ANNUAL RATE
 - 3.1 Solve For i In (see Note 1):
4. ACCRUED INTEREST
 - n = No. of Months
 - R = Annual Add-On Rate
5. DISCOUNTED NOTE
 - 4.1 $i_{360} = n PV i/36000$
 - 4.2 $i_{365} = i_{360} \times 0.98630137$

$$1 - \frac{1 + \frac{n}{12} \times \frac{R}{100}}{n} \times \frac{(1+i)^n - 1}{i(1+i)^n}$$

n = No. of Days

i = Annual Interest Rate (%)

FV = Face Value of Note

$$5.1 \ d_{360} = FV \times n \times i / 36000$$

$$5.2 \ \text{yield}_{360} = d_{360} \times 36000 / n (FV - d_{360})$$

$$5.3 \ d_{365} = d_{360} \times 360 / 365$$

$$5.4 \ \text{yield}_{365} = d_{365} \times 36500 / n (FV - d_{365})$$

6. BOND

6.1 Price of a Bond (PV) (see Note 2):

n = No. of Days (uncompensated for leap days)

i = yield

c = coupon rate

For $n \geq 182.5$:

$$PV = 100 \left(1 + \frac{i}{200} \right)^{-\frac{n}{182.5}} + \frac{100c}{i} \left(1 + \frac{i}{200} \right)^j - \left(1 + \frac{i}{200} \right)^{-\frac{n}{182.5}} - \frac{cj}{2}$$

where $j = 1 - \text{frac } n / 182.5$

For $n < 182.5$:

$$PV = \frac{200 + c}{2 + \frac{n}{180} \cdot \frac{i}{100}} - \frac{(1-n)c}{2}$$

6.2 Yield of a Bond (see Note 2):

Solve for i knowing PV, in the above 2 equations, which one depending on n .

Solution gives

$$|i_{\text{actual}} - i_{\text{calc}}| < 2 \times i_{\text{actual}}^2 \times C \times 10^{-6}$$

7. DATE (see Note 3):

7.1 Date 1 - Date 2

7.2 Date $\pm n$ Days

1900 \leq Date \leq 2099 A.D.

8. ACCUMULATED INTEREST

accumulated interest (in \$)

$$= PMT \left\{ k - j \frac{\left(1 + \frac{i}{100} \right)^{k-n}}{\frac{i}{100}} \left[1 - \left(1 + \frac{i}{100} \right)^{j-k} \right] \right\}$$

$$8.2 \ PV_k = PMT \times 100 / i \{ 1 - (1 + i / 100)^{k-n} \}$$

9. ACCUMULATION & MEANS & σ

$$9.1 \quad \text{Sum} = \sum_{i=1}^n x_i$$

$$9.2 \quad \text{Sum of Squares} = \sum_{i=1}^n x_i^2$$

$$9.3 \quad \text{Mean} = \frac{1}{n} \sum_{i=1}^n x_i$$

$$9.4 \quad \sigma = \left\{ \frac{1}{n-1} \sum_{i=1}^n x_i^2 - \text{Mean}^2 \right\}^{1/2}$$

10. TREND LINE

$$10.1 \quad \text{Slope} = \frac{2 \sum_{k=1}^n k y_k - (n+1) \sum_{k=1}^n y_k}{n(n^2-1)/6}$$

$$10.2 \quad Y \text{ Intercept} = \frac{1}{n} \sum_{k=1}^n y_k - \frac{n+1}{2} \cdot \text{slope}$$

$$10.3 \ y_{(k)} = mk + c$$

11. SUM OF DIGITS DEPRECIATION

n = Depreciable Lifetime

PV = Initial Value of Asset

$$5 \quad 11.1 \text{ Depreciation at time } (k) = [2 \text{ PV} / n (n+1)] (n-k+1)$$

$$11.2 \text{ Remaining book value at } (k) = \text{PV} [(n-k)(n-k+1)] / n (n+1)$$

12. CASH FLOW

$$10 \quad 12.1 \text{ Current Sum of Present Value of Cash Flow}$$

$$= \sum_{j=0}^n P_j (1+i)^{-j}$$

15

where

$j = j^{\text{th}}$ cash flow

and i = cost of capital

20 Note 1: FIG. 32 illustrates the algorithm used for the solution of 2.3, 2.4 and 3.1 above. The technique is a simple Newton-Raphson method for the solution of an implicit equation.

25 Note 2: FIG. 33 shows how the price of a bond is calculated, and FIG. 34 illustrates the algorithm used to compute the yield to maturity of a bond.

Note 3: FIG. 35 illustrates the date algorithm. The first half computes the dates difference and the next 30 half the date $\pm n$ days.

OPERATING INSTRUCTIONS

All of the operations described below are controlled or initiated from the keyboard input unit 12 which is 35 shown in FIG. 1.

BASIC INSTRUCTIONS

Clearing

To clear display only

press CLX CLEAR

To clear everything (except constant storage)

press CLX CLX

Constant Storage

To store a constant

press STO

To recall a constant

press RCL

NOTE: Certain important pre-programmed calculations overwrite previous contents of the constant storage. These are:

- 55 Add-on to annual Percentage Rate Conversion
- Effective yield of an annuity (Loan repayment and sinking fund)
- Accrued interest and discounted note problems
- 60 Trend lines (least squares linear regression)
- Sum-of-the-digits calculations
- Bond calculations (price and yield)
- Accumulated interest paid on a loan
- Discounted Cash Flow Analysis

65 Except where noted above, a constant remains in the machine until it is turned off or over-written by another constant.

Rounding

To round-off (the display only) press RND

then any desired numeral key between **0** and **6**. A numeral key greater than **6** will put display in so-called "scientific notation." Normal turn-on mode is automatically rounded to two decimal places.

NOTE: Rounding affects the display only. The full internal accuracy of the machine is maintained.

Arithmetic Operations

To perform simple arithmetic operations between two numbers:

—key in the first number	press	SAVE ↑
—key in the second number, press desired operation		+ , − , × , ÷

To perform chain calculations, only the first number has to be loaded through a **SAVE ↑** operation. All subsequent numbers need only be keyed in and the desired function key pressed after each one.

Automatic computation between a displayed number and a stored constant is achieved by pressing **RCL** and the desired function.

Changing Sign

To change the sign of a displayed number	press	CHS
To enter a negative number, key in number	press	CHS

Raising a Number to a Power
Key in positive base number
(to be raised to a power)
Key in power (exponent)

press	SAVE ↑
press	y^x

To Obtain Square Root of a Number

Key in Number	press	√x y^x
---------------	-------	-----------------------------------

Percentage Operations

To obtain the percent amount of a number:

—key in the base number	press	SAVE ↑
—key in the percent (as a %)	press	%

To add or subtract the percentage amount to the base number simply press **+** or **−**, respectively.

To obtain the percent difference between the two numbers:

—key in the base (or reference) number	press	SAVE ↑
—key in the second number (answer is displayed in percent)	press	Δ% %

Calendar Functions

Data entry sequence is: month, decimal point, two numeral day and four numeral year. Example: May 8, 1972 = 5.081972 Calendar range is from Jan. 1, 1900 to December 31, 2099.

To obtain difference between two dates:

—key in first date	press	SAVE ↑
—key in second date	press	DAY

5 To obtain a date from a base date:

—key in the base date	press	SAVE ↑
—key in the number of days (can be positive or negative)	press	DATE DAY

10 To obtain the day of the week of a date:

—key in today's date	press	SAVE ↑
—key in the desired date	press	DAY SAVE ↑

—key in 7	press	÷
------------------	-------	----------

—key in that portion of the display left of the decimal point	press	−
—key in 7 again	press	×

20 If the date in question is beyond today, its day of the week will be today's day plus the number shown in the display.

25 If the date in question is before today, its day of the week will be today's day minus the number shown in the display.

Error Indication

An improper or illegal operation (such as dividing by zero) will result in a steady blinking display..

30 **Battery Condition (low charge indication)**

All decimals in the display indicates low battery condition. Plug into recharger.

COMPOUND INTEREST

35 **NOTE:** To use the compound interest keys (top row) simply remember to enter your known values in left-to-right sequence and then press the key which corresponds to your answer.

40

Future Value

Key in number of time periods	press	n
-------------------------------	-------	----------

Key in interest rate per time period (in %)	press	i
---	-------	----------

45

Key in present value (principal)	press	PV
----------------------------------	-------	-----------

To obtain future value	press	FV
------------------------	-------	-----------

50 **NOTE:** Simple arithmetic operations may be performed prior to entering any value. Also, a mistaken last entry may be corrected by pressing **CLX**, then keying in the correct value and pressing the appropriate key.

Present Value

Key in number of time periods	press	n
-------------------------------	-------	----------

Key in interest rate per time period (in %)	press	i
---	-------	----------

Key in future value amount	press	FV
----------------------------	-------	-----------

60

To obtain present value	press	PV
-------------------------	-------	-----------

Rate of Return (growth rate)

Key in number of periods	press	n
--------------------------	-------	----------

65

Key in present (beginning) value	press	PV
----------------------------------	-------	-----------

Key in future (ending) value	press	FV
------------------------------	-------	-----------

To obtain effective rate per period (in %)	press	i
--	-------	----------

3,863,060

75

76

Number of Time Periods (for a compounded amount)

Key in interest rate per period (in %) press

Key in present (beginning) value press

Key in future (ending) value press 5

To obtain number of time periods press

-Continued

Key in future value press

To obtain number of time periods press

LOAN PAYMENT

Accrued (Simple) Interest Payment Due

Nominal Rate Converted to Effective Annual Rate

Key in number of time periods per year press

Key in Nominal Rate (as a %) press

Key in press

To obtain effective annual rate (in %) press

Effective Annual Rate Converted to Nominal Rate

Key in number of time periods per year press

Key in press

Key in effective annual rate (in %) press

To obtain nominal rate (in %) press

SINKING FUND

Future Value of an Annuity (sinking fund)

Key in number of time periods press 35

Key in interest rate per period (as a %) press

Key in payment (installment) amount press 40

To obtain future value press

Key in number of days press

Key in annual interest rate (in %) press

Key in the principal (present value) press

To obtain interest payment due on a 360 day basis press

To obtain interest payment due on a 365 day basis press

Sinking Fund Payment Amount

Key in number of time periods press 45

Key in interest rate per period (as a %) press

Key in future value press 50

To obtain payment amount press

Discounted Note and Effective Annual Yields

Key in number of days press

Key in annual interest (discount) rate (in %) press

Key in the face (future) value of note press

To obtain the discount amount (i.e.—the interest portion) of the note on a 360 day basis press

Effective Yield of Sinking Fund

Key in number of time periods press 55

Key in payment (installment) amount press

Key in future value press 60

To obtain interest rate per period (as a %) press

To obtain the effective annual yield on a 360 day basis press

To obtain the discount amount of the note on a 365 day basis press

To obtain the effective annual yield on a 365 day basis press

Number of Periods Required for a Sinking Fund

Key in interest rate per period (as a %) press 65

Key in payment (installment) amount press

True Equivalent Annual Yield

Key in number of days press

Key in press

Key in the principal (present value) of note press

Key in the face value (future value) of note

press

FV

To obtain true equivalent annual yield

press

i

Present Value of an Annuity (Principal Amount of a Loan)

Key in the number of time periods (months, years, etc.)

press

n

Key in interest rate per period (in %)

press

i

Key in the amount of the payment per period

press

PMT

To obtain present value (principal)

press

PV

Loan Repayment Amount

Key in number of time periods

press

n

Key in interest rate per period (in %)

press

i

Key in present value (principal)

press

PV

To obtain payment amount per period

press

PMT

True Interest Rate of a Loan

Key in number of time periods

press

n

Key in payment amount per period

press

PMT

Key in present value (principal)

press

PV

To obtain interest rate per period (in %)

press

i

NOTE: To obtain an annual rate simply key in the number of time periods per year and press **X**.

Number of Time Periods Required for a Loan

Key in the interest rate per time period

press

i

Key in the payment amount per time period

press

PMT

Key in the present value (principal)

press

PV

To obtain the number of time periods

press

n

Accumulated Interest Paid on a Loan (between two points in time)

Key in the payment number corresponding to the first point of the time span in question

press

STO

Key in the payment number corresponding to the last point of the time span in question

press

n

Key in the total number of payments of the loan

press

n

Key in the interest rate per payment (or period)

press

i

5 Key in the payment amount per period

press

PMT

To obtain the accumulated interest

press

 $\Sigma +$

Remaining Balance (Principal) of a Loan

10 As an extension of the previous problem To obtain the remaining balance (principal)

press

X Y

"Add-on" Interest Converted to a True Annual Percentage Rate

Key in the number of months of the loan

press

n

15 Key in the "add-on" rate (per annum)

press

i

To obtain true annual percentage rate

press

i

To obtain the monthly payment amount

press

XY

20 Then key in the principal amount to be loaned

press

X

Interest Rebate (Rule of 78's)

Key in last payment number

press

n

25 Key in total number of payments for the loan

press

n

Key in the total finance charge

press

PV

To obtain the unearned interest (rebate)

press

SOD

XY

30 To obtain the remaining principal due, key in the amount of each payment remaining

press

SAVE ↑

35

press

X

XY

—

DEPRECIATION AMORTIZATION

Sum-of-the-years' Digits Depreciation

40 1. Key in given year number (or beginning year number)

press

n

2. Key in life of asset (number of years)

press

n

3. Key in depreciable amount (purchase less salvage value)

press

PV

45 4. To obtain given year's depreciation

press

SOD

5. To obtain subsequent year's depreciation

press

SOD

6. Continue step 5 as desired
7. To obtain the depreciation for a particular year not in sequence, simply key in the year number desired and

press

n

SOD

50 8. Continue step 7 as desired.

NOTE: To obtain the remaining book value after each year's depreciation press **XY**. The **XY** key must be pressed again before the next **SOD** calculation (step 5).

Straight Line Depreciation

Key in depreciable amount (purchase less salvage value)

press

SAVE ↑

SAVE ↑

To obtain each year's depreciation

—key in life of asset (number of years)

press

÷

NOTE: To obtain the remaining book value after each year's depreciation, first press **STO** **—** (for book value after first year) then **RCL** **—** for each subsequent year.

Variable Rate, Declining-Balance Depreciation

1. Key in **1 0 0** and
2. Key in life of asset (number of years)
3. Key in declining factor or rate (i.e.—1.5, 2 etc.)
4. Key in depreciable amount (purchase less salvage value)
5. To obtain year's depreciation
6. To obtain remaining book value
7. Continue steps 5. and 6. for subsequent years.

press **SAVE ↑**
 press **÷**
 press **×** **STO**
 press **RCL** **%**
 press **—**

Diminishing Balance Depreciation

1. Key in life of asset (number of years)
2. Key in beginning value of asset
3. Key in ending (salvage) value of asset

press **n**
 press **PV**
 press **FV**

NOTE: Salvage value must be greater than zero.

25 place accuracy for most yield calculations.

4. To obtain and store rate of depreciation
5. Key in beginning value of asset
6. To obtain year's depreciation
7. To obtain remaining book value
8. Continue steps 6. and 7. for subsequent years.

press **i** **CHS** **STO**
 press **RCL** **%**
 press **—**

BONDS

Price of a Bond

1. Key in either maturity or purchase date
2. Key in remaining date
3. Key in yield-to-maturity (as a %)
4. Key in coupon rate (as a %)
5. To obtain bond price

press **SAVE ↑**
 press **DAY**
 press **i**
 press **PMT**
 press **BOND** **PV**

Yield-to-Maturity of a Bond

1. Key in either maturity or purchase date*
2. Key in remaining date
3. Key in coupon rate (as a %)
4. Key in the bond price
5. To obtain bond yield

press **SAVE ↑**
 press **DAY**
 press **PMT**
 press **PV**
 press **YTM** **i**

- a) Determine the number of days, months and years to maturity (in accordance with trade custom)

40

- b) Key in number of days

press **SAVE ↑**

- c) Key in **3 0** (days/month)

press **—**

- d) Key in number of months

press **—**

- e) Key in **1 2** (months/year)

press **—**

- f) Key in number of years

press **—**

- g) Key in **3 6 5** (days/year)
Continue with step 3. of bond calculations.

press **Y** **a**

50

INVESTMENT ANALYSIS

Discounted Rate of Return (for even cash flows)

- Key in number of time periods

press **n**

- Key in amount of cash flow per period

press **PMT**

- Key in original investment

press **PV**

- To obtain discounted rate of return (in %) per period

press **i**

Discounted Cash Flow Analysis (for uneven cash flows)

1. Key in discount rate (in %) per period

press **i**

2. Key in original investment

press **CHS** **PV**

3. Key in cash flow per period
4. Continue step 3. for subsequent flows.





press **PV** **Σ+**



NOTE: Normal accuracy of bond calculations is two decimal places for most cases. If further accuracy is required, the following procedure should replace steps 1 and 2 of either of the above.



NOTE: Investment is profitable (to the extent of the discount rate) if the result is positive. Furthermore, the user can determine the "break-even" period by noting the period in which step 3 first yielded a positive result.

STATISTICS



Mean and Standard Deviation

1. Clear the entire machine by pressing  CLEAR 
2. Key in data item press 
3. Continue step 2. until all data are entered.
4. To obtain mean (arithmetic average) press 




NOTE: To obtain the standard deviation after each mean calculation press . The  key must be pressed again before resuming.


5. To return to the summation mode press  $\rightarrow \Sigma$ 
6. Continue with step 2. if desired.





NOTE: To correct a data item key in its value and press




 $\Sigma -$ 





Trend Lines (Least Squares Linear Regression)

1. Clear the entire machine by pressing  CLEAR 
2. Sequentially, key in data item press 

NOTE: Each time  is pressed, the sequence number for that item is displayed.

3. Continue step 2. until all data are entered.
4. To terminate the data entry sequence press  
5. To obtain a specific value on the trend line, key in the appropriate time period number press  
6. Repeat step 5. as often as desired.

NOTE: The user may also "step-along" the trend line by simply pressing  as many times as desired. Further, the current time period number may be obtained by pressing . The  key must be pressed again before resuming.

7. To obtain the amount of change of the trend line per period (commonly called "slope") press  
8. To resume operation press  

We claim:

1. An electronic calculator comprising: keyboard input means having a plurality of numeric keys manually operable for entering numerical data into the calculator and having a plurality of non-numeric control keys manually operable for controlling the calculator, said non-numeric con-

trol keys including a plurality of function keys associated with a plurality of mathematically related variables and manually operable for designating selected numerical data as said variables and for conditioning the calculator to perform mathematical operations involving said variables;

storage means coupled to said keyboard input means for storing numerical data entered into or calculated by the calculator;

processing means coupled to said keyboard input means and to said storage means and responsive to successive actuation of one or more of said numeric keys and one or more of said non-numeric control keys in a sequence, including one of said function keys followed by one of said function keys without interruption by any of said numeric keys, for automatically performing a mathematical operation employing selected numerical data, stored in said storage means and designated as one or more of said variables by actuation of one or more of said function keys, to determine the value of a variable associated with the last of said function keys in said sequence; and

output means coupled to said processing means for indicating the value of the variable associated with the last of said function keys in said sequence.

2. An electronic calculator as in claim 1 wherein said processing means is responsive to actuation of one of said non-numeric control keys for determining the percentage difference between two numerical values stored in said storage means and for thereupon causing said output means to display that percentage difference in decimal digit form.

3. An electronic calculator as in claim 1 wherein: said processing means is operable for generating the percentage difference between two numerical values stored in said storage means; and said non-numeric control keys include a command key manually operable with another of said non-numeric control keys, when said command key and said other non-numeric control key are successively actuated in the order mentioned, for causing said processing means to generate the percentage difference between two numerical values stored in said storage means and for thereupon causing said output means to display that percentage difference in decimal digit form.

4. An electronic calculator as in claim 1 wherein said processing means includes:

first means responsive to actuation of a first one of said non-numeric control keys for generating the arithmetic average of numerical data stored in said storage means, for causing said output means to display said arithmetic average in decimal digit form, and for generating an end-of-operation signal; and

second means coupled to said first means and responsive to said end-of-operation signal for generating the standard deviation of numerical data stored in said storage means, said processing means being responsive to actuation of a second one of said non-numeric keys for causing said output means to indicate said standard deviation in decimal digit form.

5. An electronic calculator as in claim 4 wherein said non-numeric control keys include a command key

manually operable with said first and second non-numeric control keys, when said second non-numeric control key, said command key, and said first non-numeric control key are successively actuated in the order mentioned following generation of said arithmetic average and said end-of-operation signal, for causing said first and second means to generate the arithmetic average and the standard deviation for different numerical data entered into said storage means by said numeric keys.

6. An electronic calculator as in claim 1 wherein: said numeric and non-numeric control keys are manually operable for storing numerical data representing any two calendar dates in said storage means in a predetermined decimal digit format; and

said processing means includes first means responsive to actuation of one of said non-numeric control keys for determining the number of days between any two calendar dates, included within a predetermined range of calendar dates and represented by numerical data stored in said storage means in said predetermined decimal digit format, and for causing said output means to display the determined number of days in said predetermined decimal digit format.

7. An electronic calculator as in claim 6 wherein said processing means includes:

second means responsive to actuation of said one of said non-numeric control keys for causing said output means to indicate when numerical data stored in said storage means represents an erroneous calendar date, represents a calendar date outside said predetermined range of calendar dates, or is incompatible with said predetermined decimal digit format; and

third means for compensating for the extra day in leap-years occurring within said predetermined range of calendar dates.

8. An electronic calculator as in claim 1 wherein: a first one of said non-numeric control keys is manually operable with one or more of said numeric keys for storing a first numerical datum representing the number of successive payments;

a second one of said non-numeric control keys is manually operable with one or more of said numeric keys for storing a second numerical datum representing the interest rate per payment;

a third one of said non-numeric control keys is manually operable with one or more of said numeric keys for storing a third numerical datum representing the amount of each payment; and

said processing means is responsive to actuation of a fourth one of said non-numeric control keys for generating the present value of the number of successive payments represented by said first numerical datum in the amount per payment represented by said third numerical datum and at the interest rate per payment represented by said second numerical datum and for causing said output means to display the generated present value in decimal digit form.

9. An electronic calculator is in claim 8 wherein each of said first, second, third, and fourth non-numeric control keys comprises a different one of said function keys.

10. An electronic calculator as in claim 1 wherein: a first one of said non-numeric control keys is manually operable with one or more of said numeric keys for storing equally-spaced and chronologically-sequenced numerical data in said storage means; said processing means is operable for generating the least-squares linear regression of the numerical data stored in said storage means; and

said non-numeric control keys include a command key manually operable with said first non-numeric control key, when said command key and said first non-numeric control key are successively actuated in the order mentioned, for causing said processing means to generate a first value of the least-squares linear regression of the numerical data stored in said storage means and for causing said output means to indicate said first value in decimal digit form.

11. An electronic calculator as in claim 10 wherein: said first value is the value at the y-intercept in rectangular coordinate notation; and

said processing means is responsive to further successive actuations of said first non-numeric control key for generating succeeding values of the least-squares linear regression of the numerical data stored in said storage means.

12. An electronic calculator as in claim 10 wherein: a second one of said non-numeric control keys is manually operable with one or more of said numeric keys for storing the sequence number of a chronological numerical datum in said storage means; and

said processing means is responsive to actuation of said first non-numeric control key, when successively preceded by actuation of said second non-numeric control key, for generating the value of the least-squares linear regression for the chronological numerical datum designated by the sequence number stored in said storage means by actuation of said second non-numeric control key.

13. An electronic calculator as in claim 1 wherein: said numeric and non-numeric control keys are manually operable for storing numerical data representing an initial calendar date in said storage means in a predetermined decimal digit format and for storing numerical data representing a number of days from said initial calendar date; and

said processing means includes first means responsive to actuation of one of said non-numeric control keys for generating the calendar date of the day corresponding to said number of days from said initial calendar date, when said initial and generated calendar dates are included within a predetermined range of calendar dates, and for causing said output means to display the generated calendar date in said predetermined decimal digit format.

14. An electronic calculator as in claim 13 wherein said processing means includes:

second means responsive to actuation of said one non-numeric control key for causing said output means to indicate when numerical data stored in said storage means represents an erroneous calendar date, represents a calendar date outside said predetermined range of calendar dates, or is incompatible with said predetermined decimal digit format; and

third means for compensating for the extra day in leap-years occurring within said predetermined range of calendar dates.

15. An electronic calculator as in claim 13 wherein said non-numeric control keys include a command key manually operable with said one non-numeric control key, when said command key and said one non-numeric control key are successively actuated in the order mentioned, for causing said first means to generate the calendar date of the day corresponding to said number of days from said initial calendar date and to cause said output means to display the generated calendar date in said predetermined decimal digit format.

16. An electronic calculator as in claim 13 wherein said first means is responsive to actuation of said one non-numeric control key for generating the calendar date of the day corresponding to said number of days forward from said initial calendar date, when the numerical data representing said number of days from said initial calendar date is positive, and for generating the calendar date of the day corresponding to said number of days backward from said initial calendar date, when the numerical data representing said number of days from said initial calendar date is negative.

17. An electronic calculator as in claim 1 wherein:

a first one of said non-numeric control keys is manually operable with one or more of said numeric keys for storing a first numerical datum designating a particular period within the depreciable life of an asset in said storage means and is manually operable with one or more of said numeric keys for storing a second numerical datum representing the total number of such periods in the depreciable life of the asset in said storage means;

a second one of said non-numeric control keys is manually operable with one or more of said numeric keys for storing a third numerical datum representing an initial value of the asset in said storage means;

said processing means is operable for generating the value of the sum-of-the-digits depreciation and the depreciated value of the initial value of the asset represented by said third numerical datum for the period designated by said first numerical datum; and

said non-numeric control keys include a command key manually operable with a third one of said non-numeric control keys, when said command key and said third non-numeric control key are successively actuated in the order mentioned, for causing said processing means to generate said value of the sum-of-the-digits depreciation and said depreciated value and for causing said output means to display said generated value of the sum-of-the-digits depreciation in decimal digit form.

18. An electronic calculator as in claim 17 wherein said processing means is responsive to further successive actuations of said third non-numeric control key for generating the value of the sum-of-the-digits depreciation of the initial value of the asset represented by said third numerical datum for each period included within the total number of periods represented by said second numerical datum subsequent to the period designated by said first numerical datum and for causing said output means to display each generated value of the sum-of-the-digits depreciation in decimal digit form.

19. An electronic calculator as in claim 17 wherein a fourth one of said non-numeric control keys is manually operable for thereafter causing said output means to display said generated depreciated value in decimal digit form.

20. An electronic calculator as in claim 1 wherein:

a first one of said non-numeric control keys is manually operable with one or more of said numeric keys for storing a first numerical datum representing the number of days within an interest accruing period in said storage means;

a second one of said non-numeric control keys is manually operable with one or more of said numeric keys for storing a second numerical datum representing an annual interest rate in said storage means;

a third one of said non-numeric control keys is manually operable with one or more of said numeric keys for storing a third numerical datum representing a principal amount in said storage means;

said processing means is responsive to actuation of a fourth one of said non-numeric control keys for generating the values of the amount of interest on a 360-day basis and on a 365-day basis of the principal amount represented by said third numerical datum for the number of days represented by said first numerical datum and at the annual interest rate represented by said second numerical datum and for causing said output means to display the value of the amount of interest generated on one of said bases in decimal digit form.

21. An electronic calculator as in claim 20 wherein said processing means is responsive to actuation of a fifth one of said non-numeric control keys for causing said output means to display the value of the amount of interest generated on the other of said bases in decimal digit form.

22. An electronic calculator as in claim 20 wherein said non-numeric control keys include a command key manually operable with said fourth non-numeric control key, when said command key and said fourth non-numeric control key are successively actuated in the order mentioned, for causing said processing means to generate said values of the amount of interest and to cause said output means to display the value of the amount of interest generated on said one of said bases in decimal digit form.

23. An electronic calculator as in claim 22 wherein said processing means is responsive to actuation of a fifth one of said non-numeric control keys for causing said output means to display the value of the amount of interest generated on the other of said bases in decimal digit form.

24. An electronic calculator as in claim 1 wherein:

a first one of said non-numeric control keys is manually operable with one or more of said numeric keys for storing a first numerical datum representing the number of payments within an interest accruing period in said storage means;

a second one of said non-numeric control keys is manually operable with one or more of said numeric keys for storing a second numerical datum representing the annual add-on interest rate in said storage means; and

said processing means is responsive to reactivation of said second non-numeric control key, following storage of said second numerical datum in said storage means, for generating the value of an annu-

lar percentage interest rate and the value of a monthly payment factor for the number of payments represented by said first numerical datum at the annual add-on interest rate represented by said second numerical datum and for causing said output means to display the generated value of the annual percentage interest rate in decimal digit form.

25. An electronic calculator as in claim 24 wherein each of said first and second non-numeric control keys comprises a different one of said function keys.

26. An electronic calculator as in claim 24 wherein said processing means is responsive to actuation of a third one of said non-numeric control keys for causing said output means to display the generated value of the monthly payment factor in decimal digit form.

27. An electronic calculator as in claim 26 wherein said processing means is further responsive to actuation of one or more of said numeric keys for storing a third numerical datum representing the principal amount in said storage means and is thereupon responsive to actuation of a fourth one of said non-numeric control keys for generating the value of the monthly payment and for causing said output means to display the generated value of the monthly payment in decimal digit form.

28. An electronic calculator as in claim 1 wherein: said plurality of function keys comprises five financial function keys associated with five mathematically related financial function variables; and said processing means is responsive to successive actuation of one or more of said numeric keys and one or more of said non-numeric control keys in a sequence, including one of said financial function keys followed by one of said financial function keys without interruption by any of said numeric keys, for automatically performing a mathematical operation employing selected numerical data, stored in said storage means and designated as one or more of said variables by actuation of one or more of said financial function keys, to determine the value of a financial function variable associated with the last of said financial function keys in said sequence.

29. An electronic calculator as in claim 28 wherein said five financial function keys include:

a first financial function key associated with a financial function variable representing a number of periods;

a second financial function key associated with a financial function variable representing an interest rate per period;

a third financial function key associated with a finan-

cial function variable representing a periodic payment amount;

a fourth financial function key associated with a financial function variable representing a present value of principal; and

a fifth financial function key associated with a financial function variable representing a future value of principal after one or more periods have elapsed.

30. An electronic calculator as in claim 29 wherein said first, second, third, fourth, and fifth financial function keys are positioned on said keyboard input means in a lineal sequence and in the order mentioned.

31. An electronic calculator as in claim 29 wherein: said non-numeric control keys include a command key for associating one of said five financial function keys with an alternate financial function variable representing a yield to maturity of a bond, for associating another of said five financial function keys with an alternate financial function variable representing an accrued interest amount, and for associating still another of said five financial function keys with an alternate financial function variable representing a bond price; and

each of said three last-mentioned financial function keys is manually operable, when actuated successively following actuation of said command key, for designating selected numerical data as the corresponding one of said alternate financial function variables and for conditioning the calculator to perform a mathematical operation involving that alternate financial function variable.

32. An electronic calculator as in claim 31 wherein: said first, second, third, fourth, and fifth financial function keys are positioned on said keyboard input means in a lineal sequence and in the order mentioned; and

said three financial function keys associable with said alternate financial function variables are positioned on said keyboard input means in a lineal sequence and in the order mentioned.

33. An electronic calculator as in claim 31 wherein: each of said five financial function keys is provided with an indicium indicating the financial function variable with which it is associated; and

said keyboard input means is further provided with indicia indicating the alternate financial function variables with which said three financial function keys are associable.

34. An electronic calculator as in claim 33 wherein said command key and said indicia indicating the alternate financial function variables are color coded.

* * * * *

UNITED STATES PATENT OFFICE

CERTIFICATE OF CORRECTION

Patent No. 3,863,060 Dated January 28, 1975

Inventor(s) France Rode et al.

It is certified that error appears in the above-identified patent and that said Letters Patent are hereby corrected as shown below:

Col. 11, line 12, after "pointer" insert -- \neq --;

Col. 13, line 37, "16" should read -- 20 --;

Col. 15, line 10, "wave forms" should read -- waveforms --;

Col. 18, line 66, after "with" insert -- the functions it initializes. The additional functions --;

Col. 19-22, the "TABLE OF INSTRUCTION TYPES (X = DON'T CARE)" should appear as shown below:

TABLE OF INSTRUCTION TYPES (X = DON'T CARE)

TYPE	AVAILABLE INSTRUCTIONS	NAME	FIELDS
1	256 (ADDRESSES)	JUMP SUBROUTINE	<div style="border: 1px solid black; padding: 2px; display: inline-block;"> <div style="text-align: center;">8</div> <div style="display: flex; justify-content: space-between;"> SUBROUTINE ADDRESS 0 1 </div> </div>
	256 (ADDRESSES)	CONDITIONAL BRANCH	<div style="border: 1px solid black; padding: 2px; display: inline-block;"> <div style="text-align: center;">9</div> <div style="display: flex; justify-content: space-between;"> BRANCH ADDRESS 1 1 </div> </div>
2	32 x 8 = 256	ARITHMETIC/REGISTER	<div style="border: 1px solid black; padding: 2px; display: inline-block;"> <div style="display: flex; justify-content: space-around;"> <div style="text-align: center;">5</div> <div style="text-align: center;">3</div> </div> <div style="display: flex; justify-content: space-between;"> OPERATION - WORD 1 0 </div> <div style="display: flex; justify-content: space-between;"> CODE SELECT </div> </div>
3	64 (37 used)	STATUS OPERATIONS	<div style="border: 1px solid black; padding: 2px; display: inline-block;"> <div style="display: flex; justify-content: space-around;"> <div style="text-align: center;">4</div> <div style="text-align: center;">2</div> </div> <div style="display: flex; justify-content: space-between;"> N F 0 1 0 0 </div> </div> <div style="margin-top: 5px;"> $\begin{matrix} 15 & 14 & 13 & 12 & 11 & 10 \\ I_5 & I_4 & I_3 & I_2 & I_1 & I_0 \end{matrix}$ </div> <div style="margin-top: 5px;"> F = 00 F = 01 F = 10 } F = 11 } (N = 0000) </div>
		SET BIT N	
		INTERROGATE N	
		RESET N	
		CLEAR ALL	

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Patent No. 3,863,060 Dated January 28, 1975

Inventor(s) France Rode et al.

It is certified that error appears in the above-identified patent and that said Letters Patent are hereby corrected as shown below:

TABLE OF INSTRUCTION TYPES (X = DON'T CARE)

TYPE	AVAILABLE INSTRUCTIONS	NAME	FIELDS
4	64 (30 used)	POINTER OPERATIONS	<div style="display: flex; justify-content: space-around;"> 42 </div> <div style="display: flex; align-items: center;"> <div style="border: 1px solid black; padding: 2px;">F</div> <div style="border: 1px solid black; padding: 2px;">F</div> <div style="border: 1px solid black; padding: 2px;">1</div> <div style="border: 1px solid black; padding: 2px;">1</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> </div>
		SET POINTER TO P	F = 00
		INTERROGATE P	F = 10
		DECREMENT P	F = 01
5	64 (20 used)	INCREMENT P	F = 11 } P = XXXX
		DATA ENTRY/DISPLAY	<div style="display: flex; justify-content: space-around;"> 42 </div> <div style="display: flex; align-items: center;"> <div style="border: 1px solid black; padding: 2px;">N</div> <div style="border: 1px solid black; padding: 2px;">F</div> <div style="border: 1px solid black; padding: 2px;">1</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> </div>
		LOAD CONSTANT	F = 01
		IS → A	F = 1X (N = XX01)
6	32 (11 used)	BCD INPUT TO C REG	F = 1X (N = XX11)
		STACK INSTRUCTIONS AVAILABLE	F = 10 N = (---0)
		ROM SELECT, MISC.	F = 00
		SELECT ROM "N"	<div style="display: flex; justify-content: space-around;"> 32 </div> <div style="display: flex; align-items: center;"> <div style="border: 1px solid black; padding: 2px;">N</div> <div style="border: 1px solid black; padding: 2px;">F</div> <div style="border: 1px solid black; padding: 2px;">1</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> </div>
7	8	KEYBOARD ENTRY	F = 10 (N = XX1)
		EXTERNAL ENTRY	(N = XX0)
		SUBROUTINE RETURN	F = 01 (N = XXX)
8	8	{ RESERVED FOR PROGRAM STORAGE }	<div style="display: flex; justify-content: space-around;"> 4 </div> <div style="display: flex; align-items: center;"> <div style="border: 1px solid black; padding: 2px;">X</div> <div style="border: 1px solid black; padding: 2px;">X</div> <div style="border: 1px solid black; padding: 2px;">X</div> <div style="border: 1px solid black; padding: 2px;">X</div> <div style="border: 1px solid black; padding: 2px;">1</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> </div>
9	7	{ MOS CIRCUIT }	<div style="display: flex; justify-content: space-around;"> 3 </div> <div style="display: flex; align-items: center;"> <div style="border: 1px solid black; padding: 2px;">X</div> <div style="border: 1px solid black; padding: 2px;">X</div> <div style="border: 1px solid black; padding: 2px;">X</div> <div style="border: 1px solid black; padding: 2px;">1</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> </div>
10	1	AVAILABLE	<div style="display: flex; align-items: center;"> <div style="border: 1px solid black; padding: 2px;">X</div> <div style="border: 1px solid black; padding: 2px;">X</div> <div style="border: 1px solid black; padding: 2px;">X</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> </div>
10	1	NO OPERATION (NOP)	<div style="display: flex; align-items: center;"> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> <div style="border: 1px solid black; padding: 2px;">0</div> </div>

Col. 23, line 69 "+" (each occurrence) should read --

± --;

UNITED STATES PATENT OFFICE

CERTIFICATE OF CORRECTION

Patent No. 3,863,060 Dated January 28, 1975

Inventor(s) France Rode et al.

It is certified that error appears in the above-identified patent and that said Letters Patent are hereby corrected as shown below:

Col. 23, line 70, "+" (first occurrence) should read
-- \pm --;

Col. 25, the "TABLE OF STATUS INSTRUCTION DECODING" should appear as follows:

TABLE OF STATUS INSTRUCTION DECODING

Bit #	I I I I	I I	I	I I I
	9 8 7 6	5 4	3	2 1 0
FIELD	N	F	0	1 0 0

F	INSTRUCTION
0 0	SET FLAG N
0 1	INTERROGATE FLAG N
1 0	RESET FLAG N
1 1	CLEAR ALL FLAGS (N=0000)

Col. 26, the "TABLE OF POINTER INSTRUCTION DECODING" should appear as follows:

TABLE OF POINTER INSTRUCTION DECODING

BIT #	9 8 7 6	5 4	3	2 1 0
FIELD	P	F	1	1 0 0

F	INSTRUCTION
00	Set pointer to P
10	Interrogate if pointer at P
01	Decrement pointer
11	Increment pointer

} P = XXXX
i.e. don't care

UNITED STATES PATENT OFFICE

CERTIFICATE OF CORRECTION

Patent No. 3,863,060 Dated January 28, 1975

Inventor(s) France Rode et al.

It is certified that error appears in the above-identified patent and that said Letters Patent are hereby corrected as shown below:

Col. 26, the "TABLE OF TYPE 5 INSTRUCTION DECODING" should appear as follows:

TABLE OF TYPE 5 INSTRUCTION DECODING

(X = don't care, which in this context means the instruction does not depend on this bit; either a 1 or a 0 here will cause the same execution.)

I_9	I_8	I_7	I_6	I_5	I_4	1	0	0	0
-------	-------	-------	-------	-------	-------	---	---	---	---

I_9 I_8 I_7 I_6				I_5 I_4		INSTRUCTION	
	0000 → 1111			0	0	16 Available instructions	
10	0000 \xrightarrow{N} 1001			0	1	Enters 4 bit code N into C Register at pointer position (LOAD CONSTANT)	
8	0	0	0	0	1	X	Display Toggle
	0	0	1	0	1	X	Exchange Memory, C→M→C
	0	1	0	0	1	X	Up Stack, C→C→D→E→F
	0	1	1	0	1	X	Down Stack, F→F→E→D→A
	1	0	0	0	1	X	Display OFF
	1	0	1	0	1	X	Recall Memory, M→M→C
	1	1	0	0	1	X	Rotate Down, C→F→E→D→C
	1	1	1	0	1	X	Clear all registers 0→A,B,C,D,E,F,M
1	X	X	0	1	1	X	I_8 → A register (56 bits)
1	X	X	1	1	1	X	ECD → C register (56 bits)
20							

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CERTIFICATE OF CORRECTION

Patent No. 3,863,060 Dated January 28, 1975

Inventor(s) France Rode et al.

It is certified that error appears in the above-identified patent and that said Letters Patent are hereby corrected as shown below:

Col. 27, the "TABLE OF TYPE SIX INSTRUCTION DECODING" should appear as follows:

TABLE OF TYPE SIX INSTRUCTION DECODING										Instruction
Circuit Affected	I ₉	I ₈	I ₇	I ₆	I ₅	I ₄	I ₃	I ₂	I ₁	
ROM	0	0	0	0	0	1	0	0	0	ROM select. One of 8 as specified in bits I ₉ - I ₇ .
	↓					0	0	1	0	
	1	1	1	0	0	1	0	0	0	
C&T	X	X	X	0	1	1	0	0	0	Subroutine return
	X	X	0	1	0	1	0	0	0	External key code entry to C&T
	X	X	1	1	0	1	0	0	0	Keyboard entry
DATA STORAGE	1	X	0	1	1	1	0	0	0	Send Address from C Register to Data Storage Circuit
	1	0	1	1	1	1	0	0	0	Send data from C Register into Data Storage Circuit

Col. 30, lines 25-26, "instruction" should read -- to --;

Col. 33-34, the right-hand column of the ROM 0 listing at line 132 should read -- IF A > = B[X] --;

Col. 35-36, at line 242 of the ROM 0 listing, delete "L0363:";

Col. 35-36, at line 243 of the ROM 0 listing, "I0363:" should read -- L0363: --;

Col. 41-42, at line 242 of the ROM 1 listing, "JSB ROTL" should read -- JSB ROTl --;

UNITED STATES PATENT OFFICE
CERTIFICATE OF CORRECTION

Patent No. 3,863,060 Dated January 28, 1975

Inventor(s) France Rode et al.

It is certified that error appears in the above-identified patent and that said Letters Patent are hereby corrected as shown below:

Col. 49-50, at line 53 of the ROM 3 listing, "111.1.11." should read -- 111.1.111. --;

Col. 57-58 of the ROM 4 listing, line number "130" should read line number -- 230 --;

Col. 65-66, at line 50 of the ROM 6 listing, "→L6065" should read -- →L6064 --;

Col. 67-68, at line 164 of the ROM 6 listing, ".111..11..." should read -- .111..11.. --;

Col. 67-68, at line 170 of the ROM 6 listing, "IF A>=C[M]" should read -- IF A>=C[MS] --;

Col. 69-70, at line 206 of the ROM 6 listing, "L6136:" should read -- L6316: --;

Col. 69-70, at line 214 of the ROM 6 listing, "1111..11." should read -- 11111..11. --;

Col. 69-70, at line 216 of the ROM 6 listing, "0 C[W]" should read -- 0->C[W] --;

Col. 69-70, at line 224 of the ROM 6 listing, "11.111.111." should read -- 11.111.11 --;

UNITED STATES PATENT OFFICE
CERTIFICATE OF CORRECTION

Patent No. 3,863,060 Dated January 28, 1975

Inventor(s) France Rode et al.

It is certified that error appears in the above-identified patent and that said Letters Patent are hereby corrected as shown below:

Col. 69-70, at line 239 of the ROM 6 listing, "O→C[]" should read -- O→C[X] --; and

Col. 69, line 58, between "1.3" and "=" insert -- i --.

Signed and Sealed this

thirteenth Day of *April* 1976

[SEAL]

Attest:

RUTH C. MASON
Attesting Officer

C. MARSHALL DANN
Commissioner of Patents and Trademarks